



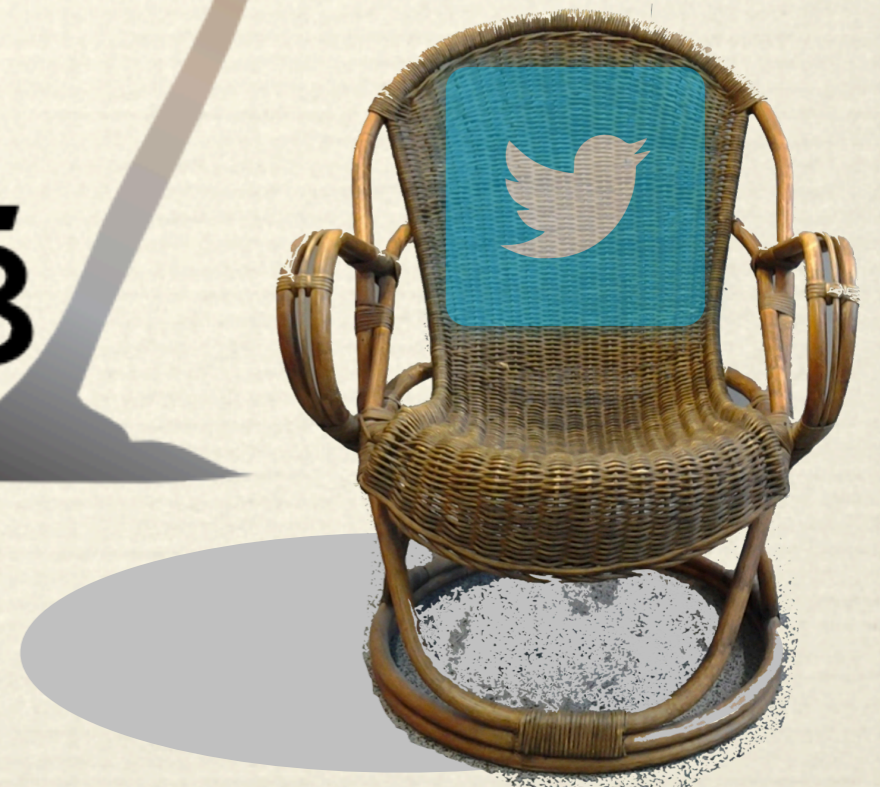
SWAT

Modeling Software Structures with GrammarLab

MoDELS 2013 Tutorial
Vadim Zaytsev, SWAT, CWI

2013

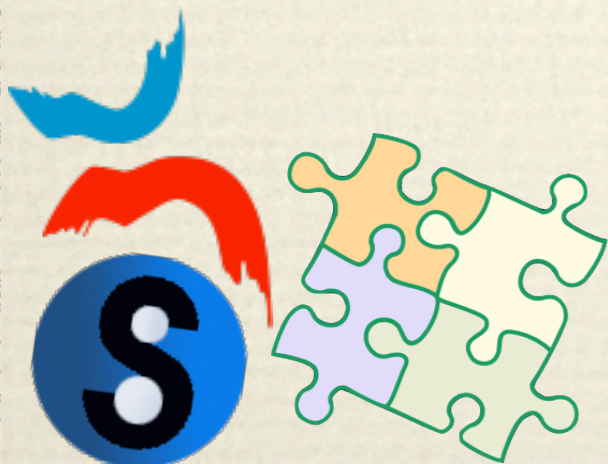
Dr. Vadim Zaytsev



Dr. Vadim Zaytsev

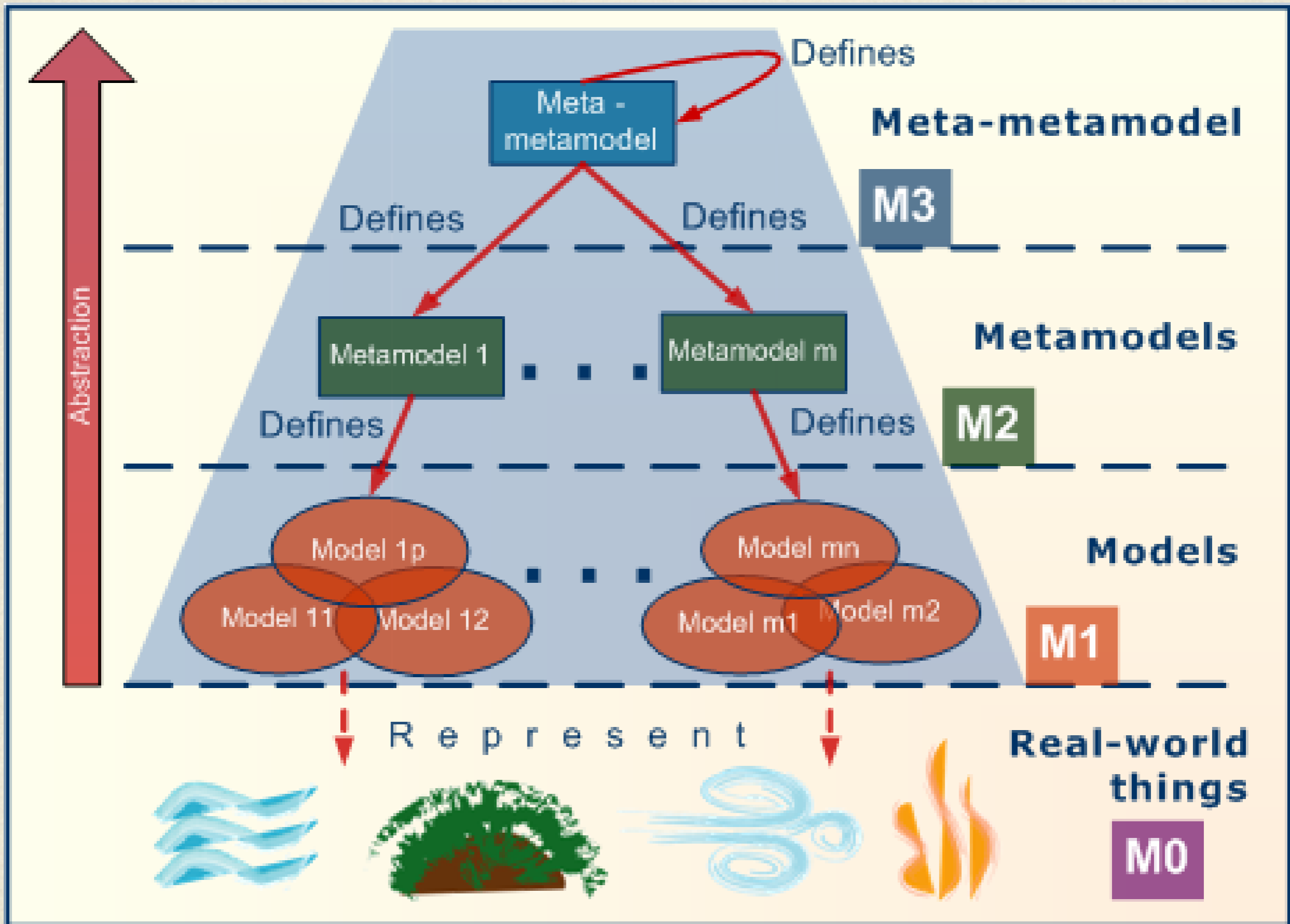


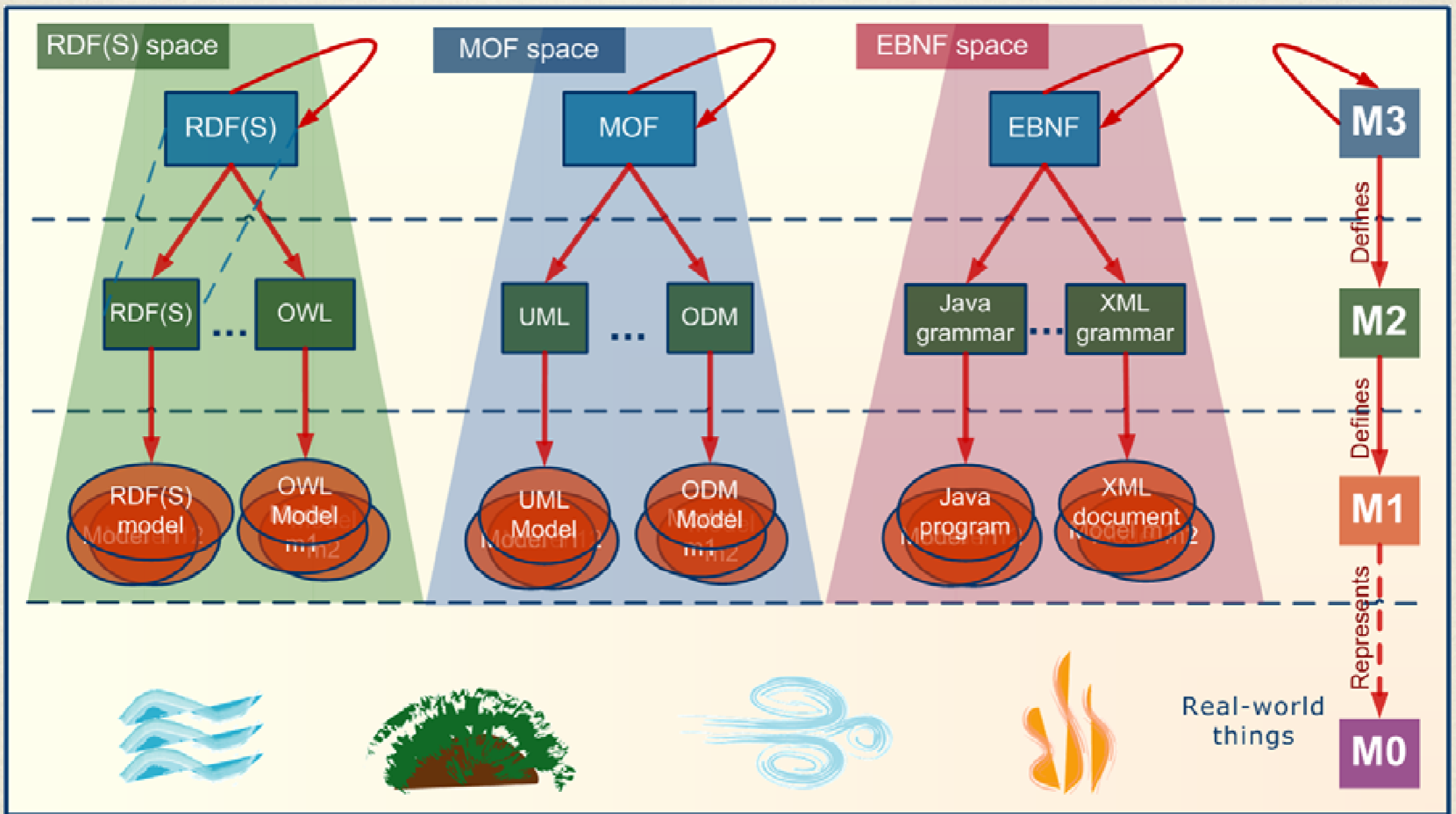
UNIVERSITÄT
KOBLENZ · LANDAU



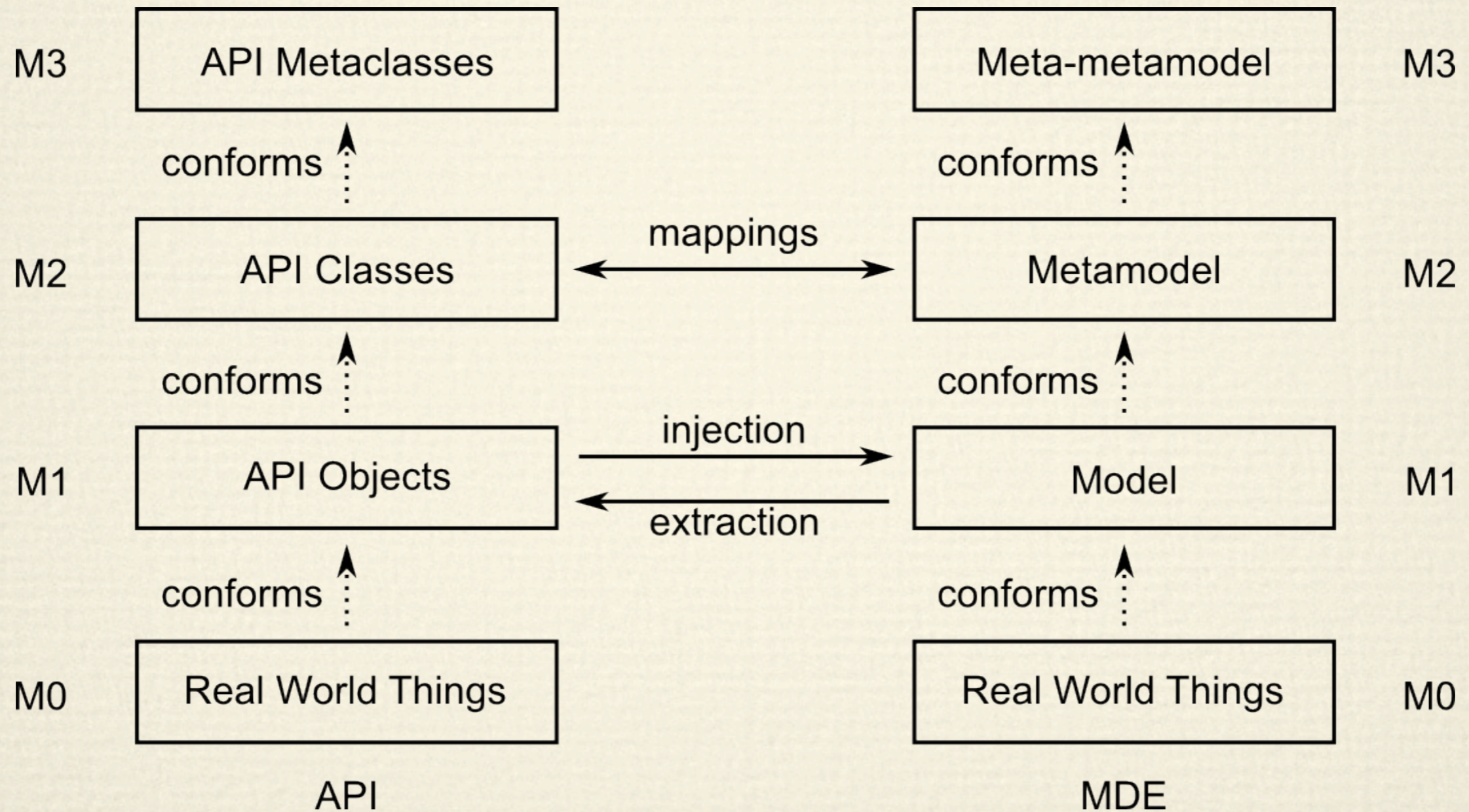
OOP
SLE
2013

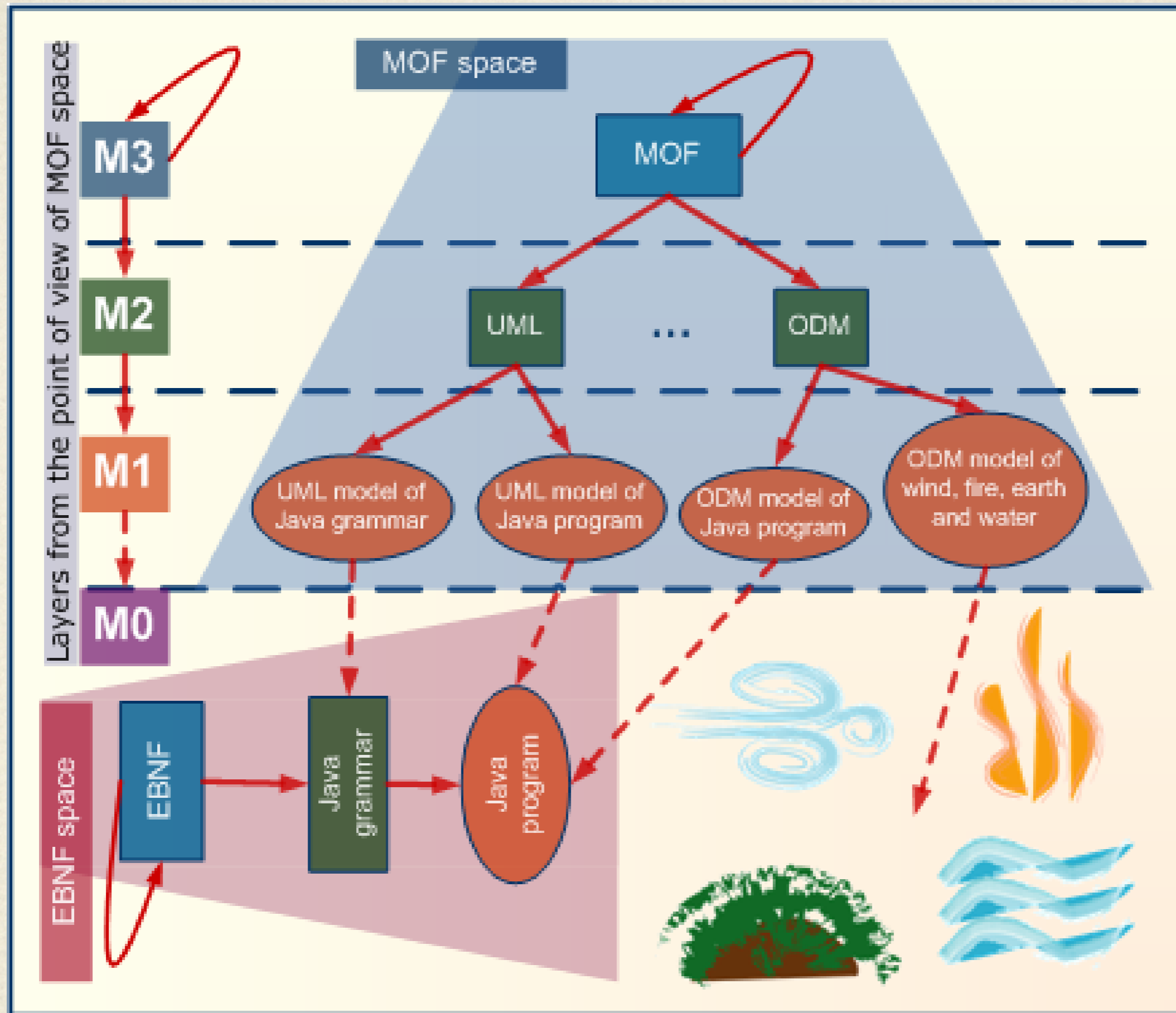






Technical side of bridging





- M. Wimmer, G. Kramler, *Bridging Grammarware and Modelware*. Workshop on Software Model Engineering at MoDELS, 2005.
- J. Bézivin, V. Devedzic, D. Djuric, J.-M. Favre, D. Gašević, and F. Jouault. *An M3-Neutral infrastructure for bridging model engineering and ontology engineering*. International Conference on Interoperability of Enterprise Software and Applications, 2005.
- D. Djuric, D. Gašević, V. Devedžic: *The Tao of Modeling Spaces*, Journal of Object Technology, 11(4), 2006.
- ...
- J. L. C. Izquierdo, F. Jouault, J. Cabot, J. G. Molina, *Automating the building of bridges between APIs and MDE*, Information & Software Technology, 54(3), 2012.
- G. Guizzardi, V. Zamborlini: *A Common Foundational Theory for Bridging Two Levels in Ontology-Driven Conceptual Modeling*. International Conference on Software Language Engineering, 2012.
- ...

To summarise



- ◆ Real world is shared

- ◆ ~~Models~~ Grammars everywhere

- ◆ Reuse across TSs

- ◆ Methodology sharing





Tutorial ::=
intro
grammars
inside
outside
rascal
grammarlab



Grammars in a broad sense

Grammars in history

- ◆ Panini (IV BC)
 - ◆ Ashtadhyayi grammar
- ◆ Chomsky (1956)
 - ◆ hierarchy of grammars
- ◆ Backus (1959), Naur (1963), Wirth (1977)
 - ◆ notation for grammars

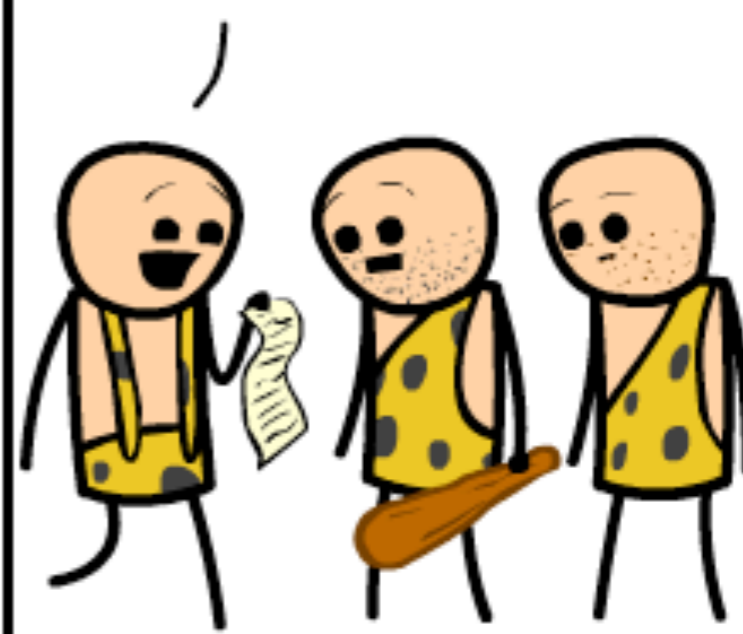
URGH! THRAK CREATE
FIRE! FIRE HOT!



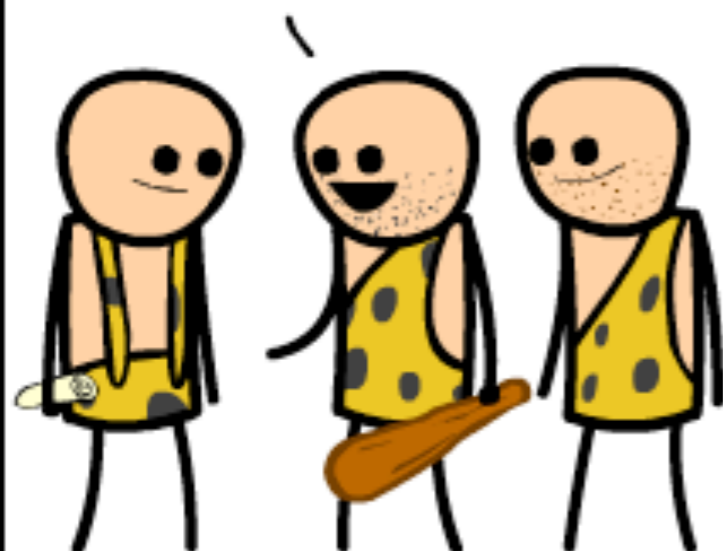
FIRE NEAT, THRAK, BUT
ME INVENT CLUB! CLUB
FOR HIT THINGS WITH.



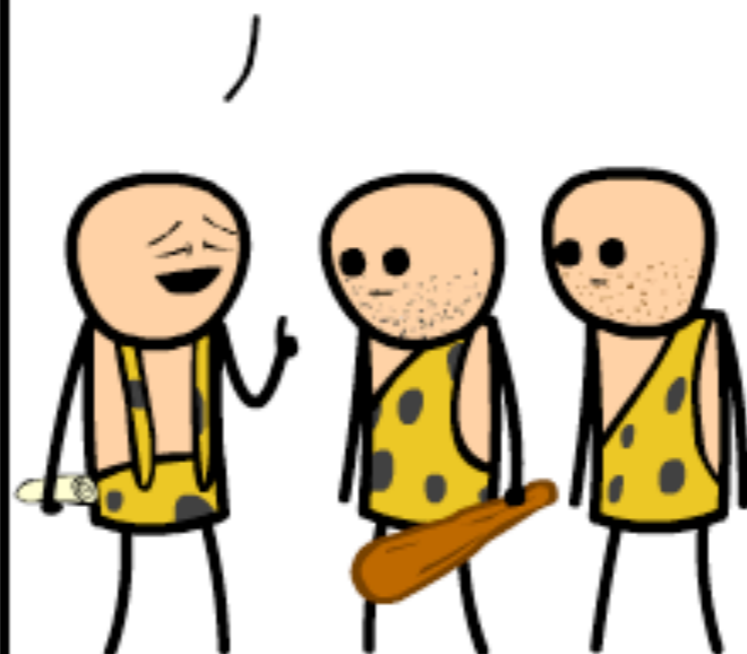
HELLO, FRIENDS! IT
SEEMS THAT I HAVE
INVENTED GRAMMAR.



THANKS! ME AND THRAK
CAN COMMUNICATE FAR
MORE EASILY NOW THAT
WE HAVE HELPING VERBS
AND ARTICLE ADJECTIVES.



"THRAK AND I"



Grammars in SE

- ◆ Definitions of languages
 - ◆ Finite definitions
 - ◆ of infinite languages
- ◆ Focus on sets of strings (“words”)
- ◆ Also, rewriting systems
- ◆ Analytic or generative

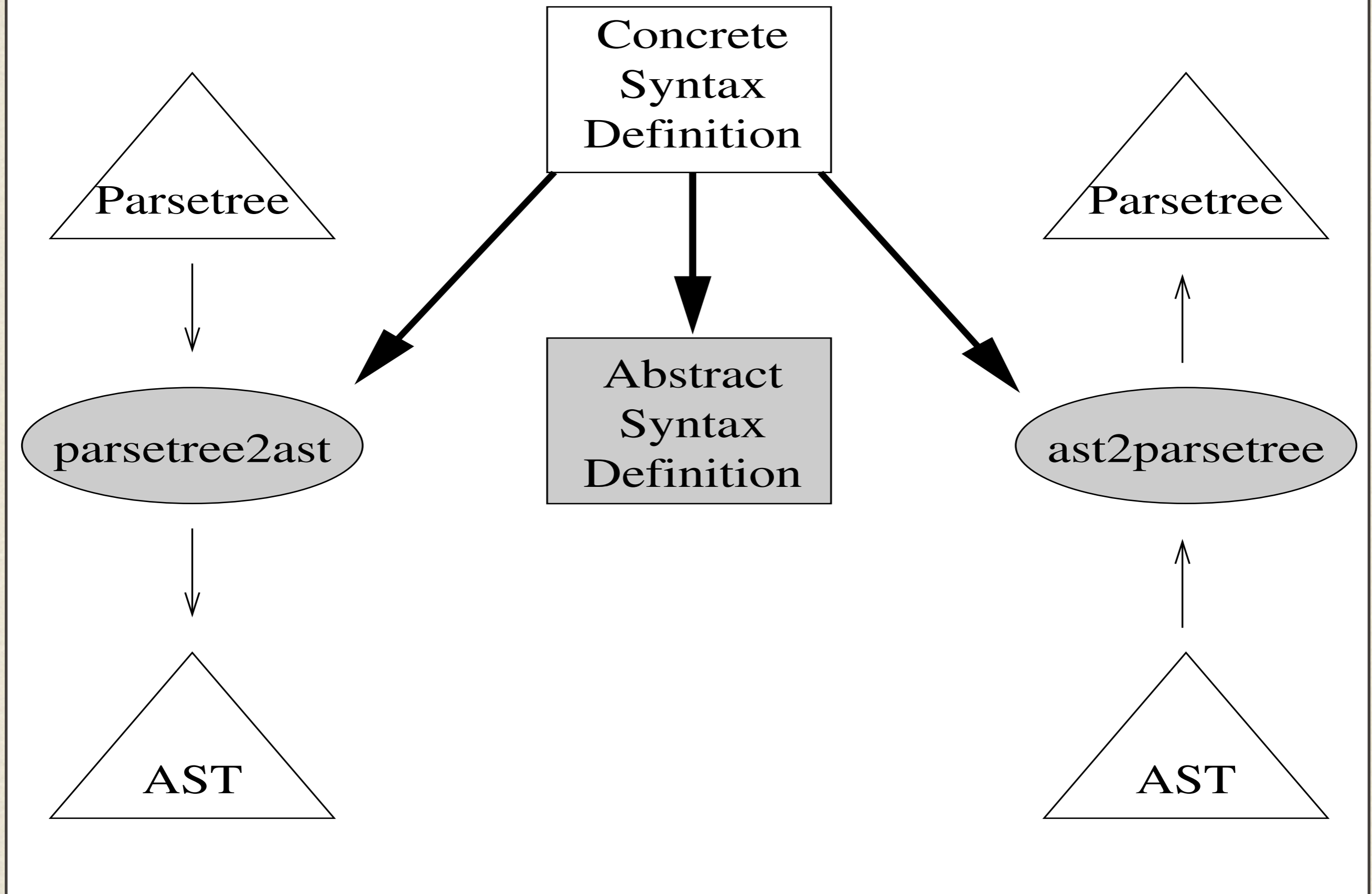
Arithmetic Expressions, Boolean Expressions, and Expressions

- $$\langle \text{factor} \rangle ::= \langle \text{number} \rangle \bar{\text{or}} \langle \text{function} \rangle \bar{\text{or}} \langle \text{variable} \rangle \bar{\text{or}} \langle \text{subscr var} \rangle$$
- $$\bar{\text{or}} (\langle \text{ar exp} \rangle) \bar{\text{or}} \langle \text{factor} \rangle \uparrow \langle \text{ar exp} \rangle \downarrow$$
- $$\langle \text{term} \rangle ::= \langle \text{factor} \rangle \bar{\text{or}} \langle \text{term} \rangle \times \langle \text{factor} \rangle \bar{\text{or}} \langle \text{term} \rangle / \langle \text{factor} \rangle$$
- $$\langle \text{ar exp} \rangle ::= \langle \text{term} \rangle \bar{\text{or}} + \langle \text{term} \rangle \bar{\text{or}} - \langle \text{term} \rangle \bar{\text{or}} \langle \text{ar exp} \rangle + \langle \text{term} \rangle$$
- $$\bar{\text{or}} \langle \text{ar exp} \rangle - \langle \text{term} \rangle$$
- $$\langle \text{ar exp A} \rangle ::= \langle \text{ar exp} \rangle$$
- $$\langle \text{relation} \rangle ::= \langle \bar{\text{or}} \rangle \bar{\text{or}} \leq \bar{\text{or}} \geq \bar{\text{or}} = \bar{\text{or}} \neq$$
- $$\langle \text{rel exp} \rangle ::= (\langle \text{ar exp} \rangle \langle \text{relation} \rangle \langle \text{ar exp A} \rangle)$$
- $$\langle \text{bool term} \rangle ::= 0 \bar{\text{or}} 1 \bar{\text{or}} \langle \text{rel exp} \rangle \bar{\text{or}} \langle \text{function} \rangle \bar{\text{or}}$$
- $$\langle \text{variable} \rangle \bar{\text{or}} \langle \text{subscr var} \rangle \bar{\text{or}} (\langle \text{bool exp} \rangle)$$
- $$\bar{\text{or}} \neg \langle \text{bool term} \rangle$$
- $$\langle \text{bool exp} \rangle ::= \langle \text{bool term} \rangle \bar{\text{or}} \langle \text{bool exp} \rangle \vee \langle \text{bool term} \rangle$$
- $$\bar{\text{or}} \langle \text{bool exp} \rangle \wedge \langle \text{bool term} \rangle \bar{\text{or}}$$
- $$\langle \text{bool exp} \rangle \equiv \langle \text{bool term} \rangle$$
- $$\langle \text{exp} \rangle ::= \langle \text{ar exp} \rangle \bar{\text{or}} \langle \text{bool exp} \rangle$$

$$\begin{array}{l}
S \rightarrow aSB \& AB \mid b \\
A \rightarrow aA \mid \varepsilon \\
B \rightarrow B_1 \mid B_2 \\
B_1 \rightarrow B_1B_3 \& B_2B_2 \mid b \\
B_2 \rightarrow B_1B_1 \& B_2B_6 \mid bb \\
B_3 \rightarrow B_1B_2 \& B_6B_6 \mid bbb \\
B_6 \rightarrow B_1B_2 \& B_3B_3
\end{array}$$

Grammars as contracts

- ◆ Many metatools
 - ◆ each implies its own language
- ◆ Interoperability
 - ◆ depends on conforming to the same language

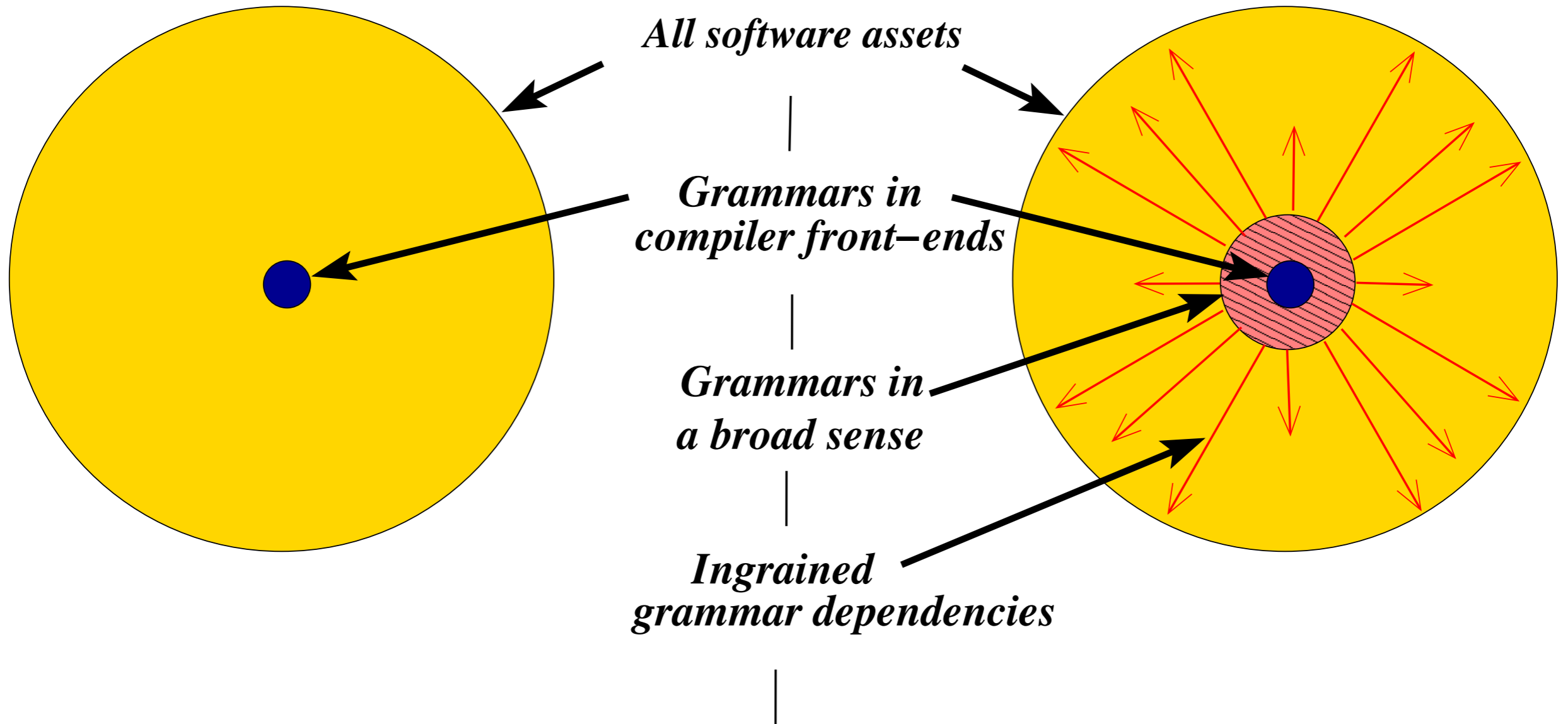


Grammars as commitments

- ◆ Intercommunication purposes
- ◆ Bridging-in-the-small
 - ◆ Assume two independent collaborators
 - ◆ Each side has its own infrastructure
 - ◆ E.g., XML and ANTLR
 - ◆ Information exchange is crucial
 - ◆ The system is evolving

Perceived view "*Lines of code*"

Proposed view "*Impact ratio*"



Grammars as interfaces

- ◆ Algebraic data types
- ◆ Domain models
- ◆ Class diagrams
- ◆ Libraries
- ◆ API
- ◆ Questionnaires

JavaScript

Browser

CoffeeScript

runs

compiles to

JavaScript

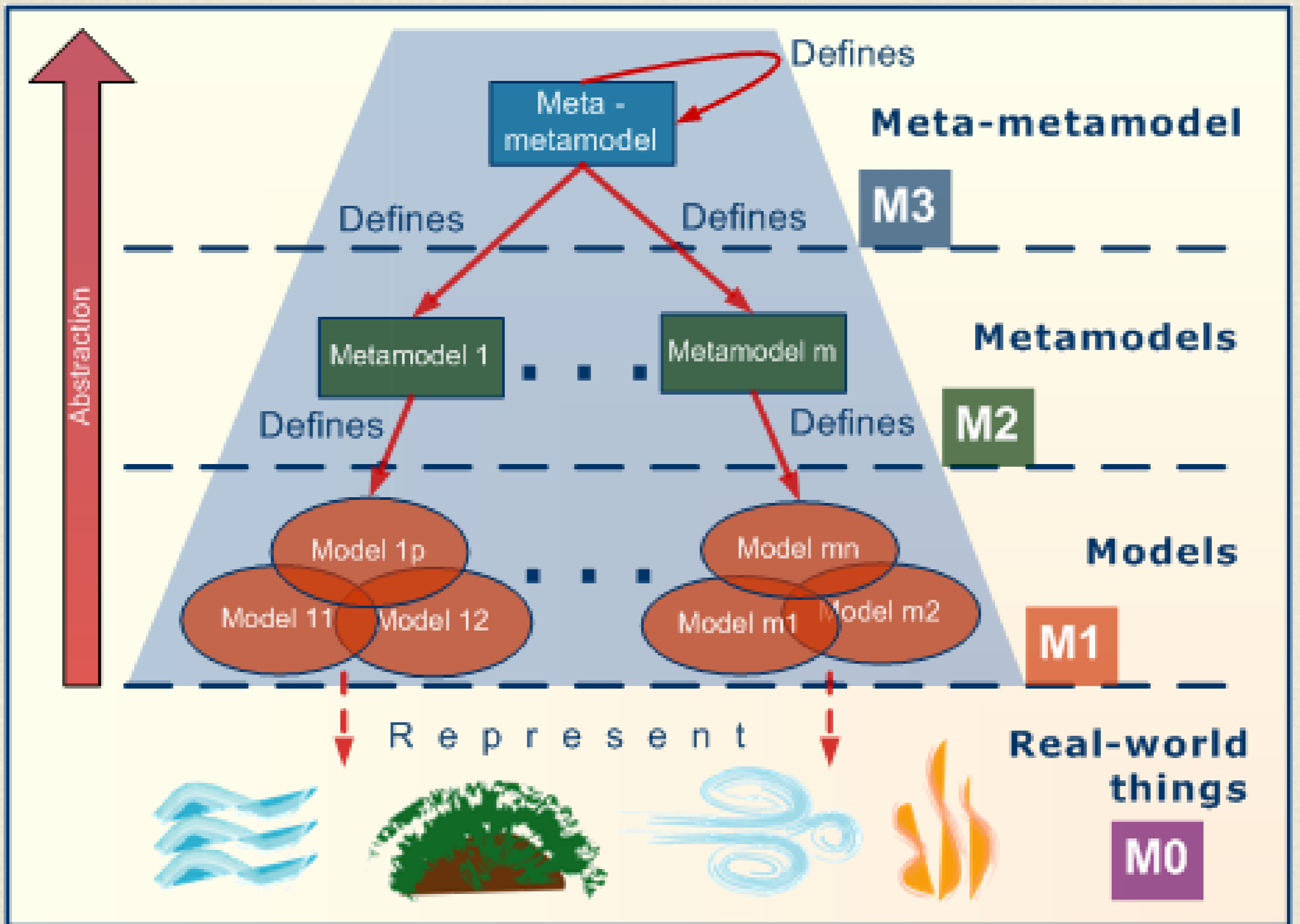
is a library for

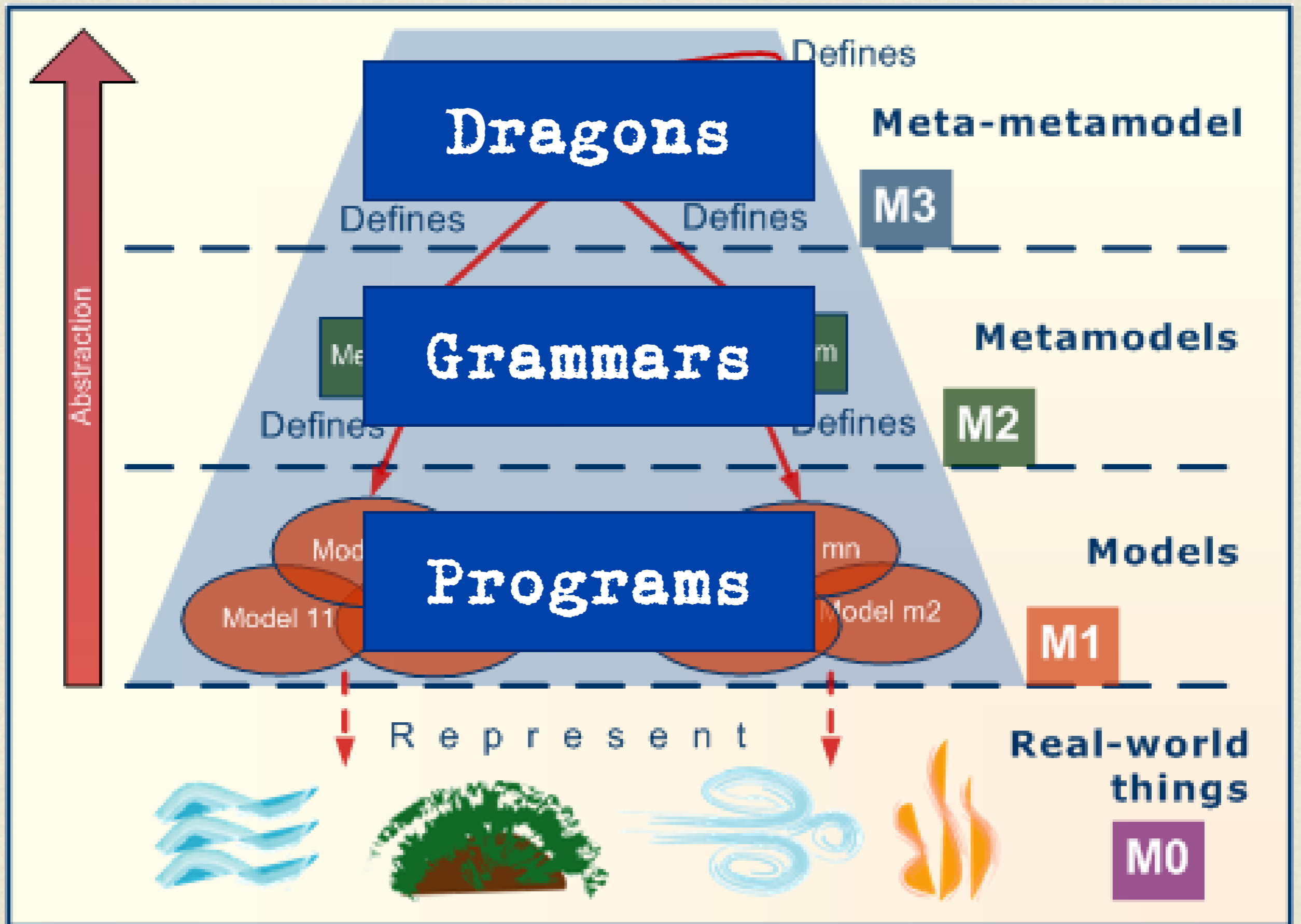
generates

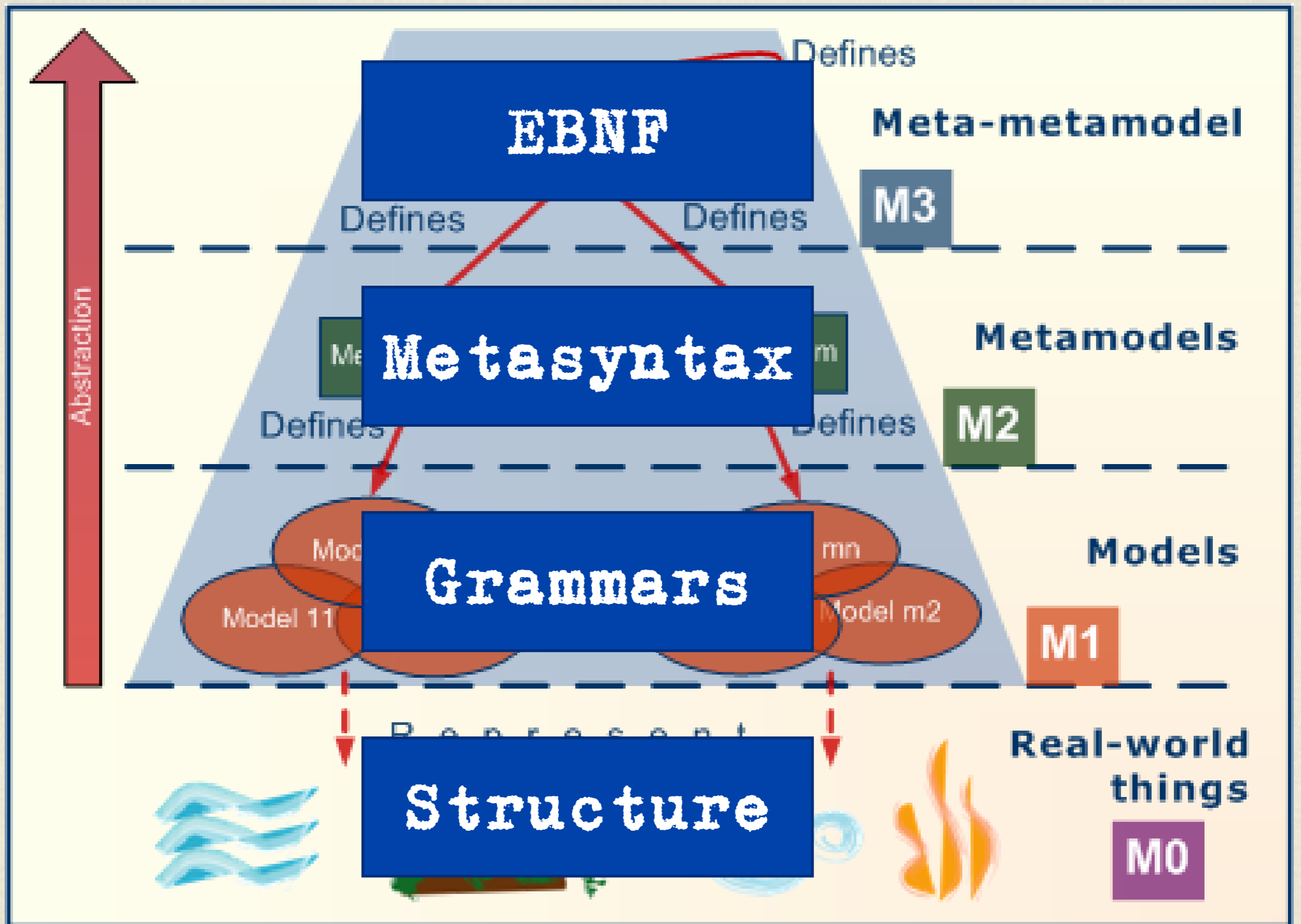
jQuery

dart2js











To summarise



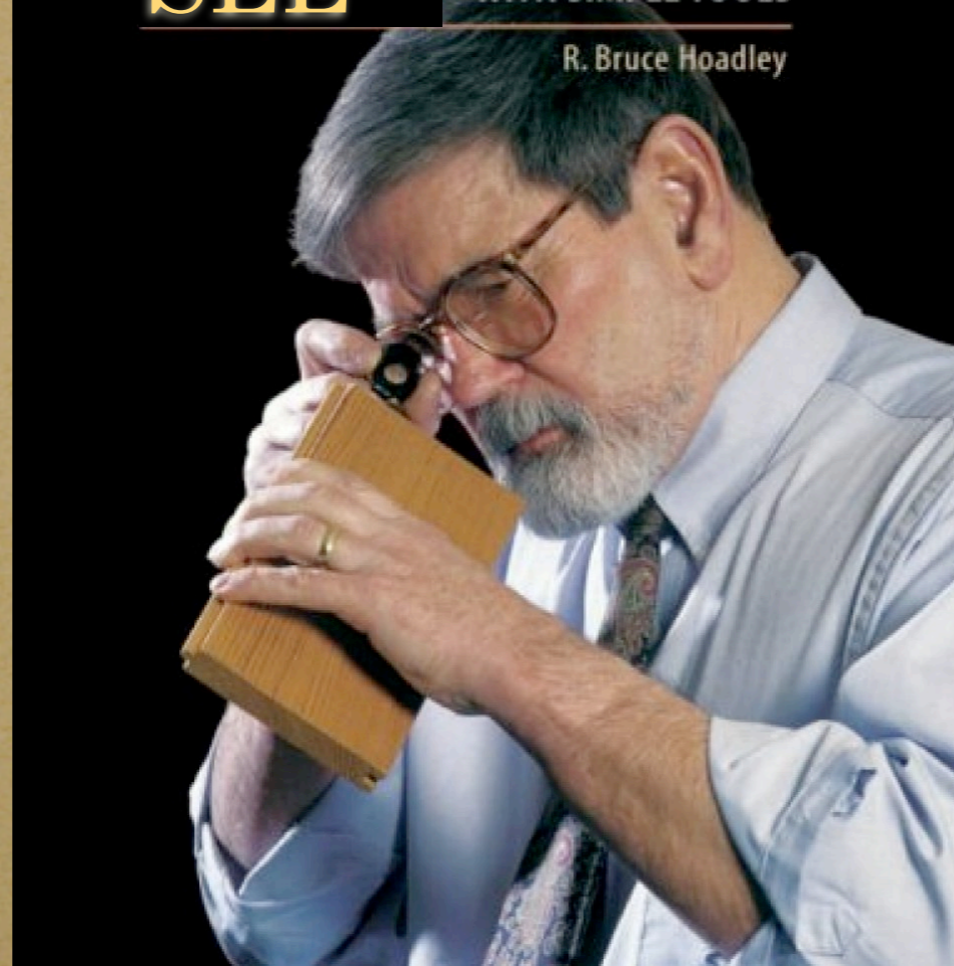
- ◆ Grammars define languages
- ◆ Represent structural commitments
- ◆ Contracts / interfaces
- ◆ Bridgebuilding material
- ◆ *Grammars in a broad sense model software languages*



**IDENTIFYING
SLE**

ACCURATE RESULTS
WITH SIMPLE TOOLS

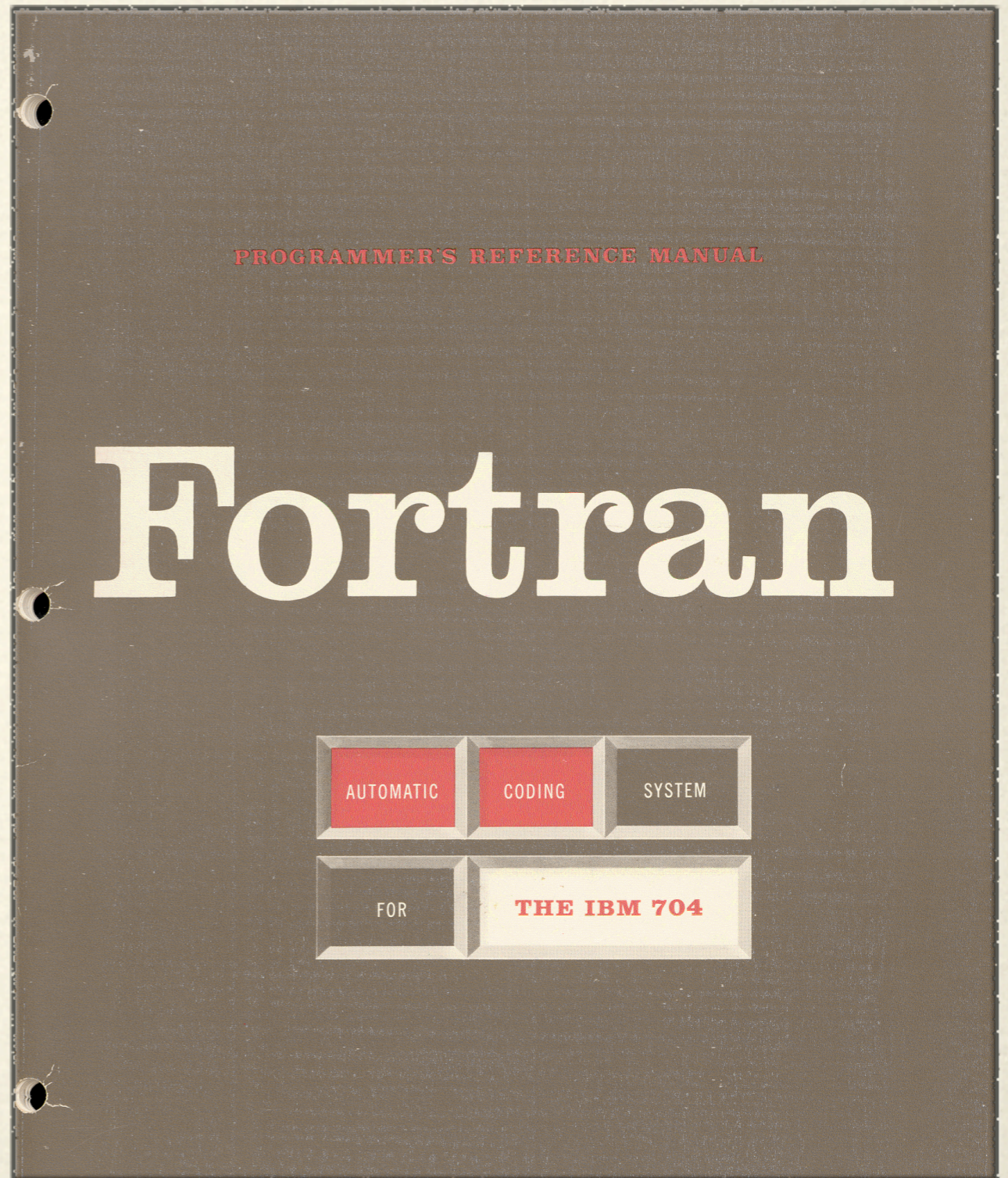
R. Bruce Hoadley



What is a software
language?

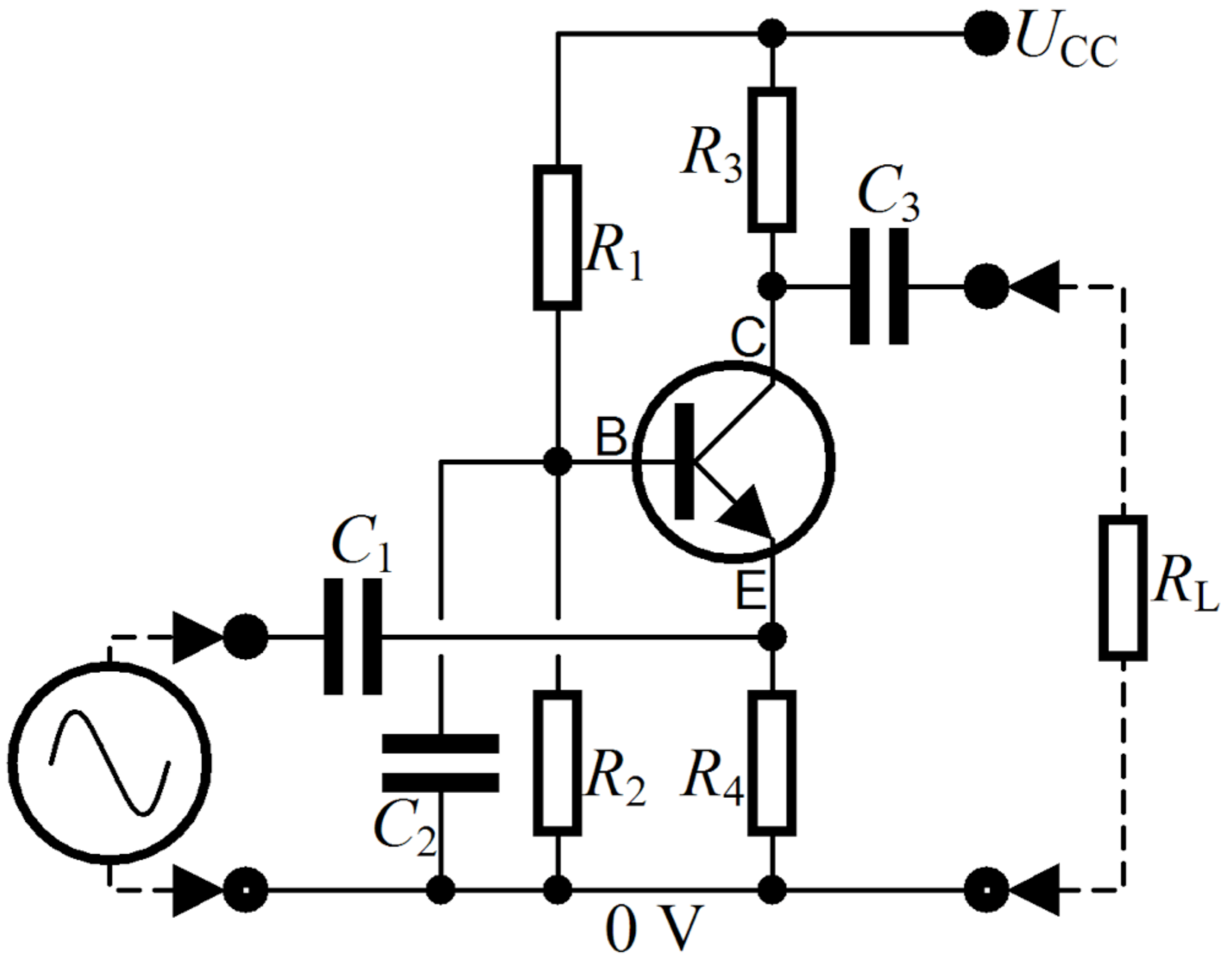
SL is GPL

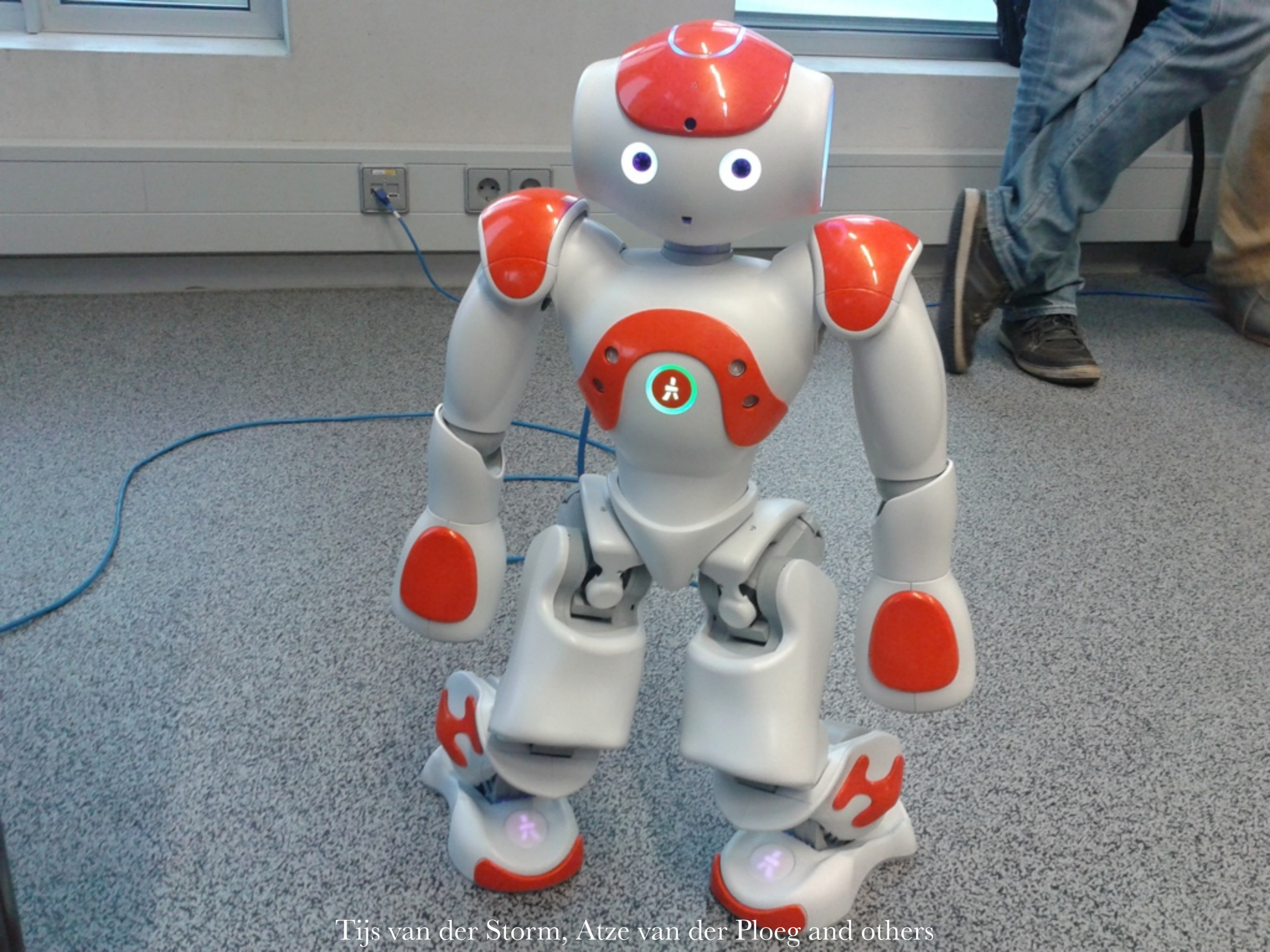
- ◆ Programming language
- ◆ General purpose
- ◆ High level constructs
- ◆ Syntactic sugar
- ◆ Can be compiled/
interpreted/computed/
evaluated/...





Hackers & Designers





Tijs van der Storm, Atze van der Ploeg and others

GET statuses/mentions_timeline

https://dev.twitter.com/docs/api/1.1/get/statuses/mentions_timeline

Developers API Health Blog Discussions Documentation Search Sign in

Home → Documentation → REST API v1.1 Tweet

GET statuses/mentions_timeline

View What links here

Updated on Thu, 2013-06-20 14:39 API version 1.1

Returns the 20 most recent mentions (tweets containing a users's @screen_name) for the authenticating user.

The timeline returned is the equivalent of the one seen when you view [your mentions](#) on twitter.com.

This method can only return up to 800 tweets.

See [Working with Timelines](#) for instructions on traversing timelines.

Resource URL

https://api.twitter.com/1.1/statuses/mentions_timeline.json

Parameters

count optional	Specifies the number of tweets to try and retrieve, up to a maximum of 200. The value of <code>count</code> is best thought of as a limit to the number of tweets to return because suspended or deleted content is removed after the count has been applied. We include retweets in the count, even if <code>include_rts</code> is not supplied. It is recommended you always send <code>include_rts=1</code> when using this API method.
since_id optional	Returns results with an ID greater than (that is, more recent than) the specified ID. There are limits to the number of Tweets which can be accessed through the API. If the limit of Tweets has occurred since the <code>since_id</code> , the <code>since_id</code> will be forced to the oldest ID available.

Resource Information

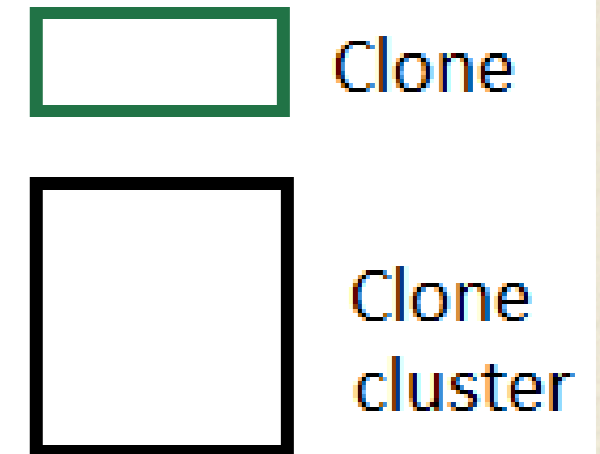
Rate Limited?	Yes
Requests per rate limit window	15/user
Authentication	Requires user context
Response Formats	json
HTTP Methods	GET
Resource family	statuses
Response Object	Tweets
API Version	v1.1

OAuth tool

Please [Sign in](#) with your Twitter account in order to use the OAuth tool.



B10		✕ ✓ f_x		0.1156		
	A	B	C	D	E	F
7						
8	Variance/Covariance Matrix:					
9		A	B	C	D	
10	A	0.1156	0.0320	0.0384	0.0448	
11	B	0.0320	0.1300	0.0480	0.0560	
12	C	0.0384	0.0480	0.1476	0.0672	
13	D	0.0448	0.0560	0.0672	0.1684	
14						



D25		✕ ✓ f_x		=B4*D\$20*\$B25				
	A	B	C	D	E	F	G	H
22								
23				Covariance Matrix				
24		<u>stdev</u>		1	2	3	4	
25		0.34	1	0.1156	0.032	0.0384	0.0448	
26		0.3606	2	0.032	0.13	0.048	0.056	
27		0.3842	3	0.0384	0.048	0.1476	0.0672	
28		0.4104	4	0.0448	0.056	0.0672	0.1684	
29								

To summarise



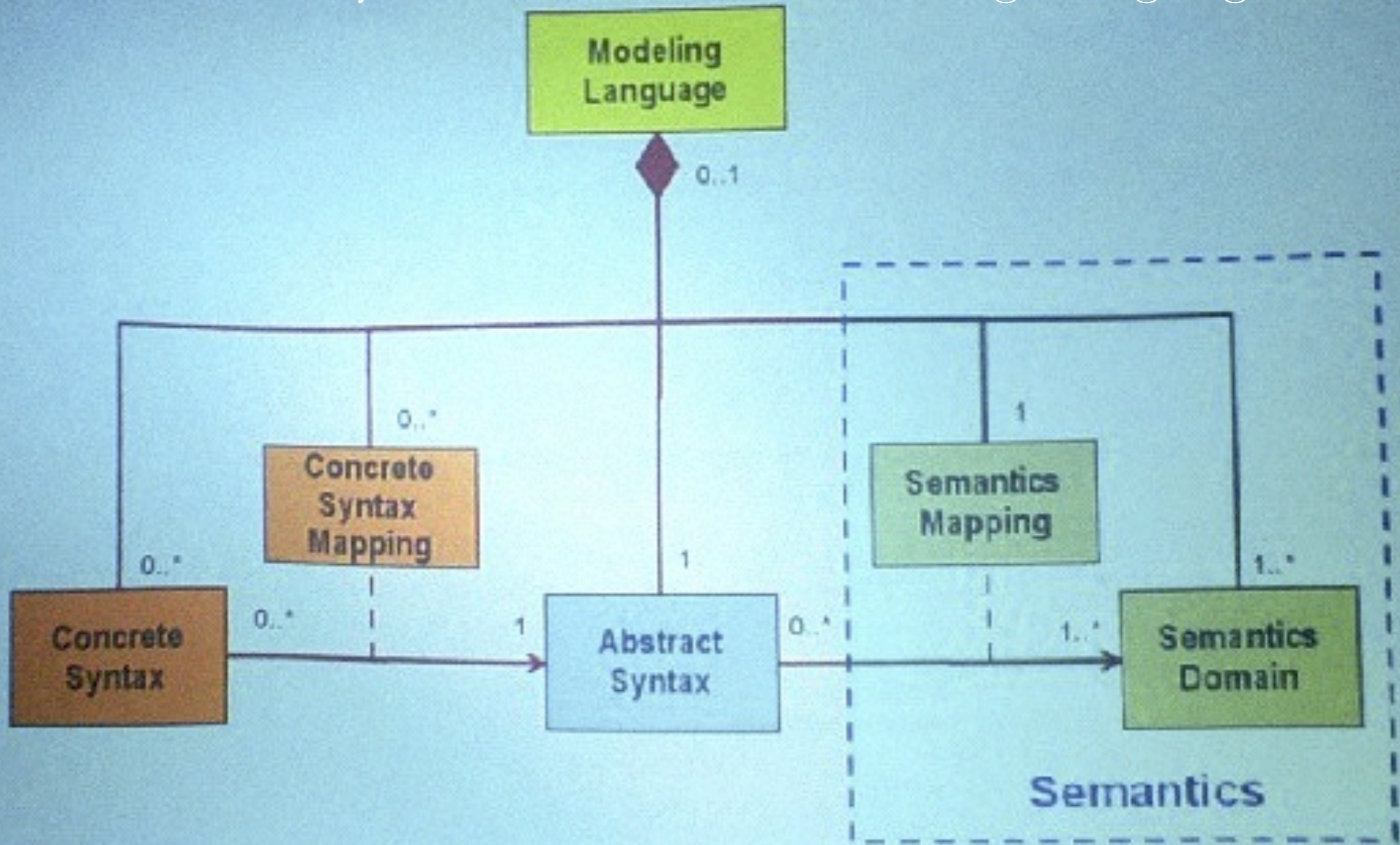
- ◆ Language is a set of instances
- ◆ Language is modelled by a grammar
- ◆ Languages are everywhere
- ◆ Not necessarily used by experts



**Abstract
&
concrete**

Elements of a Modeling Language

Bran Selic, Theory and Practice of Modeling Language Design



Ralf Lämmel, <http://instagram.com/p/e5P9r4TGHd/>

Both are grammars!

- ◆ Concrete syntax
 - ◆ elements for visualisation & programming
- ◆ Abstract syntax
 - ◆ conceptual structure
- ◆ Beyond dychotomy!

DSL concrete syntax in C#

```
public void Configure(Reader target) {
    target.AddStrategy(ConfigureServiceCall());
    target.AddStrategy(ConfigureUsage());
}
private ReaderStrategy ConfigureServiceCall() {
    ReaderStrategy result = new ReaderStrategy("SVCL", typeof (ServiceCall));
    result.AddFieldExtractor(4, 18, "CustomerName");
    result.AddFieldExtractor(19, 23, "CustomerID");
    result.AddFieldExtractor(24, 27, "CallTypeCode");
    result.AddFieldExtractor(28, 35, "DateOfCallString");
    return result;
}
private ReaderStrategy ConfigureUsage() {
    ReaderStrategy result = new ReaderStrategy("USGE", typeof (Usage));
    result.AddFieldExtractor(4, 8, "CustomerID");
    result.AddFieldExtractor(9, 22, "CustomerName");
    result.AddFieldExtractor(30, 30, "Cycle");
    result.AddFieldExtractor(31, 36, "ReadDate");
    return result;
}
```

DSL concrete syntax in XML

```
<ReaderConfiguration>
  <Mapping Code = "SVCL" TargetClass = "dsl.ServiceCall">
    <Field name = "CustomerName" start = "4" end = "18"/>
    <Field name = "CustomerID" start = "19" end = "23"/>
    <Field name = "CallTypeCode" start = "24" end = "27"/>
    <Field name = "DateOfCallString" start = "28" end = "35"/>
  </Mapping>
  <Mapping Code = "USGE" TargetClass = "dsl.Usage">
    <Field name = "CustomerID" start = "4" end = "8"/>
    <Field name = "CustomerName" start = "9" end = "22"/>
    <Field name = "Cycle" start = "30" end = "30"/>
    <Field name = "ReadDate" start = "31" end = "36"/>
  </Mapping>
</ReaderConfiguration>
```

DSL concrete syntax in text

```
mapping SVCL dsl.ServiceCall
  4-18: CustomerName
  19-23: CustomerID
  24-27 : CallTypeCode
  28-35 : DateOfCallString
```

```
mapping USGE dsl.Usage
  4-8 : CustomerID
  9-22: CustomerName
  30-30: Cycle
  31-36: ReadDate
```

POLICE TELEPHONE

FREE
FOR USE OF
PUBLIC

ADVICE & ASSISTANCE
OBTAINABLE IMMEDIATELY

OFFICER & CARS
RESPOND TO ALL CALLS

PULL TO OPEN

Grammars
on the inside

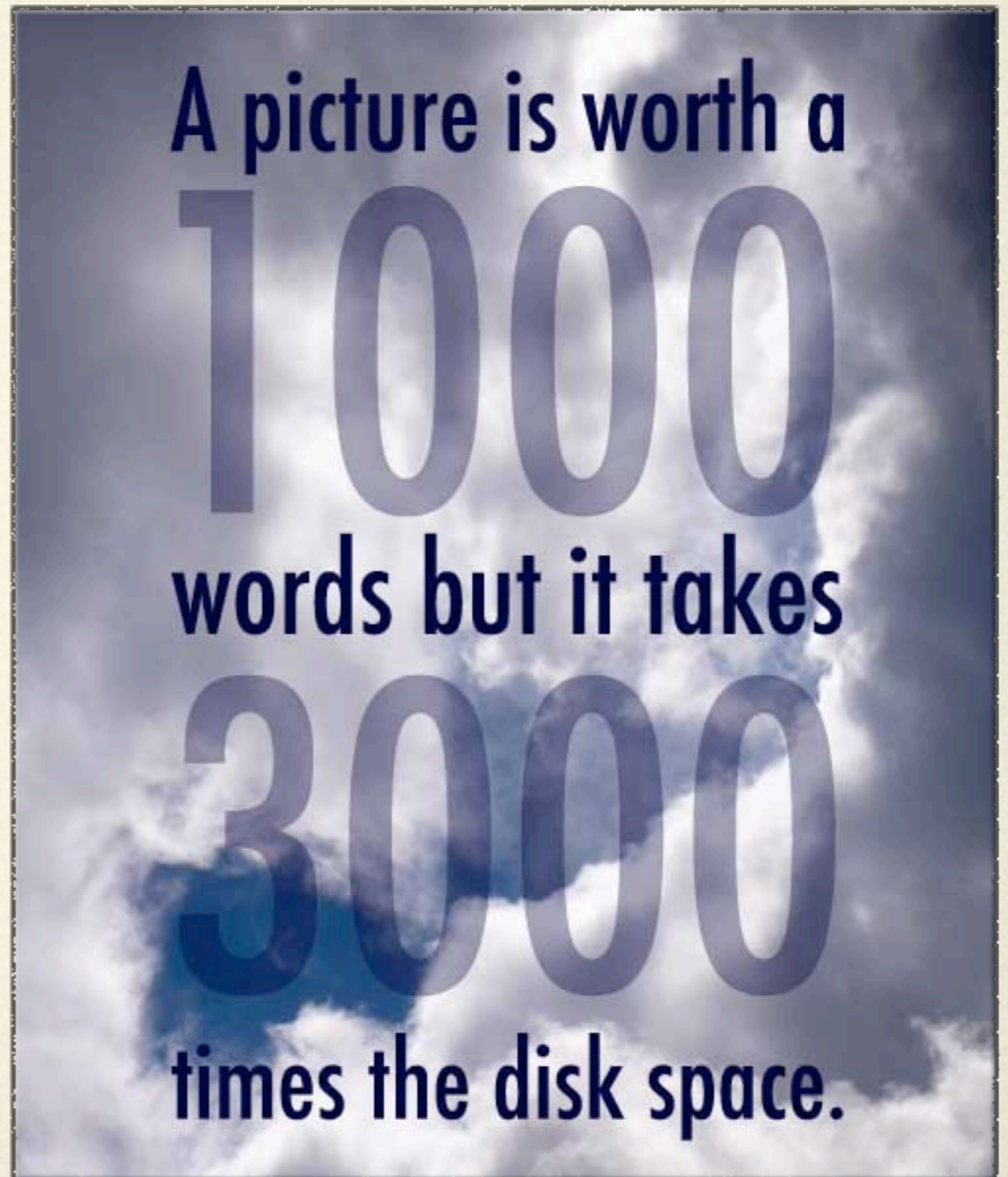
Backus Naur Form

- ◆ Historically useful
- ◆ Extended over and over
- ◆ “Backus-Wirth Notation”
 - ◆ (not a normal form)
- ◆



Backus Naur Form

- ◆ Historically useful
- ◆ Extended over and over
- ◆ “Backus-Wirth Notation”
 - ◆ (not a normal form)
- ◆ Textual



Inside a grammar

production rule

$A ::= B \text{ “C” } D? E^* F^+ (G|H);$

Inside a grammar

“left hand side”

A ::= B “C” D? E* F+ (G|H);

Inside a grammar

$A ::= B \text{ "C" } D? E^* F^+ (G|H);$

nonterminal symbol

Inside a grammar

$A ::= B \text{ “C” } D? E^* F^+ (G|H);$

terminal symbol

Inside a grammar

optional symbol

A ::= B “C” D? E* F+ (G|H);

BNE

Inside a grammar

Kleene star

A ::= B “C” D? E* F⁺ (G|H);

Inside a grammar

one or more

A ::= B “C” D? E* F⁺ (G|H);

Inside a grammar

disjunction

$A ::= B \text{ “C” } D? E^* F^+ (G|H);$

Inside a grammar

markers

$X ::= \langle a \rangle : b [c] :: d ;$

labels

Inside a grammar

$X ::= \varepsilon ;$

$X ::= \varphi ;$

$X ::= \alpha ;$



nothing

Inside a grammar

$X ::= \varepsilon ;$

$X ::= \varphi ;$

$X ::= \alpha ;$

error

A red callout box with a white border and a pointer pointing to the second production rule, $X ::= \varphi ;$. The box contains the word "error" in white, bold, lowercase letters.

Inside a grammar

$X ::= \varepsilon ;$

$X ::= \varphi ;$

$X ::= \alpha ;$



anything

Inside a grammar

$X ::= \{ Y Z \}^+ ;$

$X ::= A \& B ;$

$X ::= \neg N ;$



Y
YZY
YZYZY
...

Inside a grammar

$X ::= \{ Y Z \}^+ ;$

conjunction

$X ::= A \& B ;$

$X ::= \neg N ;$

Inside a grammar

$X ::= \{ Y Z \}^+ ;$

$X ::= A \& B ;$

$X ::= \neg N ;$



negation

Ada 95 (Magnus Kempe, HTML)

compilation:

compilation_unit*

compilation_unit:

context_clause library_item

context_clause subunit

library_item:

"private"? library_unit_declaration

library_unit_body

"private"? library_unit_renaming_declaration

library_unit_declaration:

subprogram_declaration

package_declaration

generic_declaration

generic_instantiation

library_unit_renaming_declaration:

package_renaming_declaration

generic_renaming_declaration

subprogram_renaming_declaration

<http://slps.github.io/zoo/ada/kempe.html>

BibTeX (Guillaume Hillairet, Ecore)

Bibtex:

entries::[Entry](#)+

Entry:

[Article](#)

Entry:

[Book](#)

Article:

key::[String](#) fields::[Field](#)+

Book:

key::[String](#) fields::[Field](#)+

Field:

[Authors](#)

Field:

[AuthorUrls](#)

Field:

[Title](#)

<http://slps.github.io/zoo/bibtex/bibtex-1.html>

ANSI C90 (Vinju, Kooiker, van den Brand, SDF)

TranslationUnit:
ExternalDeclaration+

ExternalDeclaration:
FunctionDefinition

ExternalDeclaration:
Declaration

FunctionDefinition:
Specifier* Declarator Declaration* "{" Declaration* Statement* "}"

Declaration:
Specifier+ (InitDeclarator (", " InitDeclarator)*) ";"

Declaration:
Specifier+ ";"

InitDeclarator:
Declarator

InitDeclarator:
Declarator "=" Initializer

GNU C grammar (TXL, Andrew J. Malton, James R. Cordy)

preprocessor:

```
"#define" id "(" id+ ")" expression NL  
"#define" id expression NL  
EX "#else" (IN NL)  
EX "#endif" (NL NL)  
NL "#if" expression (IN NL)  
NL "#ifdef" id (IN NL)  
NL "#ifndef" id (IN NL)  
"#ident" stringlit NL  
"#include" stringlit NL  
"#include" "<" SPOFF filepath ">" SPON NL  
"#line" integernumber stringlit? NL  
"#undef" id NL  
"#LINK" stringlit NL
```

filepath:

```
file_id slash_fileid*
```

file_id:

```
id  
key
```

OASIS DocBook (RELAX-NG)

bookcomponent.content:

```
divcomponent.mix+ (sect1* | refentry.class* | simplesect* | section.class*)  
sect1+  
refentry.class+  
simplesect+  
section.class+
```

local.set.attrib:

ϵ

set.role.attrib:

role.attrib

set:

```
set::(set.attlist div.title.content? setinfo? toc? book.class+ setindex?)
```

set.attlist:

```
fpi::string? status.attrib common.attrib set.role.attrib local.set.attrib
```

local.setinfo.attrib:

ϵ

<http://slps.github.io/zoo/markup/docbook-walsh.html>

To summarise



- ◆ Nonterminals
- ◆ Terminals
- ◆ Metaproperties (?, *, +)
- ◆ Special cases
- ◆ Labels and markers
- ◆ Disjunction
- ◆ ...

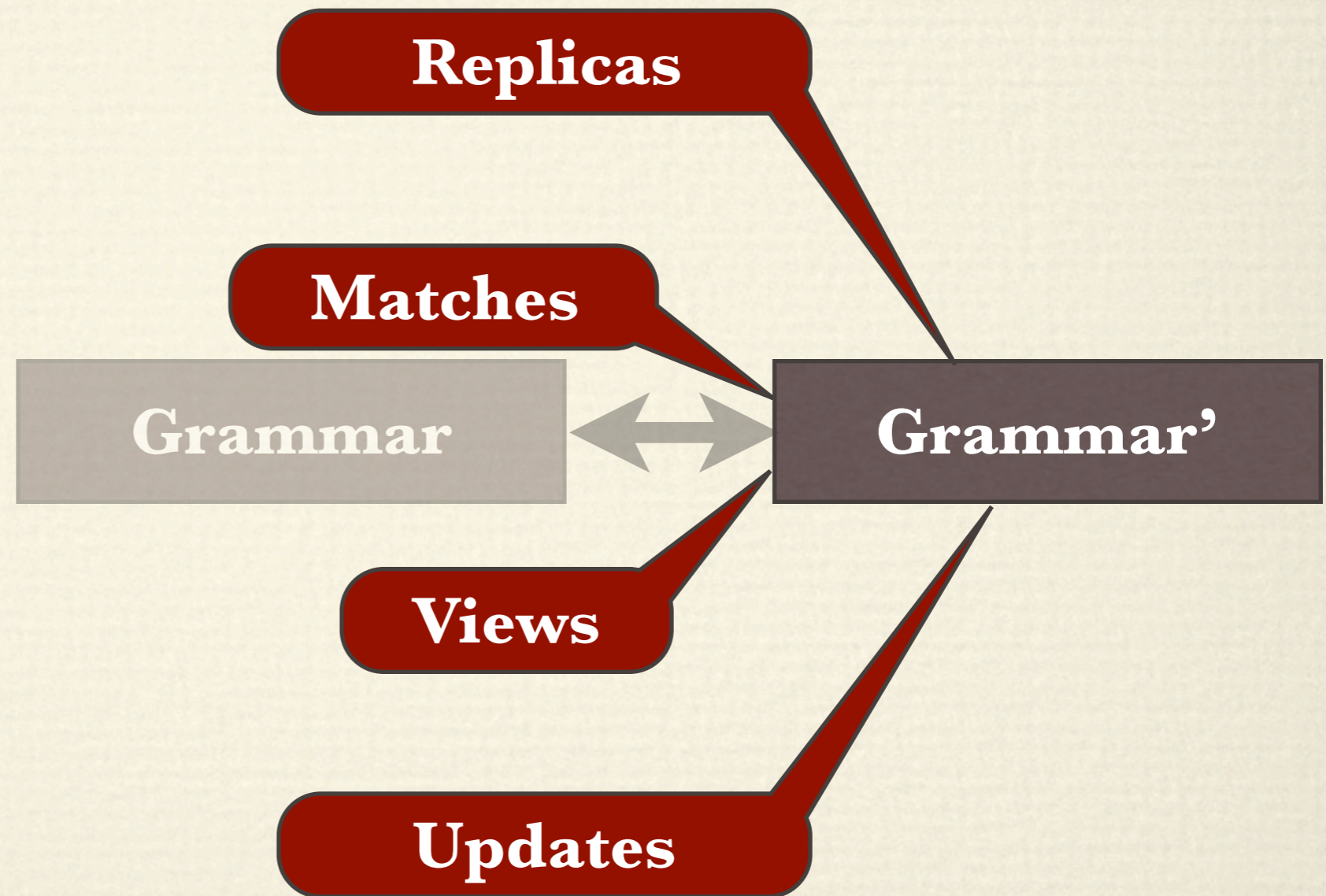


Grammars
from the outside

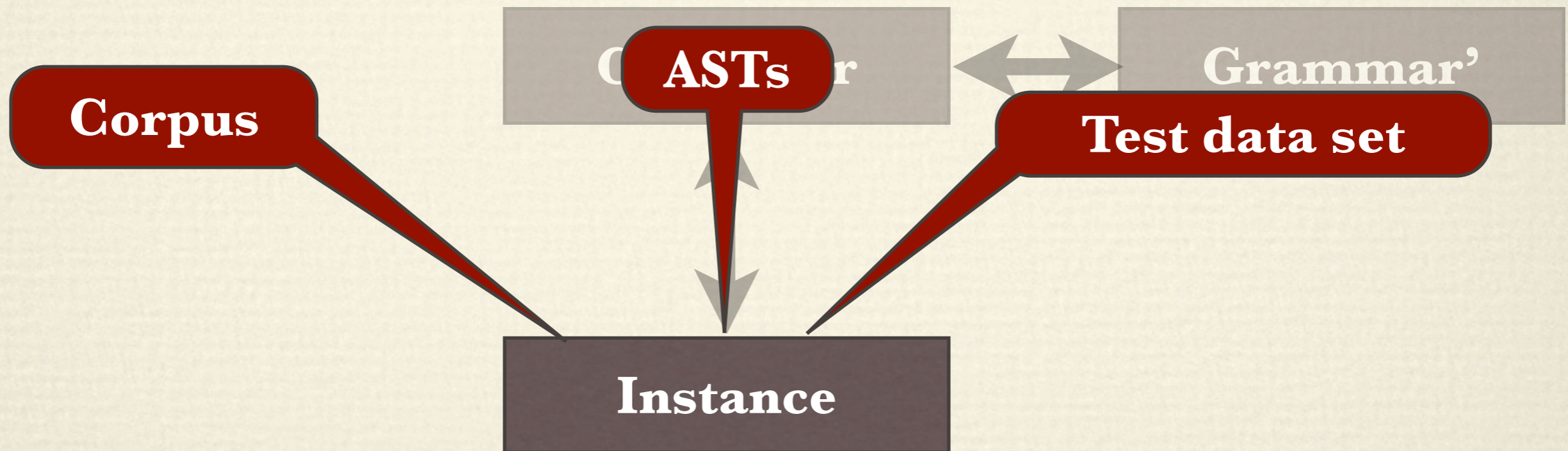
Grammar world

Grammar

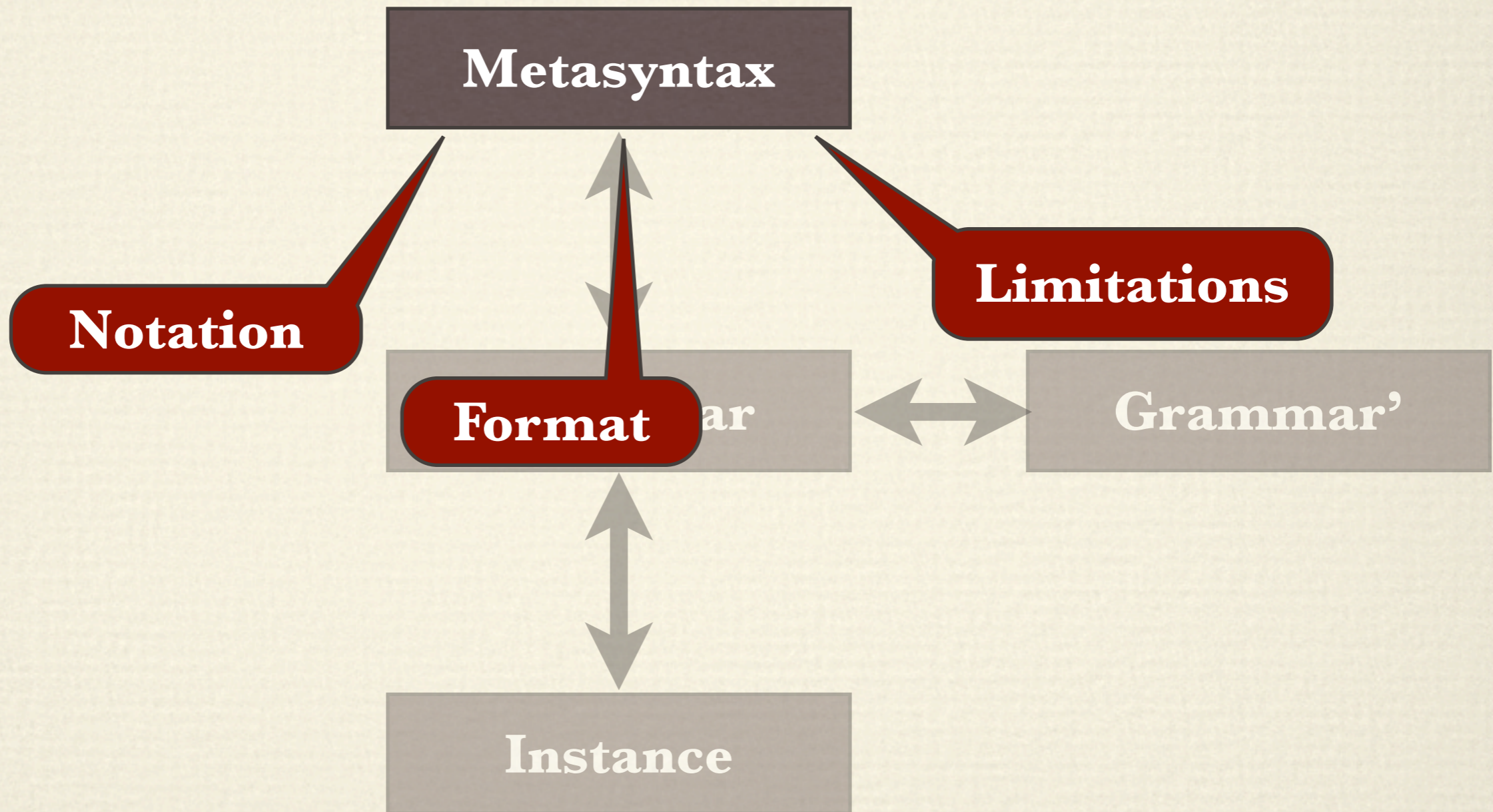
Grammar world: grammars



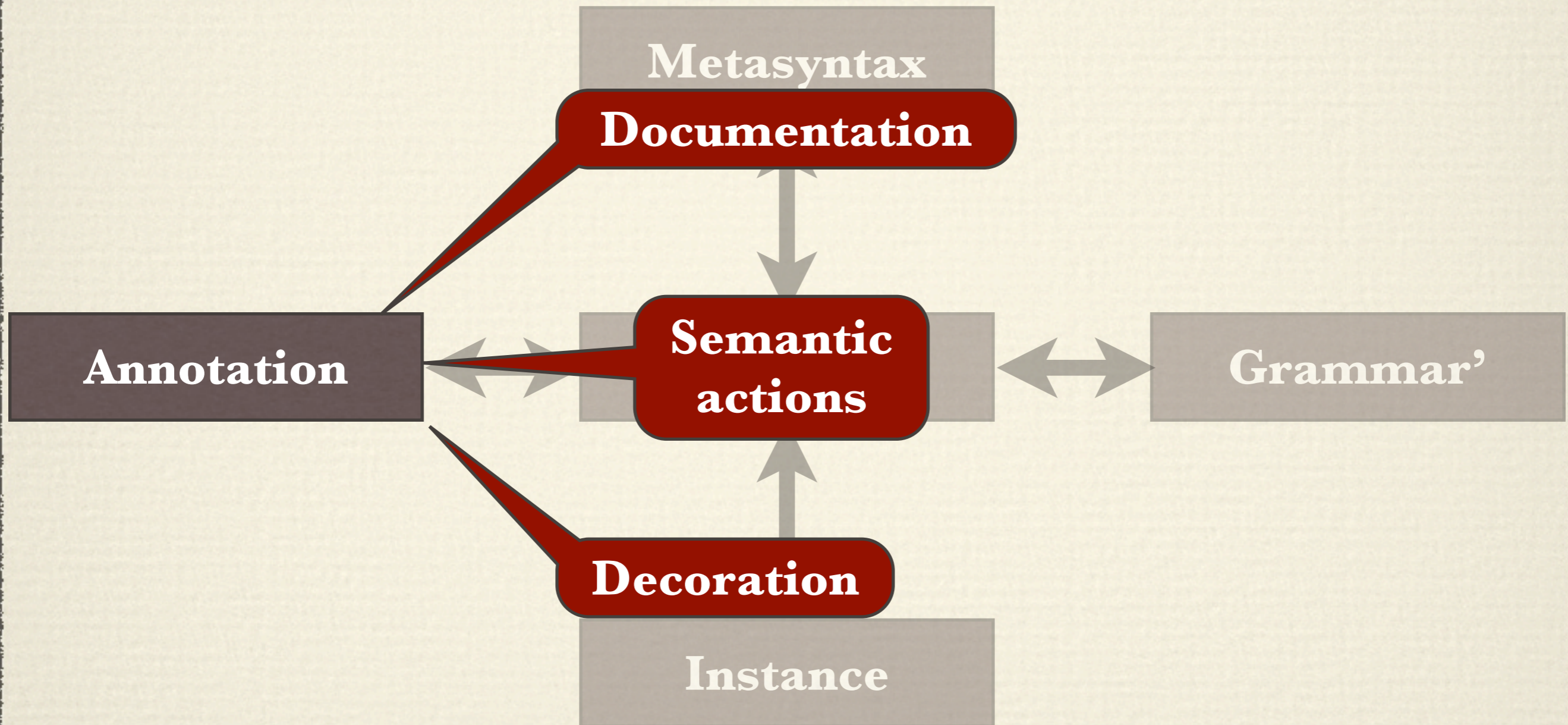
Grammar world: instances



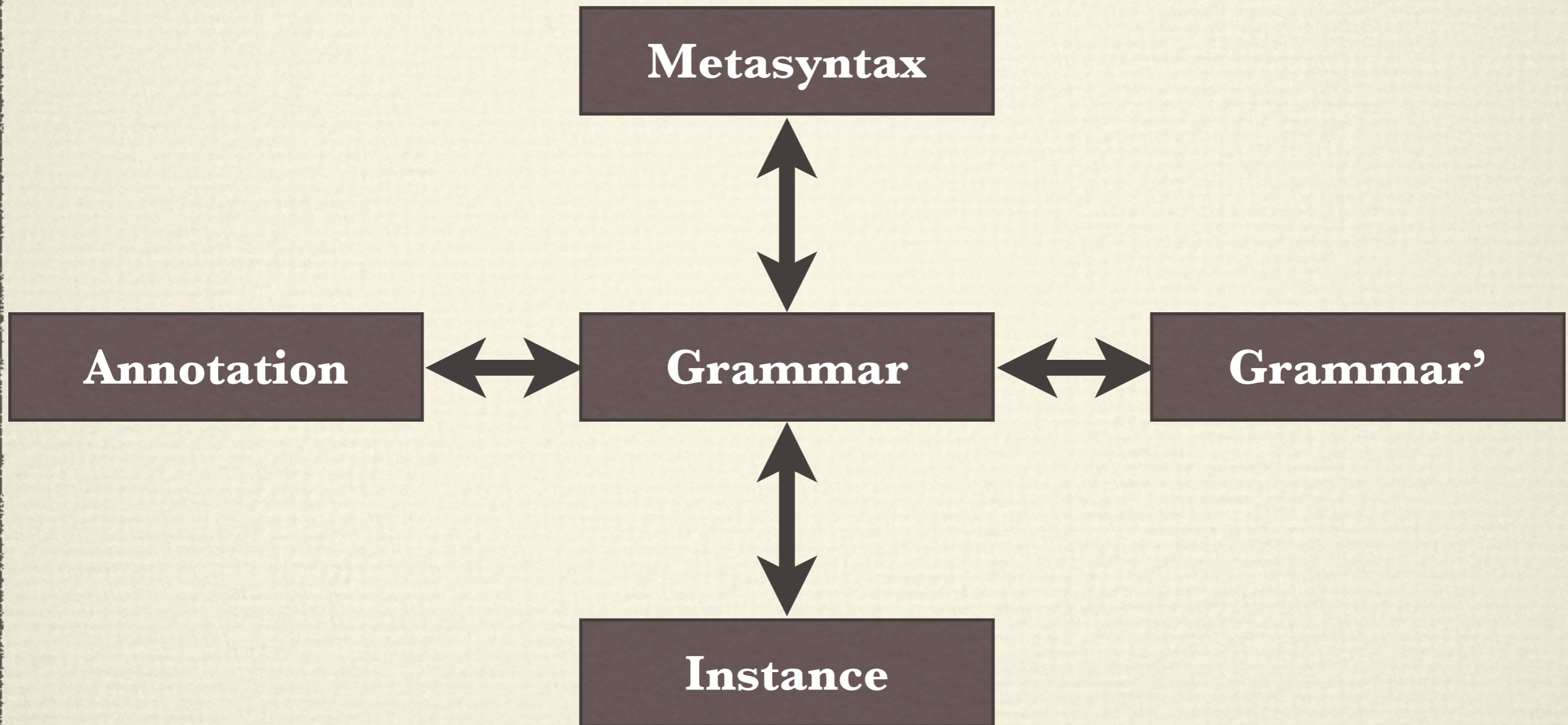
Grammar world: metamodels



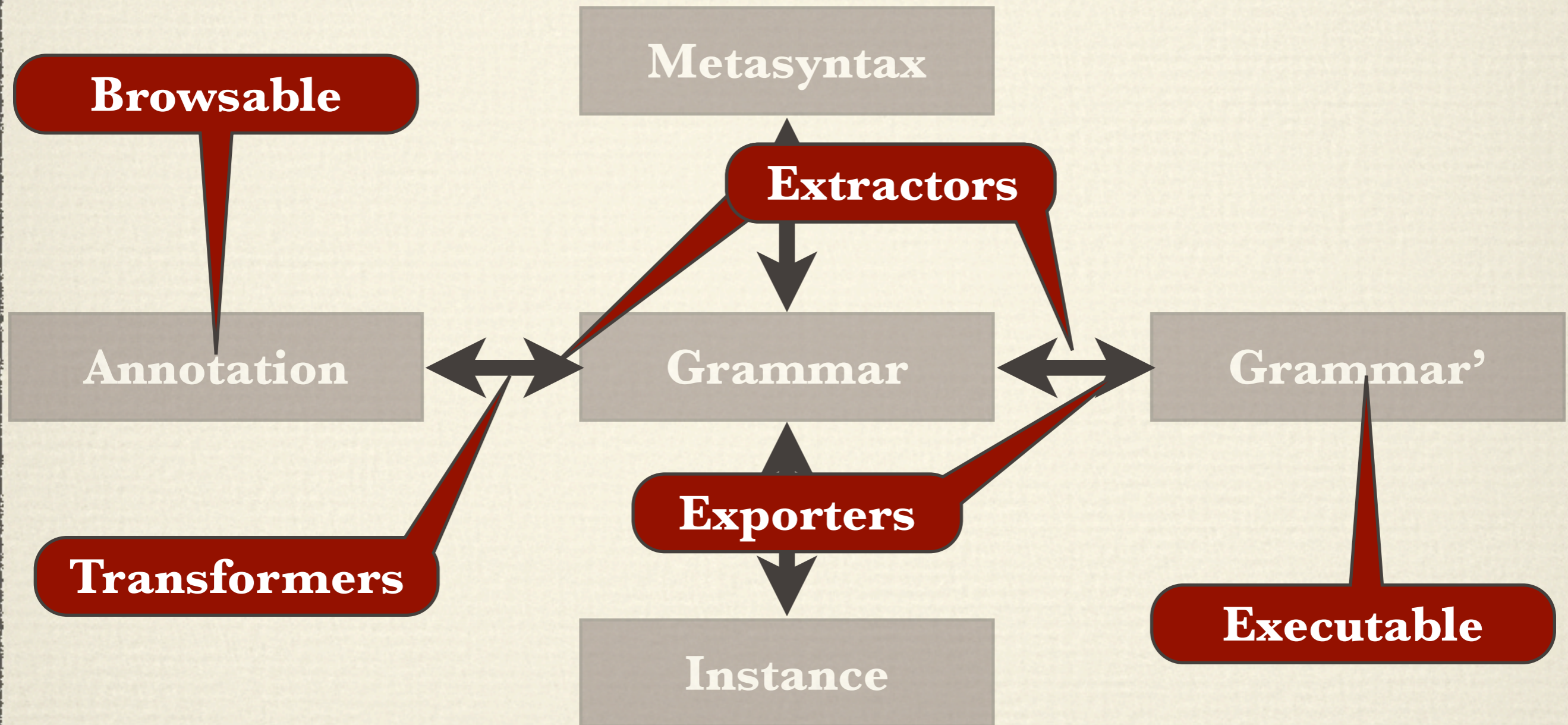
Grammar world: annotations



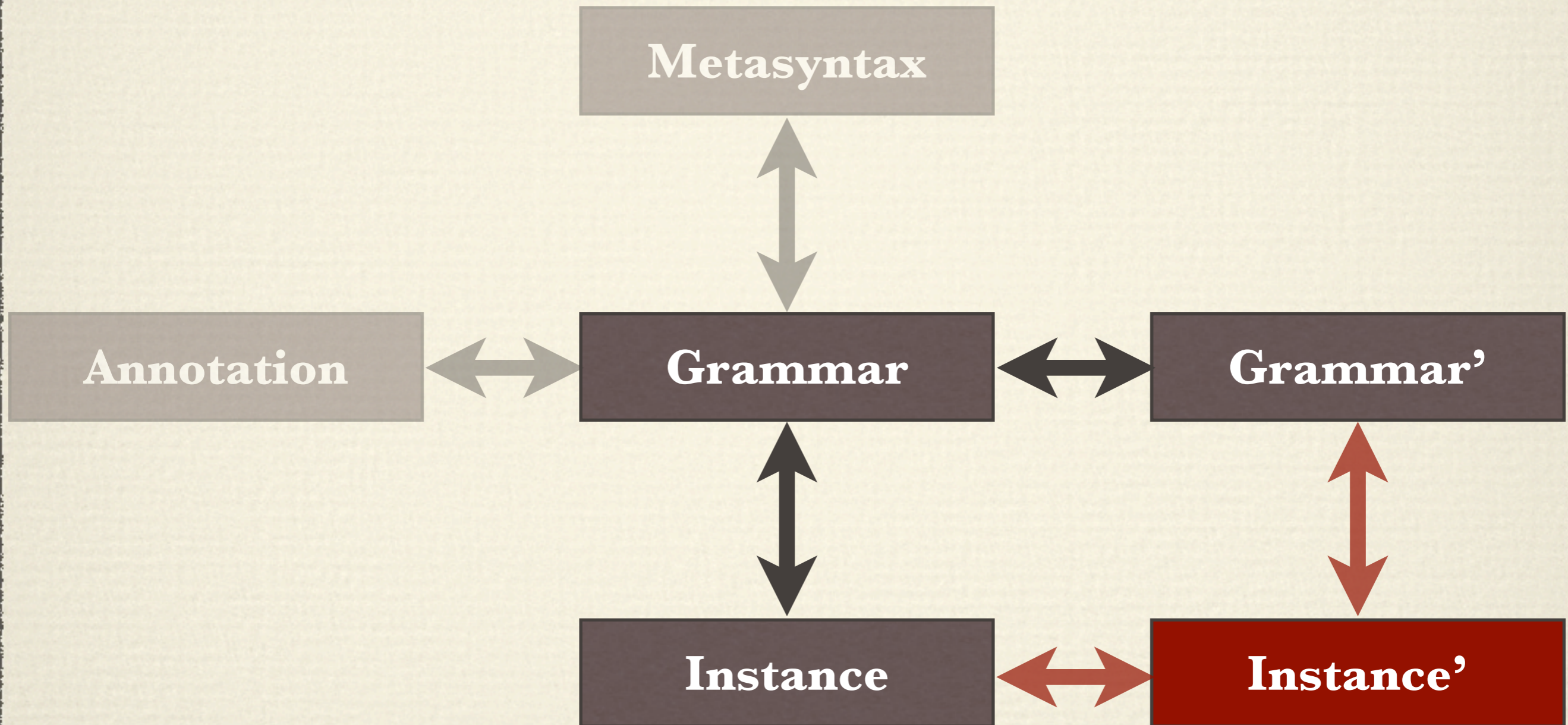
Grammar world



Grammar world: transformations



Grammar world: synchronisations



To summarise



- ◆ Diversity
- ◆ Evolution
- ◆ Extraction
- ◆ Documentation



Grammarware

Grammarware

- Parser
- Compiler
- Interpreter
- Pretty-printer
- Scanner
- Browser
- Static checker
- Structural editor
- IDE
- DSL framework
- Preprocessor
- Postprocessor
- Model checker
- Refactorer
- Code slicer
- API
- XMLware
- Modelware
- Language workbench
- Reverse eng. tool
- Benchmark
- Recommender
- Renovation tool

“Creative Space”



Parser [Aho et al]

- ◆ Syntactic analyser
- ◆ Traverses text
- ◆ Validates conformance to grammar
- ◆ Recognises language instances (“words”)
- ◆ Constructs a parse tree

Parser [Lesk/Schmidt, Johnson]

- ◆ Two separate phases
 - ◆ lexical analysis (lexer)
 - ◆ syntactic analysis (parser)
- ◆ Write your grammar in a very specific style
- ◆ The parser is generated automatically
 - ◆ ...and maybe will work
- ◆ Want a good parser?
 - ◆ forget about readability and reuse

Lesk, Schmidt, *Lex – A Lexical Analyzer Generator*, 1975.

Johnson, *YACC — Yet Another Compiler Compiler*, 1975.

Parser [Cook et al]

- ◆ Recognises syntactic features
- ◆ Creates object structure
- ◆ Results in a graph
 - ◆ object grammars allow crosslinks
- ◆ Naturally bidirectional (parsing/rendering)
 - ◆ object grammars contain pretty-printing info

Parser [Meyers/Vangheluwe]

- ◆ A mapping between
 - ◆ one of concrete syntaxes
 - ◆ the abstract syntax
- ◆ Assumes the “spanning tree” is available or easy

Parser [Ophel]

- ◆ Parser has three tasks
 - ◆ verify that the input is valid
 - ◆ provide semantical framework
 - ◆ if invalid, explain why

To summarise



- ◆ Disregard textual parsing?
- ◆ Separate lexical from syntactic?
- ◆ Link to semantics?



Syntactic Structures (1957)

Despite the undeniable interest and importance of *semantic* and statistical studies of language, they appear to have *no direct relevance* to the problem of determining or characterizing the set of grammatical utterances. I think that we are forced to conclude that grammar is *autonomous* and independent of meaning.

Noam Chomsky



Search



Andreas Wortmann

@andwor



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#models2013

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*Grammars
define structure
&
can assume
different semantics*



<http://instagram.com/p/e0pF3HlDFb/>



RASCAL

Language Workbench

Joint work of:

Ali Afroozeh, Jeroen van den Bos, Magiel Bruntink, Jan van Eijck, Paul Griffioen, Mark Hills, Anastasia Izmaylova, Paul Klint, Wouter Kwakernaak, Dimitrios Kyritsis, Mattijn Lahuis, Davy Landman, Robert van Liere, Bert Lisser, George Marmanidis, Chris Mulder, Atze van der Ploeg, Riemer van Rozen, Alexander Serebrenik, Ashim Shahi, Sunil Simon, Michael Steindorfer, Tijs van der Storm, Ioannis Tzanellis, Jurgen Vinju, Kevin van der Vlist, Vadim Zaytsev (and others)

Language workbench

- ◆ Language oriented programming
 - ◆ Little languages: grep, sed, awk
 - ◆ Homoiconic languages: LISP, REBOL, XSLT
 - ◆ Configuration files
- ◆ Other workbenches
 - ◆ MPS (Völter et al), Intentional Software (Simonyi et al), Epsilon (Paige et al)



Rascal Metaprogramming Language

[⬇ Eclipse update site](#)[🔗 Commandline REPL Jar File](#)

Disclaimer

We currently only release alpha versions of Rascal, which are subject to frequent changes

Eclipse plugin

The Eclipse update site for Rascal is: <http://update.rascal-mpl.org/stable/>

You need *Eclipse for RCP and RAP Developers (Juno or Kepler)* version of Eclipse available at www.eclipse.org/downloads/ to run Rascal. It has been reported recently that the latest release also works with normal (non RCP/RAP) versions of Eclipse. Not thoroughly tested though!

Please note that:

- Rascal now needs a JDK because it uses the Java compiler, so please download a JDK, not just a JRE.
- You may have to [edit the Eclipse init file](#) so that Eclipse can find your Java installation and Eclipse can allocate enough resources.
- For generating parsers, Rascal uses quite a bit of memory. Please use `-vmargs -Xmx1G -Xss32m`

Follow these steps to install the plugin into Eclipse

1. Start Eclipse.
2. Select Help -> Install New Software.
3. Make sure that the tick for "Contact all update sites during install to find required software" is enabled.
4. Type: <http://update.rascal-mpl.org/stable/> in the "Work with" edit box.
5. Select the feature *Rascal*.
6. Select Next (several times) and accept the software license. The process may take a few minutes!
7. Once these features have been installed, restart Eclipse.



Technical challenges

- ◆ How to parse source code/data files/models?
- ◆ How to extract facts from them?
- ◆ How to perform computations on these facts?
- ◆ How to generate new source code (transform, refactor, compile)?
- ◆ How to synthesize other information?

EASY: Extract-Analyze-SYnthesize Paradigm

Metaprogramming is EASY

◆ Extract

- ◆ Fast context-free general top-down parsing
- ◆ Pattern matching & generic traversal

◆ Analyze

- ◆ Relational queries and comprehensions
- ◆ Backtracking, fixed point computation, ...

◆ SYnthesize

- ◆ String templates
- ◆ Concrete syntax
- ◆ Interactive visualization generator



Metaprogramming is EASY

◆ Extract

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◆ SYnthesize

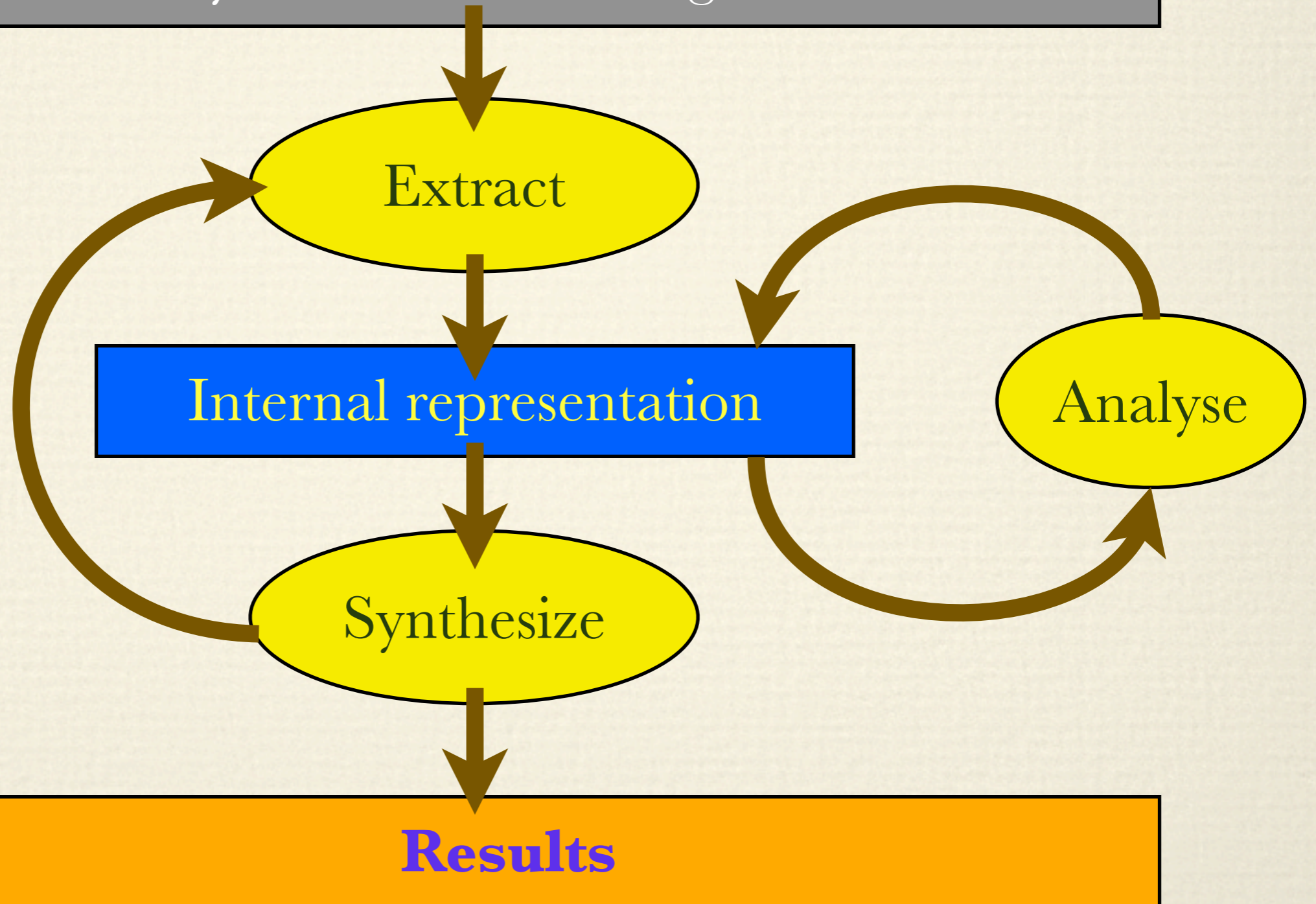
- ◆ String templates
- ◆ Concrete syntax
- ◆ Interactive visualization generator





The EASY paradigm

System under investigation



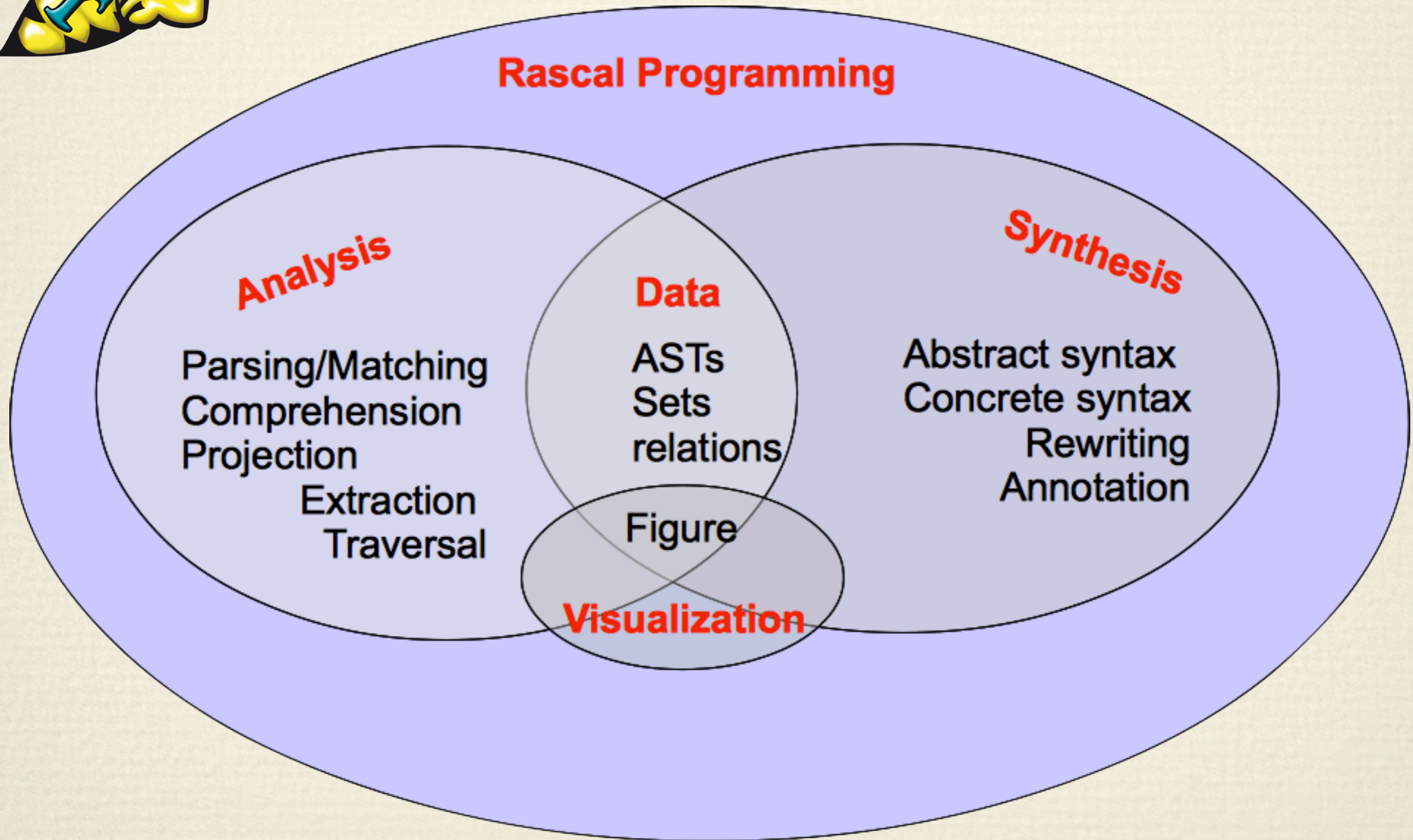


Why a new language?

- ◆ No current technology spans the full range of **EASY** steps
- ◆ There are many fine technologies but they are
 - ◆ highly specialized with steep learning curves
 - ◆ hard to learn unintegrated technologies
 - ◆ not integrated with a standard IDE
 - ◆ hard to extend
- ◆ Goal: keep all benefits of ASF+SDF and Rscript
 - ◆ in a new, unified, extensible, teachable framework



Bridging the gaps





Rascal keywords

- ◆ Complex built-in data types
- ◆ Immutable data
- ◆ Static safety
- ◆ Generic types
- ◆ Local type inference
- ◆ Pattern matching
- ◆ Syntax definitions & parsing
- ◆ Concrete syntax
- ◆ Visiting/traversal
- ◆ Comprehensions
- ◆ Higher-order
- ◆ Familiar syntax
- ◆ Java and Eclipse integration
- ◆ Read-Eval-Print (REPL)



Rascal design

- ◆ Java-like syntax presumably familiar
- ◆ Embedded in Eclipse
- ◆ Layered design
- ◆ Syntax analysis
- ◆ Term rewriting
- ◆ Relational calculus



Rascal design

- ◆ Java-like syntax
- ◆ Embedded in Eclipse
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- ◆ Term rewriting
- ◆ Relational calculus

installs as a plugin



Rascal design

- ◆ Java-like syntax
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- ◆ Term rewriting
- ◆ Relational calculus

low barrier to entry,
learn features as you go



Rascal design

- ◆ Java-like syntax
- ◆ Embedded in Eclipse
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- ◆ Relational calculus

concrete syntax matching



Rascal design

- ◆ Java-like syntax
- ◆ Embedded in Eclipse
- ◆ Layered design
- ◆ Syntax analysis
- ◆ Term rewriting
- ◆ Relational calculus

traversals, matching, ...



Rascal design

- ◆ Java-like syntax
- ◆ Embedded in Eclipse
- ◆ Layered design
- ◆ Syntax analysis
- ◆ Term rewriting
- ◆ Relational calculus

relations for sharing/merging
of facts for different languages

To summarise



- ◆ Rascal is a language workbench
- ◆ EASY paradigm
 - ◆ extract
 - ◆ analyse
 - ◆ synthesise
- ◆ Metaprogramming
- ◆ <http://rascal-mpl.org>





Rascal features



Rich (immutable) data

- ◆ Built-in types:

- ◆ lists

- ◆ sets

- ◆ maps

- ◆ tuples

- ◆ relations

- ◆ with comprehensions and many operators

```
rascal> [1..10]
```

```
list[int]: [1,2,3,4,5,6,7,8,9,10]
```

```
rascal> [x/2 | x <- [1..10]]
```

```
list[int]: [0,1,1,2,2,3,3,4,4,5]
```

```
rascal> {x/2 | x <- [1..10]} + {4,5,6}
```

```
set[int]: {6,5,4,3,2,1,0}
```



Syntax definitions

- ◆ Define lexical syntax
- ◆ Define context-free syntax
- ◆ Define whitespace/layout/...
- ◆ Get GLL parser for free
- ◆ Define an algebraic data type
- ◆ Automatically implode parse trees to ASTs



Syntax definitions

```
lexical Id = [A-Za-züäöß]+ !>> [A-Za-züäöß];  
lexical Num = [0-9]+ !>> [0-9];
```

- ◆ Define lexical syntax
- ◆ Define context-free syntax
- ◆ Define whitespace/layout/...
- ◆ Get GLL parser for free
- ◆ Define an algebraic data type
- ◆ Automatically implode parse trees to ASTs



Syntax definitions

```
start syntax System = Line+;  
syntax Line = Num ":" {Id ",", "}" + "." ;
```

- ◆ Define lexical syntax
- ◆ Define context-free syntax
- ◆ Define whitespace/layout/...
- ◆ Get GLL parser for free
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Syntax definitions

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```
layout WS = [\ \t\n\r]* !>> [\ \t\n\r];
```



Syntax definitions

- ◆ Define lexical syntax
- ◆ Define context-free syntax
- ◆ Define whitespace/layout/...
- ◆ Get GLL parser for free
- ◆ Define an algebraic data type
- ◆ Automatically implode parse trees to ASTs



Patterns

- ◆ Pattern matching
 - ◆ on concrete syntax
 - ◆ on lists
 - ◆ on sets
- ◆ on trees
- ◆ ...
- ◆ Pattern-driven dispatch

```
rascal> {int x, str y} := {2}
```

```
bool: false
```

```
rascal> {int x, str y} := {2,"3"}
```

```
bool: true
```

```
rascal> {int x, *y, str z} := {2,2,2,"3",4,"2"}
```

```
bool: true
```



Other pattern kinds

- ◆ **Regular:** grep/Perl like regular expressions
`/^<before: \W*><word: \w+><after: .*$/`
- ◆ **Abstract:** match data types
`whileStat(Exp, Stats*)`
- ◆ **Concrete:** match parse trees
`` while <Exp> do <Stats*> od ``



Pattern-directed invocation

Prolog?

```
bool eqfp(fpnt(), fpnt()) = true;
```

```
bool eqfp(fpopt(), fpopt()) = true;
```

```
bool eqfp(fpplus(), fpplus()) = true;
```

```
bool eqfp(fpstar(), fpstar()) = true;
```

```
bool eqfp(fpempty(), fpempty()) = true;
```

```
bool eqfp(fpmany(L1), fpmany(L2)) = multiseteq(L1, L2);
```

```
default bool eqfp(Footprint pi, Footprint xi) = false;
```



Switch/case

```
switch(p)
```

```
{
```

```
  case (DCGFun) `[]` => ["ε"];
```

```
  case (DCGFun) `<Word n>` =>
```

```
    ["<n>" | "<n>"==toLowerCase("<n>")];
```

```
  case (DCGFun) `(<{DCGFun " ,"}* args>` =>
```

```
    [*getTags(a) | a <- args];
```

```
  case (DCGFun) `<Word f> (<{DCGFun " ,"}* as>` =>
```

```
    ["<f>"] + [*getTags(a) | a <- as];
```

```
  default ...
```

```
}
```




Visitor and foldr

```
@contributor{Bas Basten - Bas.Basten@cwi.nl (CWI)}
```

```
@contributor{Mark Hills - Mark.Hills@cwi.nl (CWI)}
```

```
module Operations
```

```
import AST;
```

```
import IO;
```

```
public Company cut(Company c) {
```

```
  return visit (c) {
```

```
    case employee(name, [*ep,ip:intProp("salary",salary),*ep2])
```

```
      => employee(name, [*ep,ip[intVal=salary/2],*ep2])
```

```
  }}
```

```
public int total(Company c) {
```

```
return (0 | it+salary | /employee(name, [*ep,ip:intProp  
("salary",salary),*ep2]) <- c);
```

```
}
```

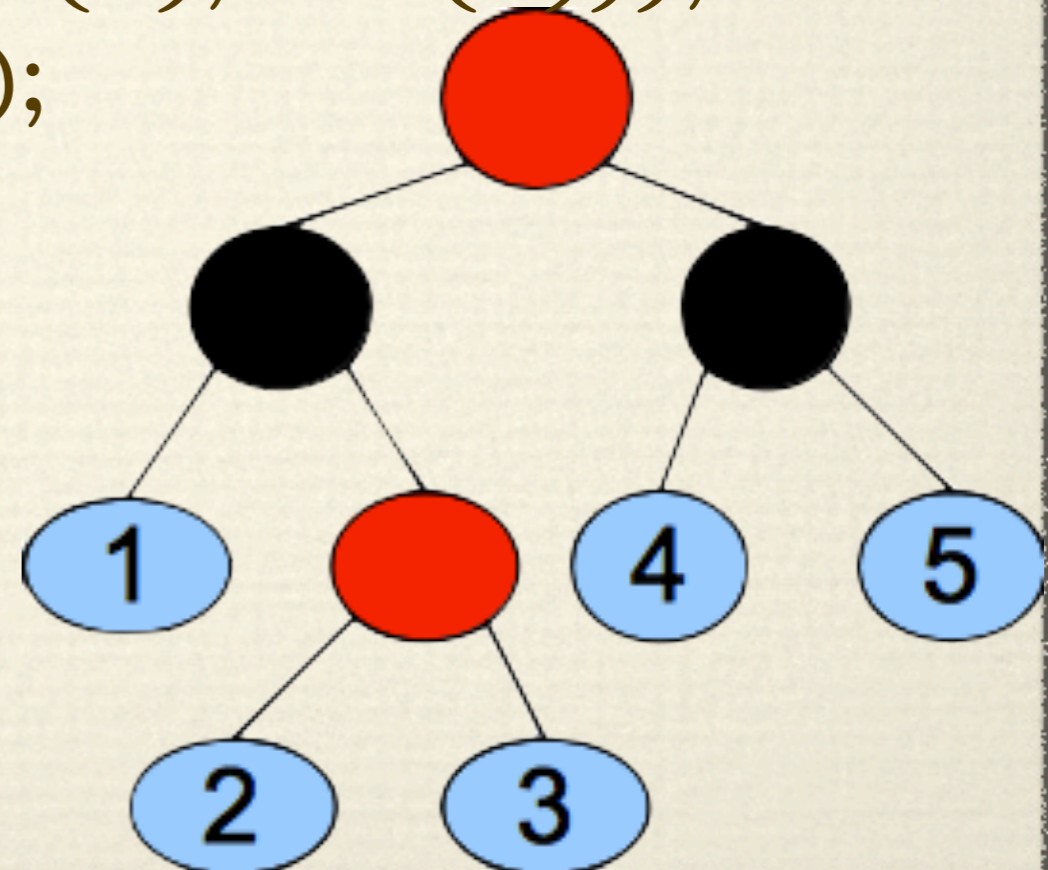


Polychromatic trees

```
data CTree = leaf(int N)  
    | red(CTree left, CTree right)  
    | black(CTree left, CTree right) ;
```

```
rb = red(black(leaf(1), red(leaf(2), leaf(3))),  
        black(leaf(4), leaf(5)));
```

```
// TODO: count red nodes
```

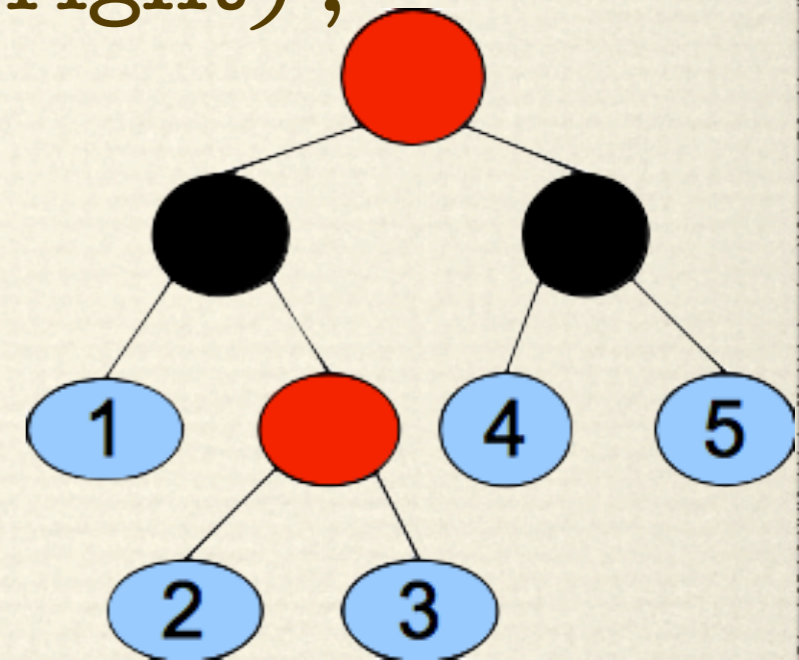




Polychromatic trees

```
data CTree = leaf(int N)
          | red(CTree left, CTree right)
          | black(CTree left, CTree right);
```

```
public int cntRed(CTree t) {
switch(t) {
  case leaf(_): return 0;
  case red(L,R):
    return 1 + cntRed(L) + cntRed(R);
  case black(L,R):
    return cntRed(L) + cntRed(R);
}}
```

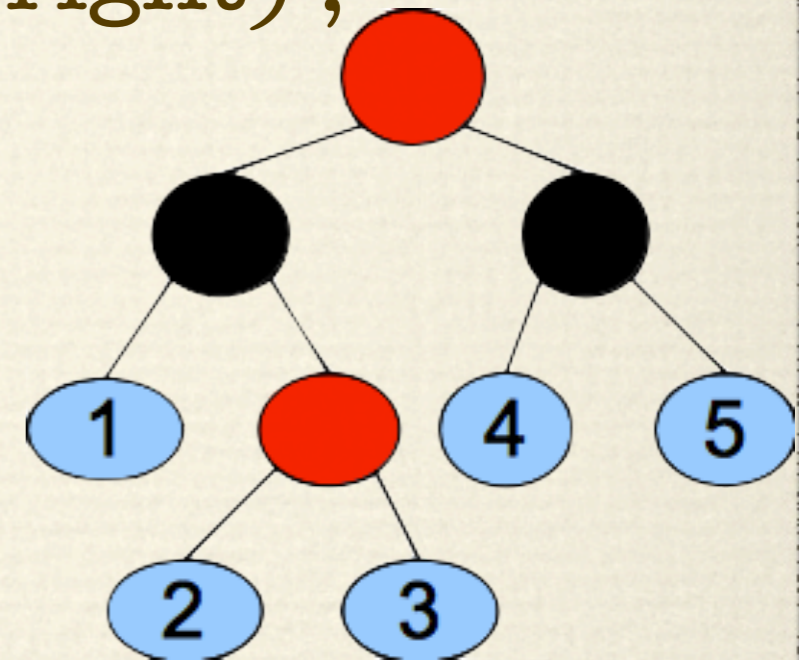




Polychromatic trees

```
data CTree = leaf(int N)
           | red(CTree left, CTree right)
           | black(CTree left, CTree right);
```

```
public int cntRed(CTree t) {
  int c = 0;
  visit(t) { case red(_, _): c += 1; };
  return c;
}
```

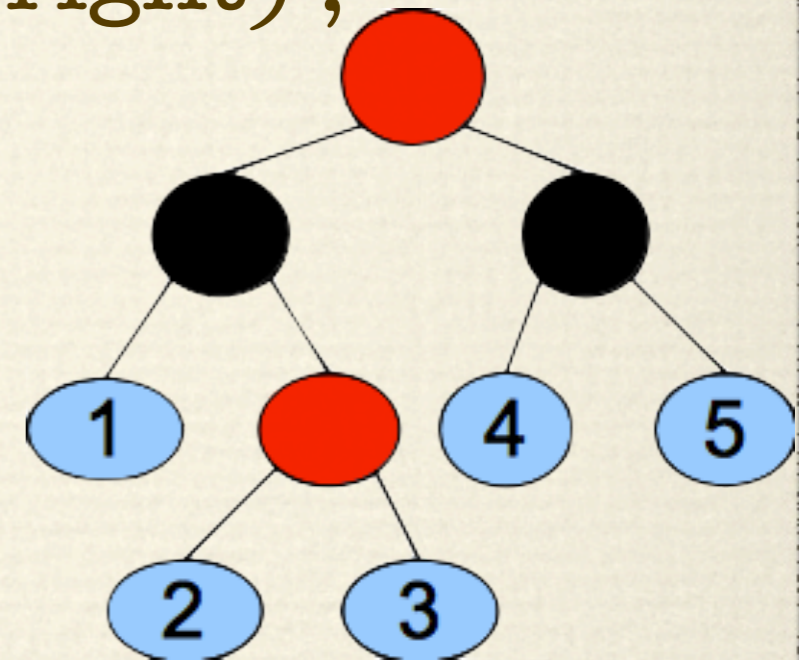




Polychromatic trees

```
data CTree = leaf(int N)  
          | red(CTree left, CTree right)  
          | black(CTree left, CTree right) ;
```

```
public int cntRed(CTree t) =  
size([n | /n:red(_,_) := t]);
```



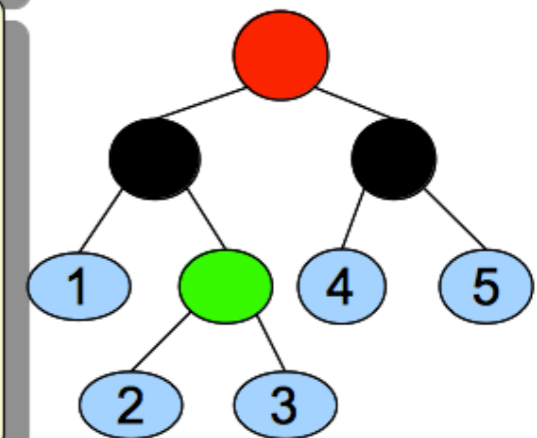
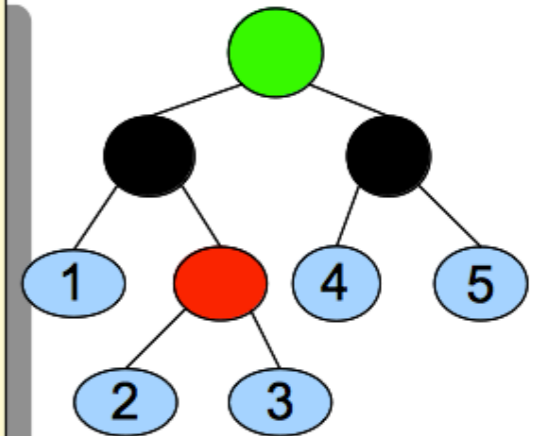
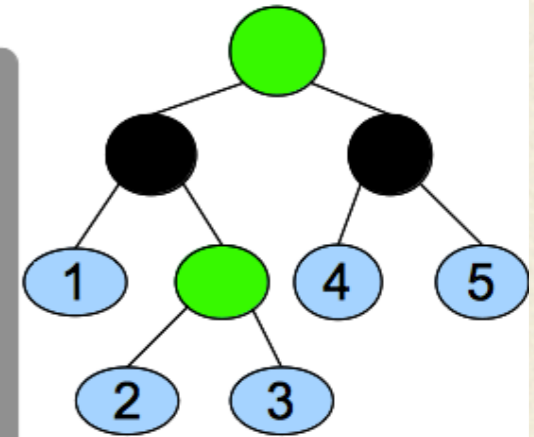


Full/shallow/deep

```
public CTree frepl(CTree T) {  
    return visit (T) {  
        case red(CTree T1, Ctree T2) => green(T1, T2)  
    };  
}
```

```
public Ctree srepl(CTree T) {  
    return top-down-break visit (T) {  
        case red(Ctree T1, CTree T2) => green(T1, T2)  
    };  
}
```

```
public Ctree drepl(Ctree T) {  
    return bottom-up-break visit (T) {  
        case red(CTree T1, CTree T2) => green(T1, T2)  
    };  
}
```




To summarise



- ◆ Rascal is a language workbench
- ◆ Sets, relations, maps...
- ◆ Grammars, ADTs...
- ◆ Cool comprehensions
- ◆ Cool pattern matching
- ◆ Many more features
- ◆ <http://rascal-mpl.org>



Work in progress

A photograph showing the interior of the Eiffel Tower during renovation. The image is dominated by a complex, dark metal lattice structure that forms the tower's framework. In the center, a vertical shaft is lined with scaffolding and a red lift basket is visible. The lighting is dramatic, with strong highlights and deep shadows, emphasizing the geometric patterns of the metalwork.

POWERED BY

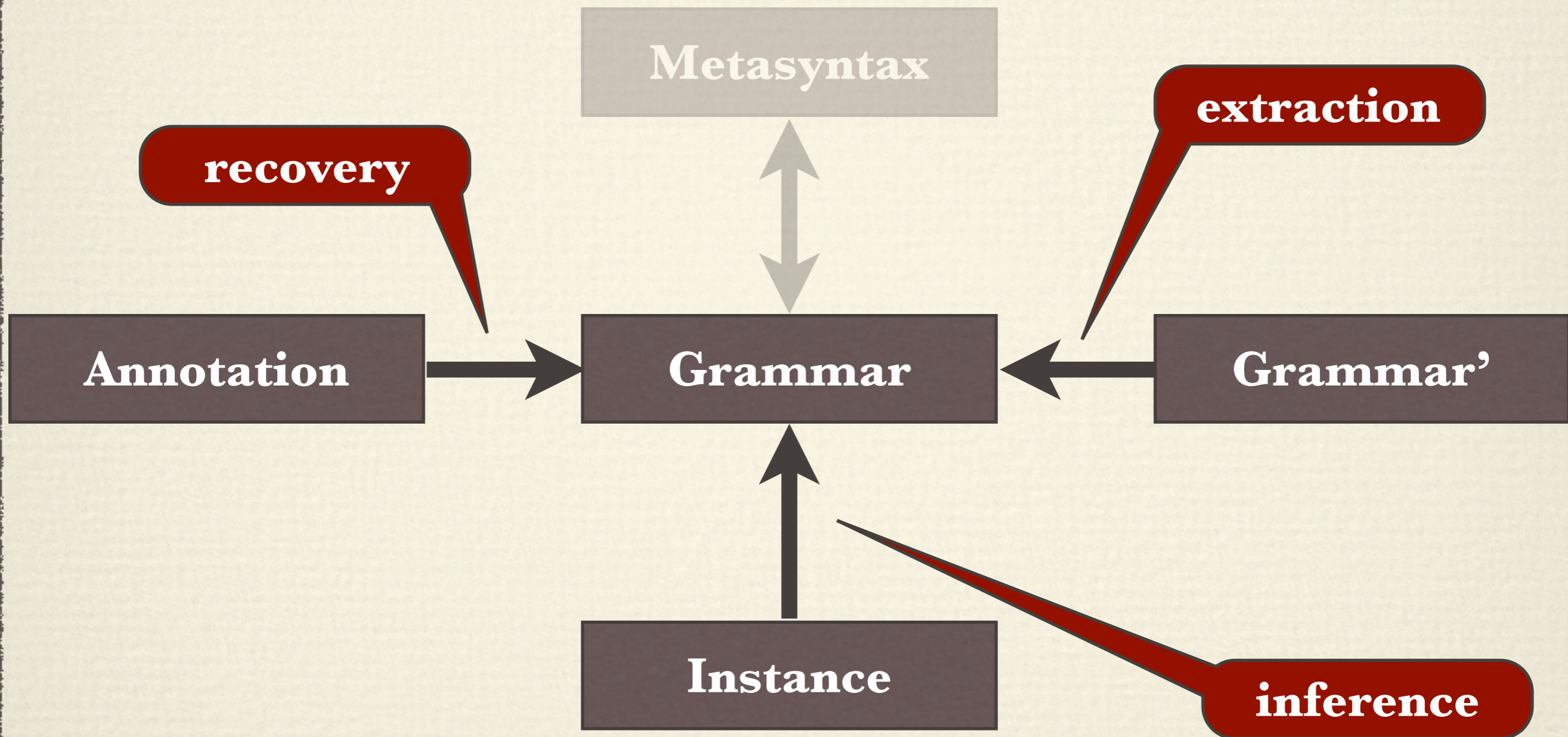
GRAMMARLAB

Grammar extraction

Grammar extraction

- ◆ Given is an artefact containing grammar knowledge:
 - ◆ a grammar
 - ◆ a parser specification
 - ◆ a metamodel
 - ◆ grammarware source code
 - ◆ a data schema
 - ◆ documentation
- ◆ How to extract a grammar from it?

Grammar extraction



Variety of notations

program ::=

function+;

function ::=

name argument* "=" expr?;

- ◆ Such details can be modeled
- ◆ We compose a notation specification

Grammar notations: BNF

Arithmetic Expressions, Boolean Expressions, and Expressions

$\langle \text{factor} \rangle ::= \langle \text{number} \rangle \mid \langle \text{function} \rangle \mid \langle \text{variable} \rangle \mid \langle \text{subscr var} \rangle$
 $\mid (\langle \text{ar exp} \rangle) \mid \langle \text{factor} \rangle \uparrow \langle \text{ar exp} \rangle \downarrow$

$\langle \text{term} \rangle ::= \langle \text{factor} \rangle \mid \langle \text{term} \rangle \times \langle \text{factor} \rangle \mid \langle \text{term} \rangle / \langle \text{factor} \rangle$

$\langle \text{ar exp} \rangle ::= \langle \text{term} \rangle \mid \langle \text{term} \rangle + \langle \text{term} \rangle \mid \langle \text{term} \rangle - \langle \text{term} \rangle$
 $\mid \langle \text{ar exp} \rangle + \langle \text{term} \rangle$
 $\mid \langle \text{ar exp} \rangle - \langle \text{term} \rangle$

$\langle \text{ar exp A} \rangle ::= \langle \text{ar exp} \rangle$

$\langle \text{relation} \rangle ::= \langle \text{ar exp} \rangle < \mid \geq \mid = \mid \neq$

$\langle \text{rel exp} \rangle ::= (\langle \text{ar exp} \rangle \langle \text{relation} \rangle \langle \text{ar exp A} \rangle)$

$\langle \text{bool term} \rangle ::= 0 \mid 1 \mid \langle \text{rel exp} \rangle \mid \langle \text{function} \rangle$
 $\mid \langle \text{variable} \rangle \mid \langle \text{subscr var} \rangle \mid \langle \text{bool exp} \rangle$
 $\mid \neg \langle \text{bool term} \rangle$

$\langle \text{bool exp} \rangle ::= \langle \text{bool term} \rangle \mid \langle \text{bool exp} \rangle \vee \langle \text{bool term} \rangle$
 $\mid \langle \text{bool exp} \rangle \wedge \langle \text{bool term} \rangle$
 $\langle \text{bool exp} \rangle ::= \langle \text{bool term} \rangle$

$\langle \text{exp} \rangle ::= \langle \text{ar exp} \rangle \mid \langle \text{bool exp} \rangle$

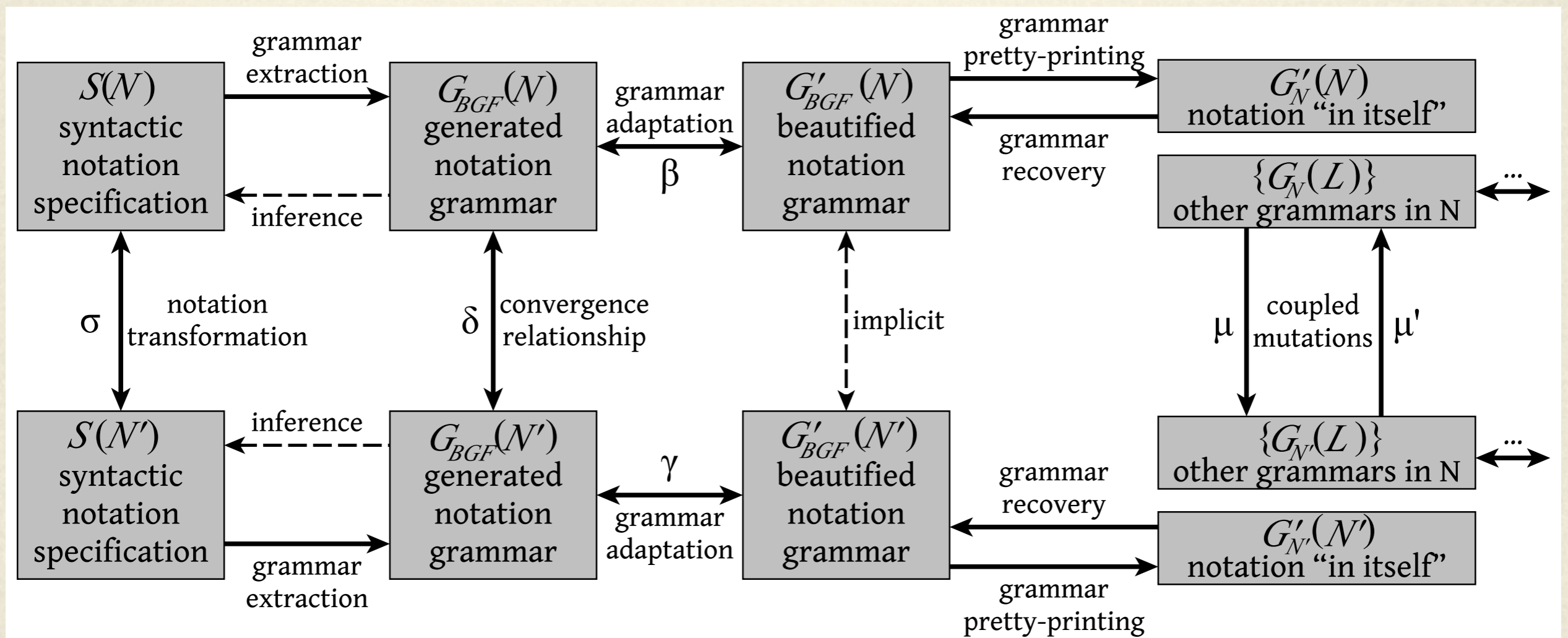
Rascal - gtrafo/src/bgf2rsc/III/Final.III - Eclipse Platform

Transformer.rsc Rascal Tutor GeneratedLLL.rsc Hunter.rsc III.edd Final.III

```
1 # Idents:      153
2 # - top:      1
3 # - used:     152
4 # - defined:  145
5 # - undefined: 8
6 # Literals:   121
7
8 ref-or-out
9   : "ref"
10  | "out"
11  ;
12
13 expression-unary-operator
14   : lex-csharp-extra/plus
15   | lex-csharp-extra/minus
16   | increment-decrement
17   | "|"
18   | "~"
19   | "*"
20   ;
21
22 increment-decrement
23   : "++"
24   | "--"
25   ;
26
27 expression-shift-operator
28   : "<<"
29   | ">>"
30   ;
31
32 expression-relational-operator
33   : lex-csharp-extra/less-than
34   | lex-csharp-extra/greater-than
35   | "<="
36   | ">="
37   | "is"
38   | "as"
39   ;
40
41 expression-equality-operator
42   : "=="
43   | "!="
44   ;
45
46 conversion-kind
```

Writable Smart Insert 1

Metasyntactic evolution





EBNF in GrammarLab

```
alias EBNF = map[Metasymbol, str];
data Metasymbol
  = start_grammar_symbol()
  | end_grammar_symbol()
  | epsilon_metasymbol()
  | defining_symbol()
  | multiple_defining_symbol()
  | terminator_symbol()
  | definition_separator_symbol()
  | start_option_symbol()
  | end_option_symbol()
  | start_nonterminal_symbol()
  | end_nonterminal_symbol()
```

...



EBNF in GrammarLab

```
public EBNF DefaultEBNF = (  
    epsilon_metasybol(): "ε",  
    defining_symbol(): " ::= ",  
    terminator_symbol(): " ;\n",  
    definition_separator_symbol(): " | ",  
    start_group_symbol(): "(",  
    end_group_symbol(): ")",  
    start_terminal_symbol(): "\"",  
    end_terminal_symbol(): "\"",  
    postfix_option_symbol(): "?",  
    ...  
);
```



IDE support

```
EBNF ebnf = ...;  
edd2rsc(ebnf, |cwd:///EBNFGrammar.rsc|);
```

...

```
Tree getEBNF(str s, loc z)  
    = parse(#EBNFGrammar, z);  
public void registerEBNF()  
    = registerLanguage("EBNF", "bnf", getEBNF);
```


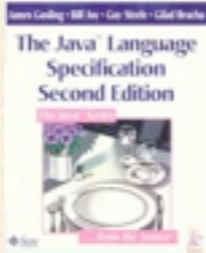
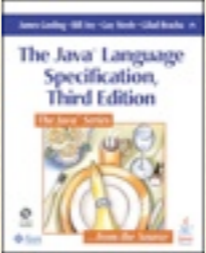

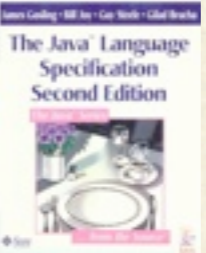
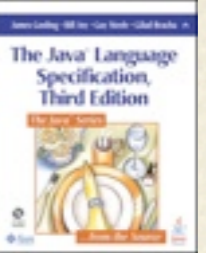
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18   | "~"
19   | "*"
20   ;
21
22 increment-decrement
23   : "++"
24   | "--"
25   ;
26
27 expression-shift-operator
28   : "<<"
29   | ">>"
30   ;
31
32 expression-relational-operator
33   : lex-csharp-extra/less-than
34   | lex-csharp-extra/greater-than
35   | "<="
36   | ">="
37   | "is"
38   | "as"
39   ;
40
41 expression-equality-operator
42   : "=="
43   | "!="
44   ;
45
46 conversion-kind
```

Writable Smart Insert 1

Extraction of Java grammars

	 impl1	 impl2	 impl3	 read1	 read2	 read3	Total
Arbitrary lexical decisions	2	109	60	1	90	161	423
Well-formedness violations	5	0	7	4	11	4	31
Indentation violations	1	2	7	1	4	8	23
Recovery rules	3	12	18	2	59	47	141
○ Match parentheses	0	3	6	0	0	0	9
○ Metasymbol to terminal	0	1	7	0	27	7	42
○ Merge adjacent symbols	1	0	0	1	1	0	3
○ Split compound symbol	0	1	1	0	3	8	13
○ Nonterminal to terminal	0	7	3	0	8	11	29
○ Terminal to nonterminal	1	0	1	1	17	13	33
○ Recover optionality	1	0	0	0	3	8	12
Purge duplicate definitions	0	0	0	16	17	18	51
Total	11	123	92	24	181	238	669



Extraction & recovery

```
EBNF ebnf = ...;
```

```
GGrammar bgf
```

```
= ebnf2bgf(ebnf, |cwd:///grammar.ebnf|);
```

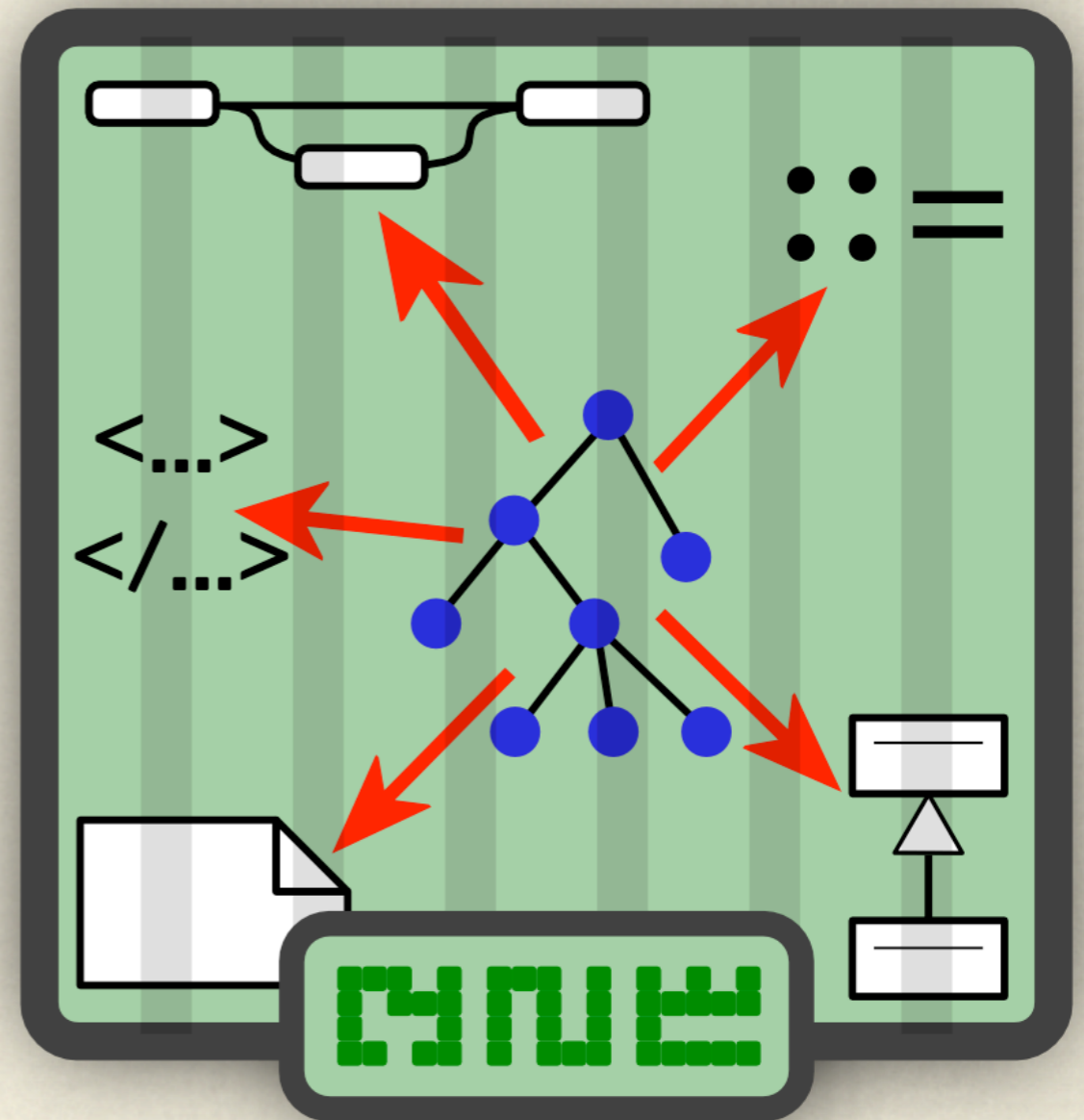
```
GGrammar bgf
```

```
= hunter(ebnf, |cwd:///grammar.ebnf|);
```

```
// possibly with errors
```

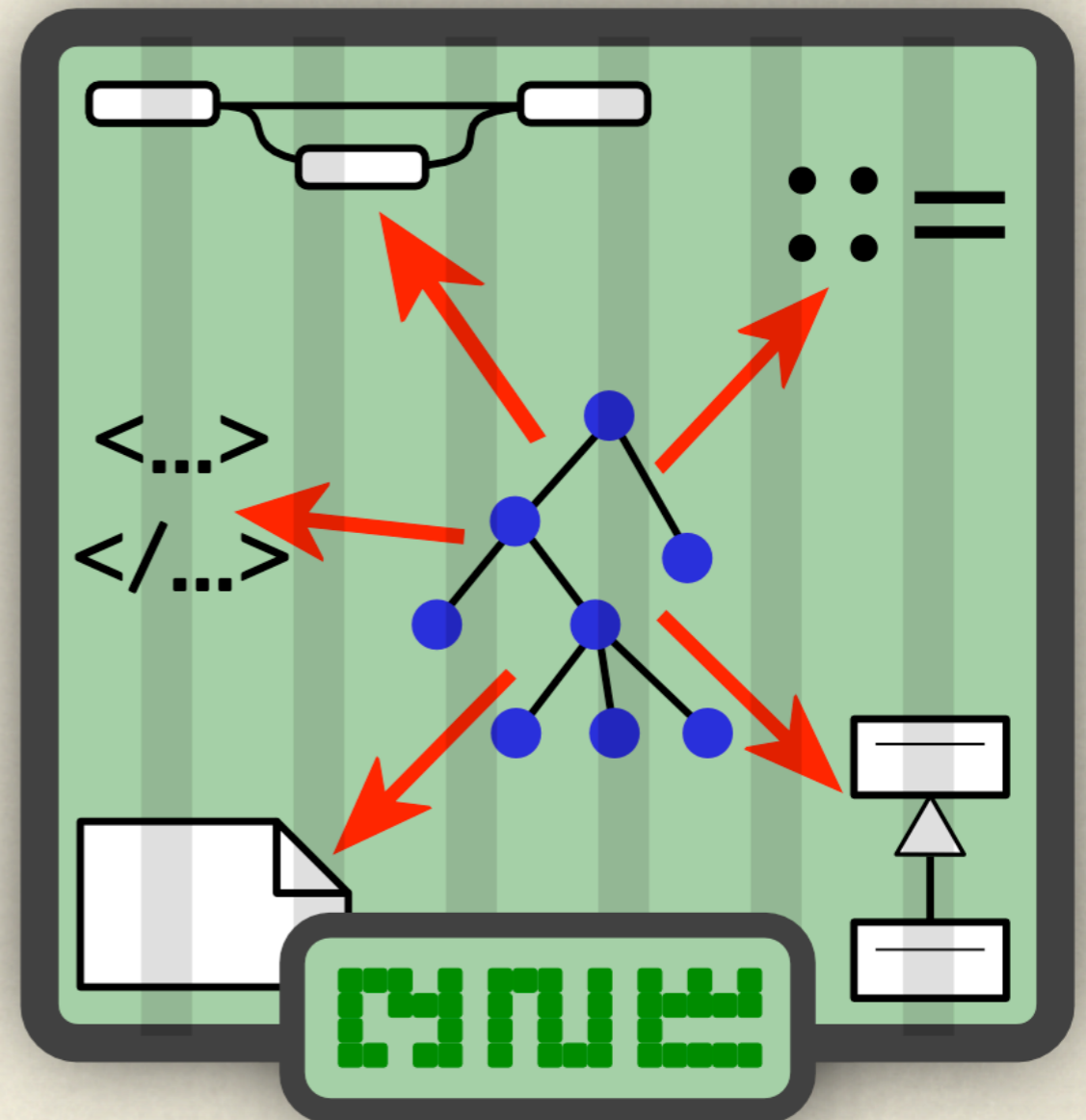
Notation-parametric recovery

- ◆ Perform a robust heuristic-based recovery process
- ◆ Successful for grammars of Ada, C, C++, C#, Dart, Eiffel, Modula, MediaWiki, LLL, EBNF, etc.



Notation-parametric recovery

- ◆ Perform a robust heuristic-based recovery process
- ◆ Successful for grammars of Ada, C, C++, C#, Dart, Eiffel, Modula, MediaWiki, LLL, EBNF, etc.



**Grammar
extraction
from other sources**

Other available extractors

- ◆ Eclipse Modeling Framework

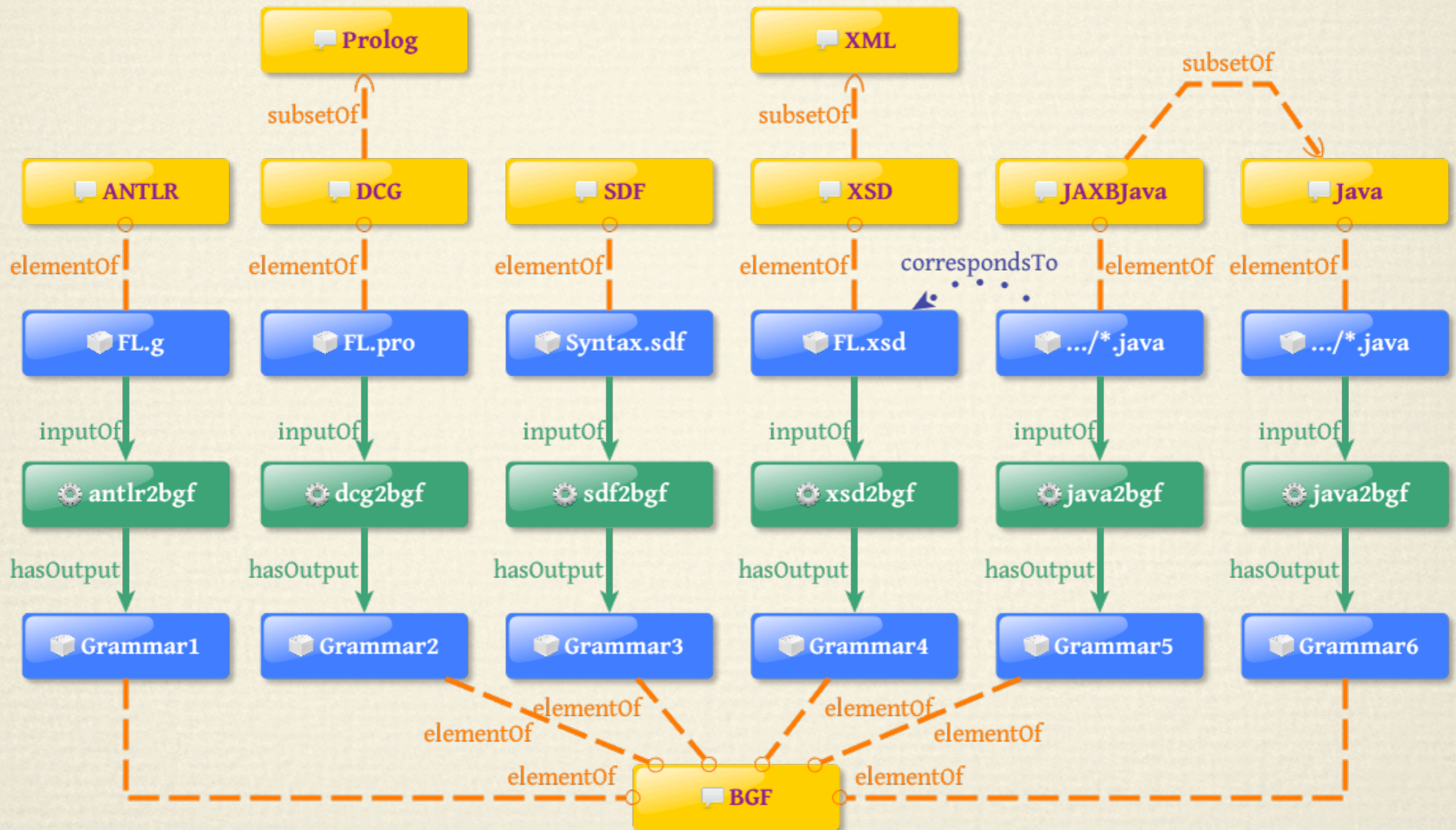
```
GGrammar bgf  
    = ecore2bgf( | cwd:///metamodel.ecore | );
```

- ◆ XML Schema Definition

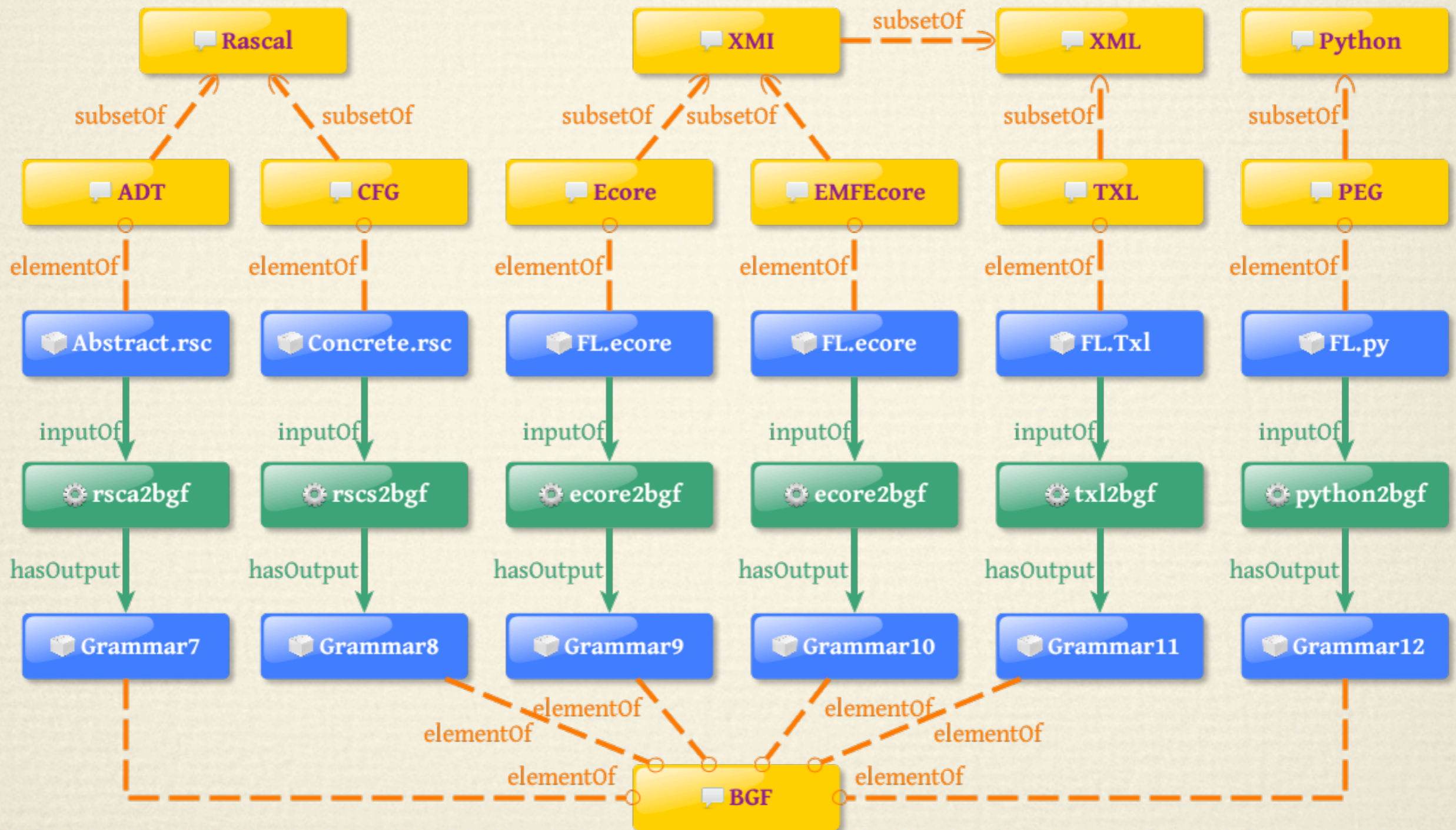
```
GGrammar bgf  
    = xsd2bgf( | cwd:///schema.xsd | );
```

Not all extractors have already been migrated to use Rascal

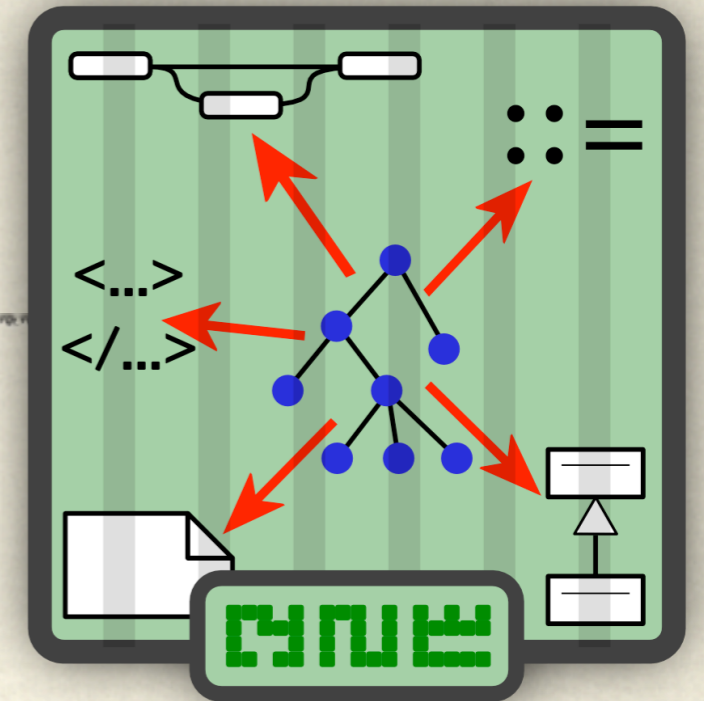
Flashy megamodel!



Flashy megamodel!!!



Grammar Zoo



◆ 569 grammars

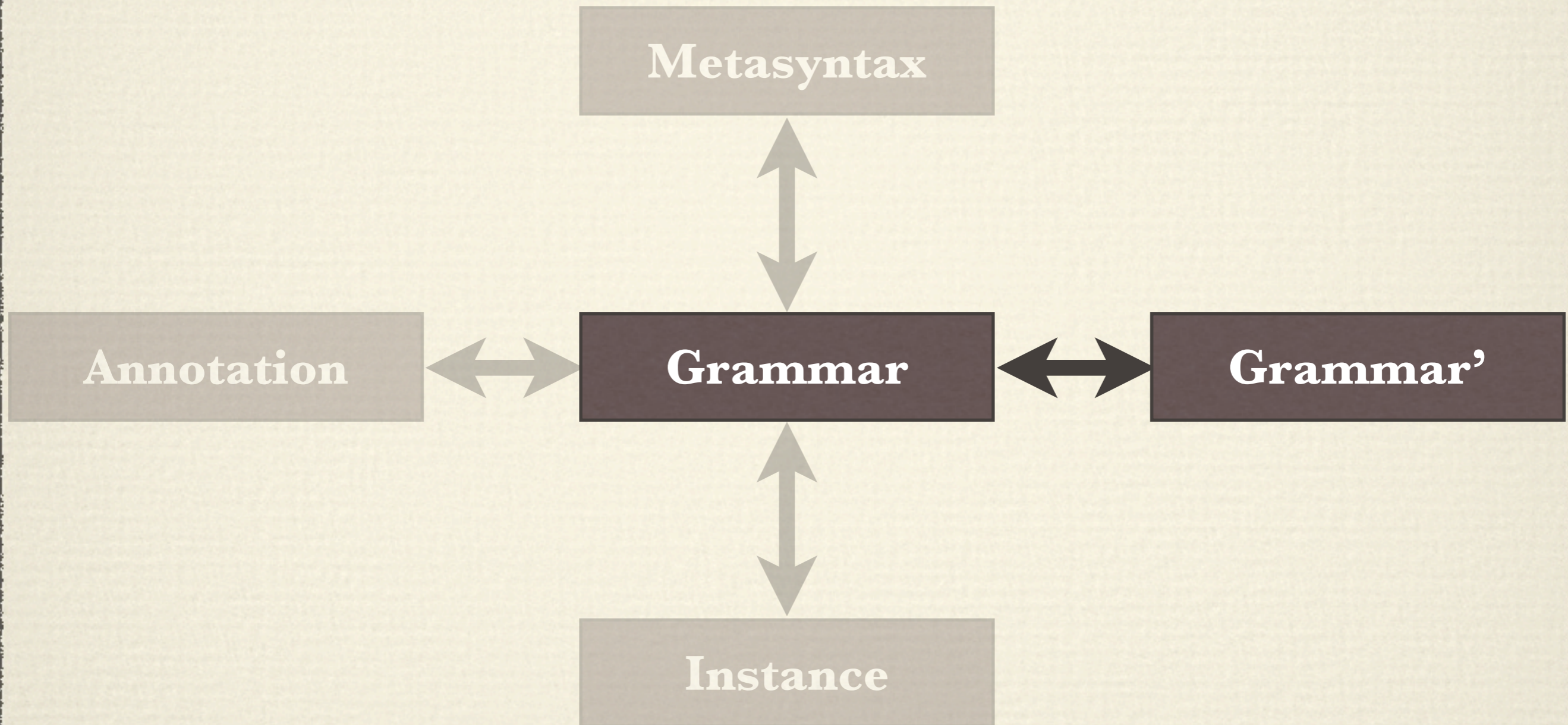
◆ Ada, ASM, ATL, ATOM, BASIC, BibTeX, C, C++, C#, CSV, Dart, Dot, DTD, EBNF, Eiffel, Fortran, HTML, Java, JavaScript, JSON, Modula, Occam, ODF, Pascal, PIF, PL/I, SVG, UML, Wiki, XML, XPath, XSLT, XQuery, ...

Grammar evolution

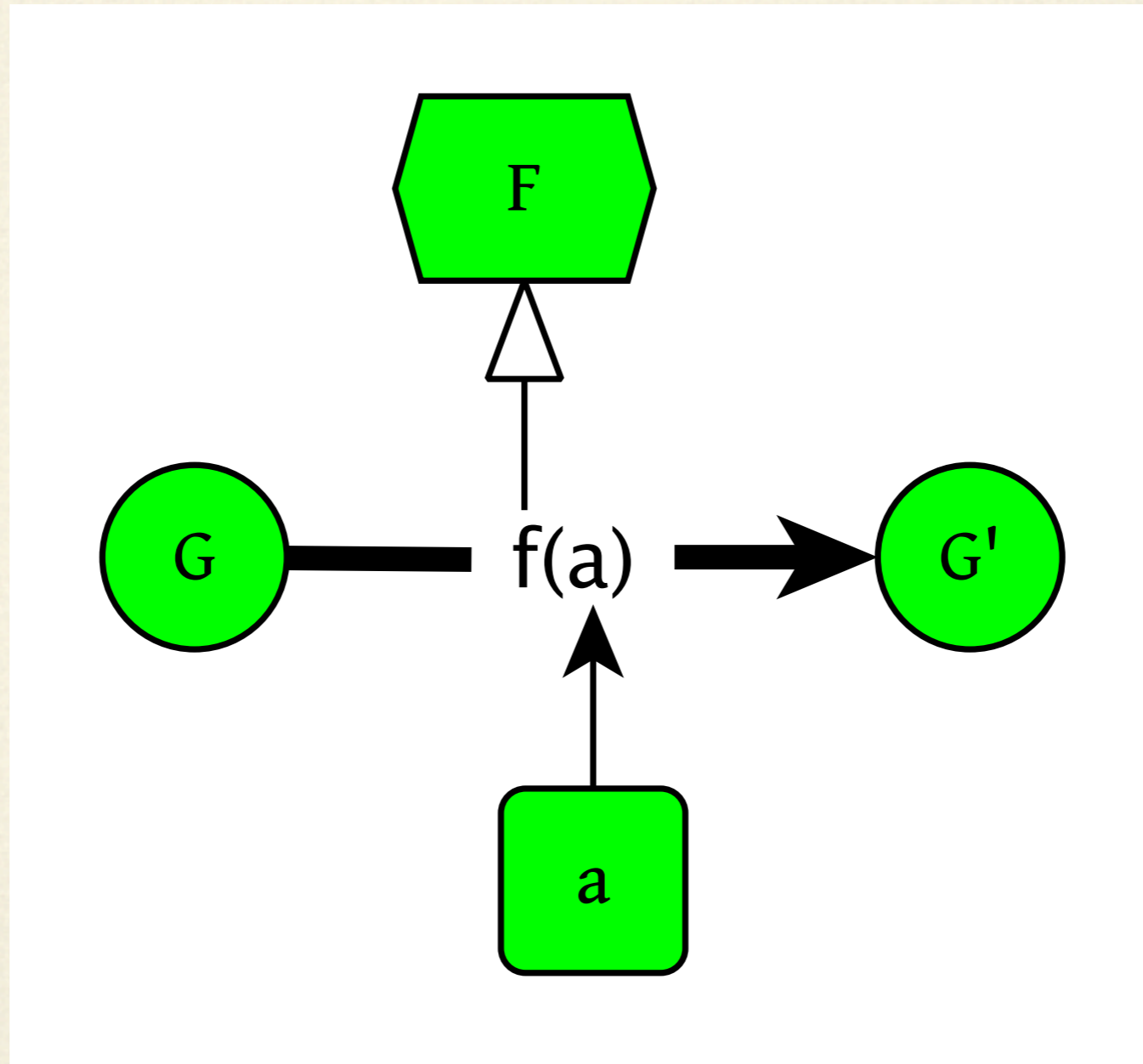
Grammar evolution

- ◆ Any change is a transformation
- ◆ Why transform?
 - ◆ adaptation
 - ◆ normalisation
 - ◆ beautification
 - ◆ inconsistency management
 - ◆ version control
- ◆ Documented, well-understood, compositional change
- ◆ Good for representing relationships

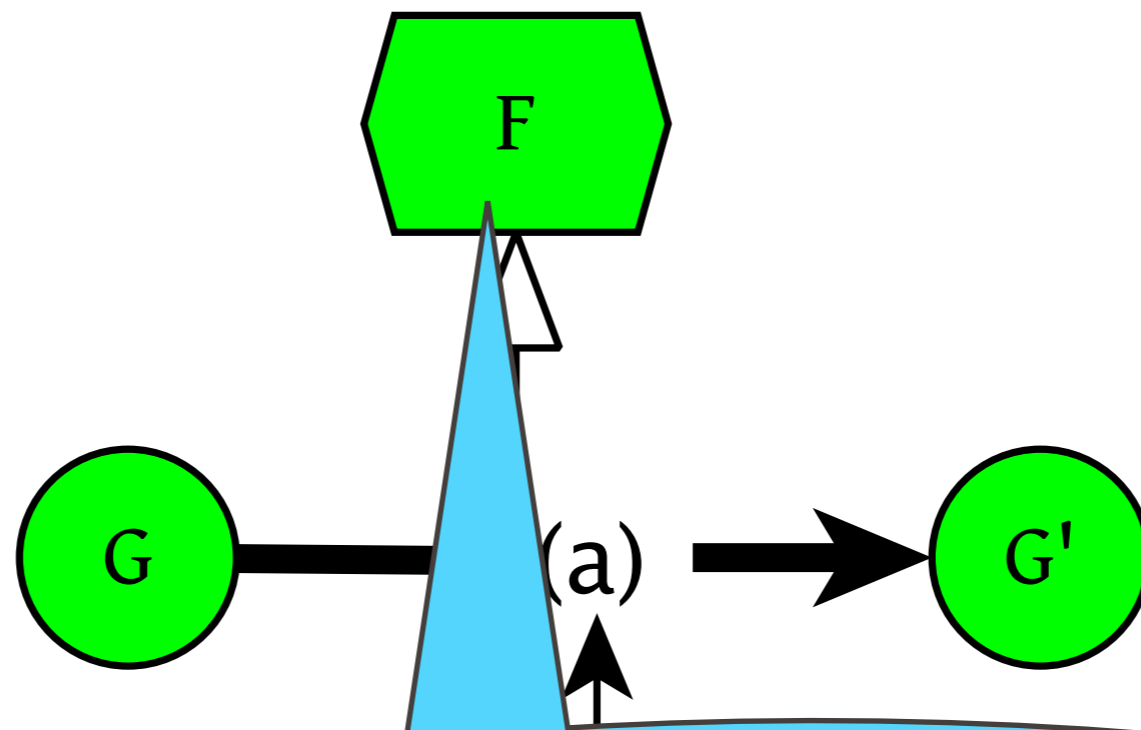
Grammar evolution



Transformation components



Transformation components

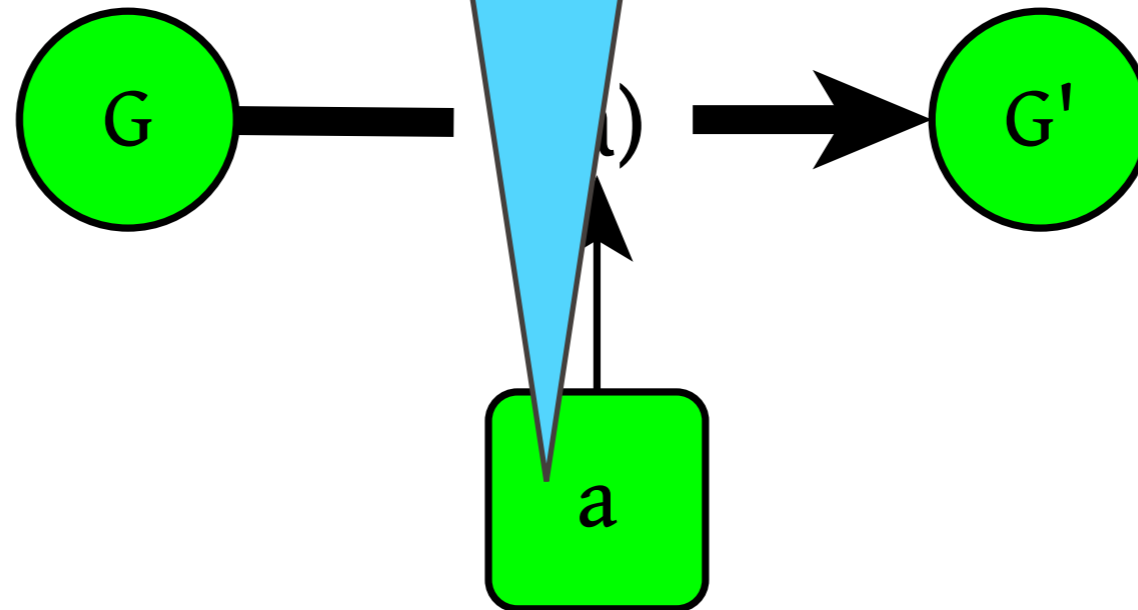


- known semantics, well-defined algorithm
- rename, fold, factor, inject, remove, ...

Transformation components

Arguments

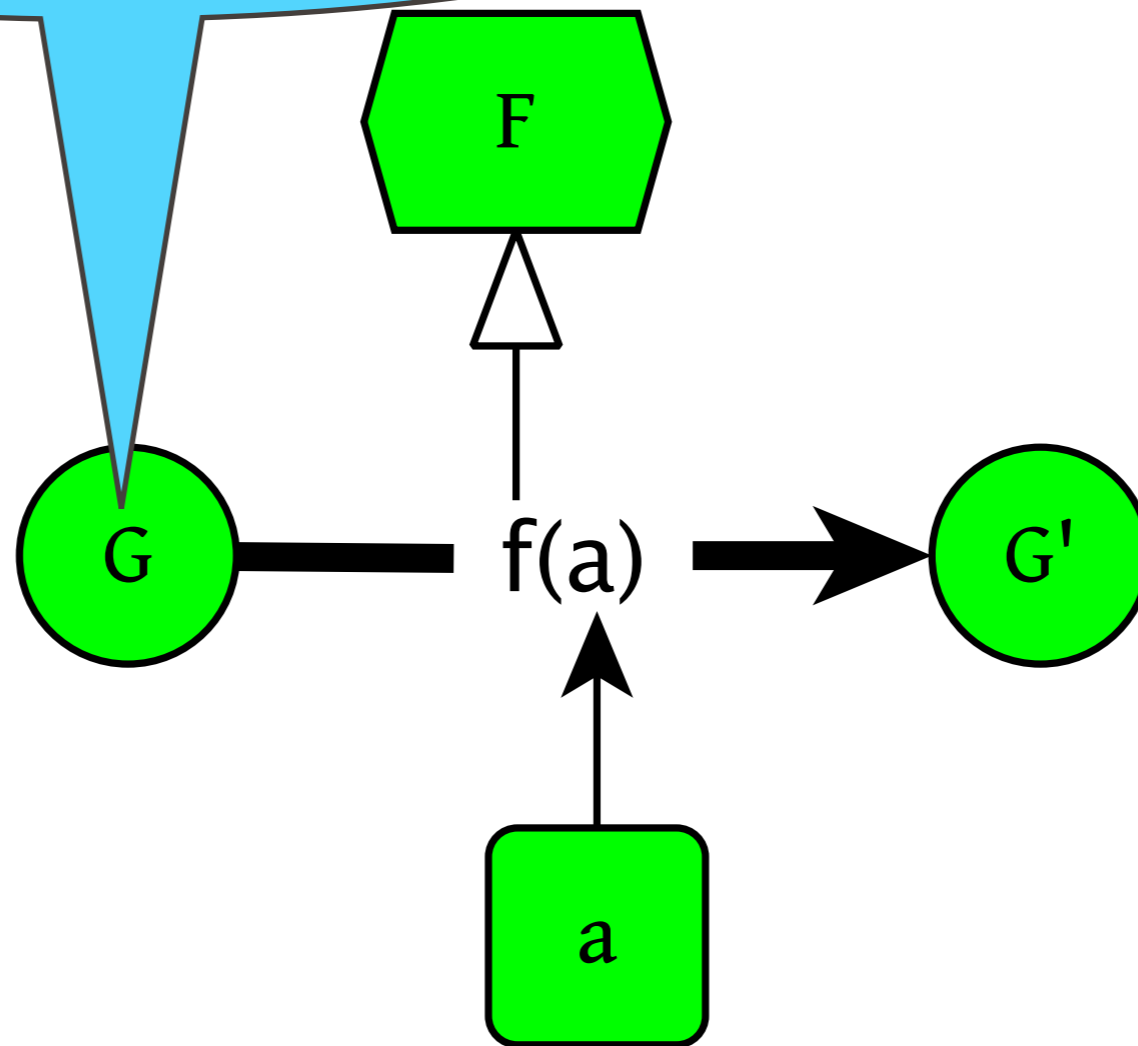
- what exactly to rename/factor/inject/...?



Transformation components

Input grammar

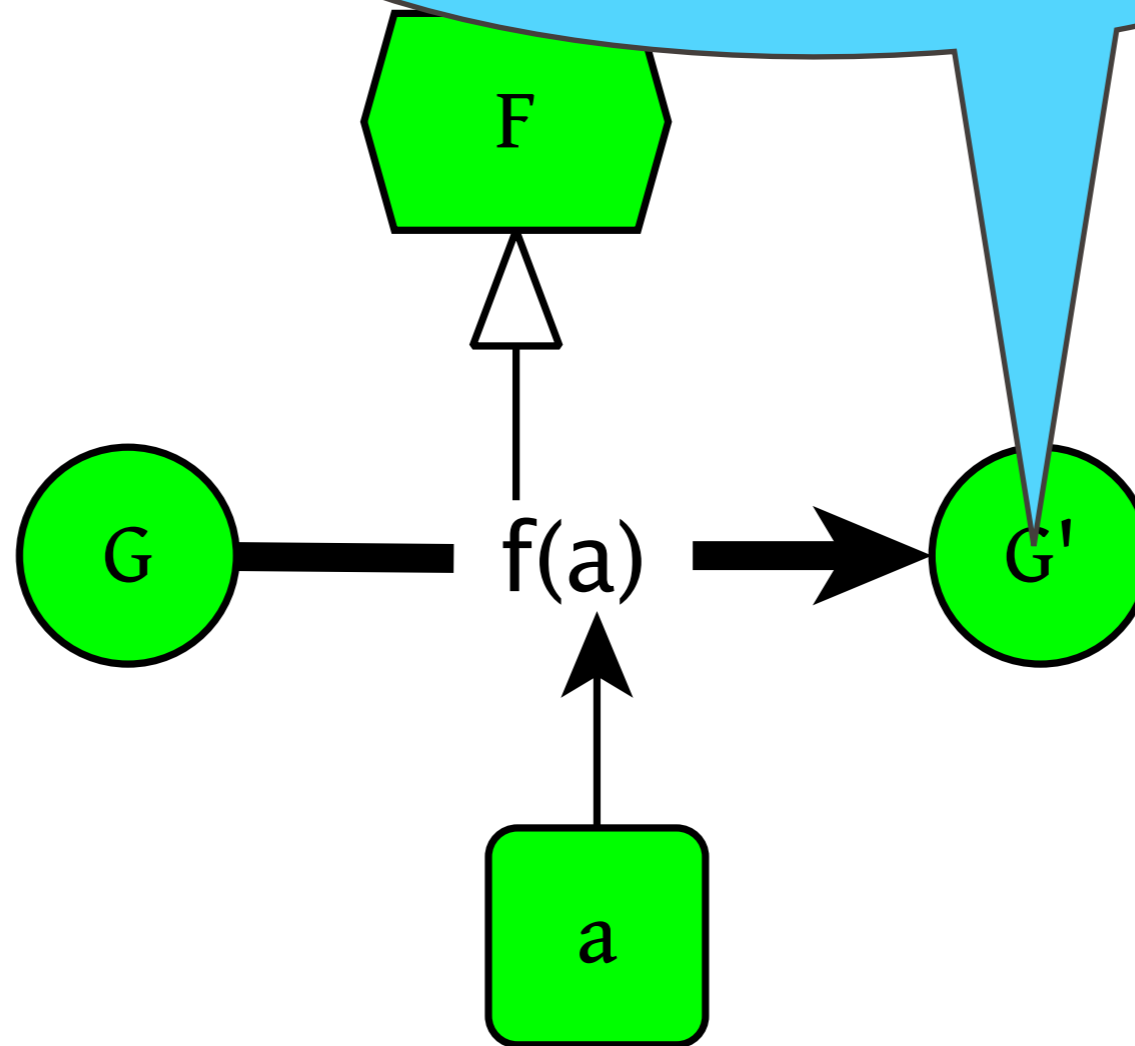
- determines applicability



Transformation components

Output grammar

- inferred result



Excerpt from the operator suite

Operator [LZ09, LZ11, ...]	Group [LZ11]	Class [CREP08]	Delta [GKP07]
<code>bypass()</code>	preserving	update	not breaking
<code>factor(x, y)</code>	preserving	update	not breaking
<code>introduce(n ::= rhs)</code>	preserving	additive	not breaking
<code>eliminate(n)</code>	preserving	subtractive	not breaking
<code>widen(x, y)</code>	increasing	additive	not breaking
<code>define(n ::= rhs)</code>	revising	additive	not breaking
<code>renameN(a, b)</code>	preserving	update	resolvable
<code>fold(n)</code>	preserving	additive	resolvable
<code>extract(n ::= rhs)</code>	preserving	additive	resolvable
<code>unfold(n)</code>	preserving	subtractive	resolvable
<code>disappear(p, m)</code>	decreasing	subtractive	resolvable
<code>permute(p, q)</code>	revising	update	resolvable
<code>concretize(p, m)</code>	revising	additive	resolvable
<code>abstractize(p, m)</code>	revising	subtractive	resolvable
<code>project(p, m)</code>	revising	subtractive	resolvable
<code>narrow(x, y)</code>	decreasing	subtractive	unresolvable
<code>inject(p, m)</code>	revising	additive	unresolvable
<code>replace(x, y)</code>	revising	—	unresolvable

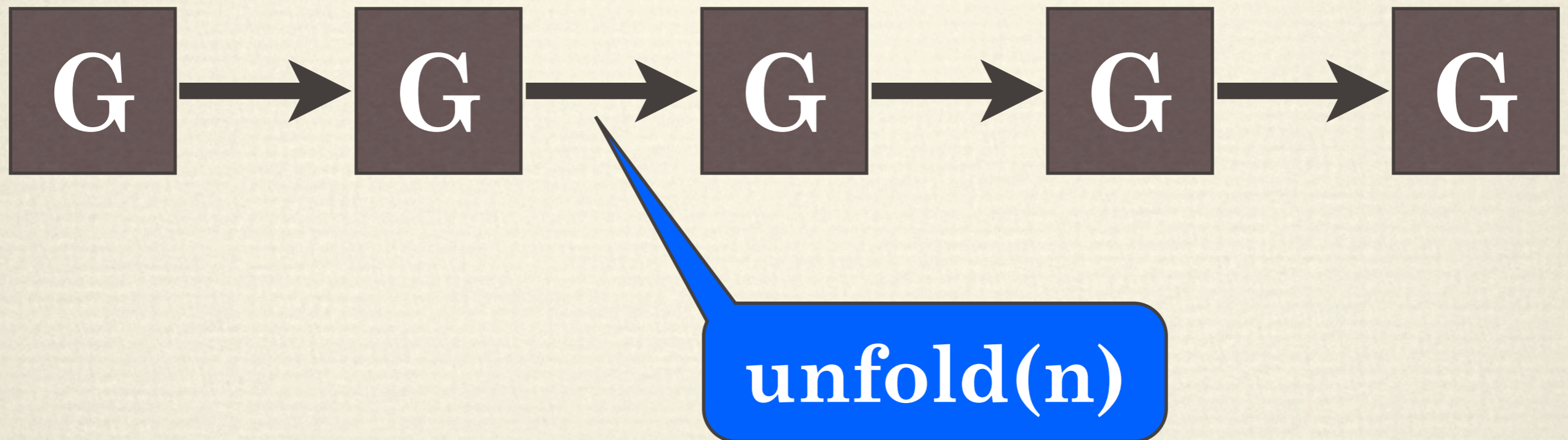
Transform to converge

	jls1	jls12	jls123	jls2	jls3	read12	read123	Total
Number of lines	682	5114	2847	6774	10721	1639	3082	30859
Number of transformations	67	290	111	387	544	77	135	1611
○ Semantics-preserving (§4.2.2)	45	231	80	275	381	31	78	1121
○ Semantics-increasing/-decreasing	22	58	31	102	150	39	53	455
○ Semantics-revising	—	1	—	10	13	7	4	35
Preparation phase (§4.2.1)	1	—	—	15	24	11	14	65
○ Known bugs	—	—	—	1	11	—	4	16
○ Post-extraction	—	—	—	7	8	7	5	27
○ Initial correction	1	—	—	7	5	4	5	22
Resolution phase	21	59	31	97	139	35	43	425
○ Extension (§4.2.3)	—	17	26	—	—	31	38	112
○ Relaxation (§4.2.4)	18	39	5	75	112	—	2	251
○ Correction (§4.2.5)	3	3	—	22	27	4	3	62

Grammar programming



Grammar programming

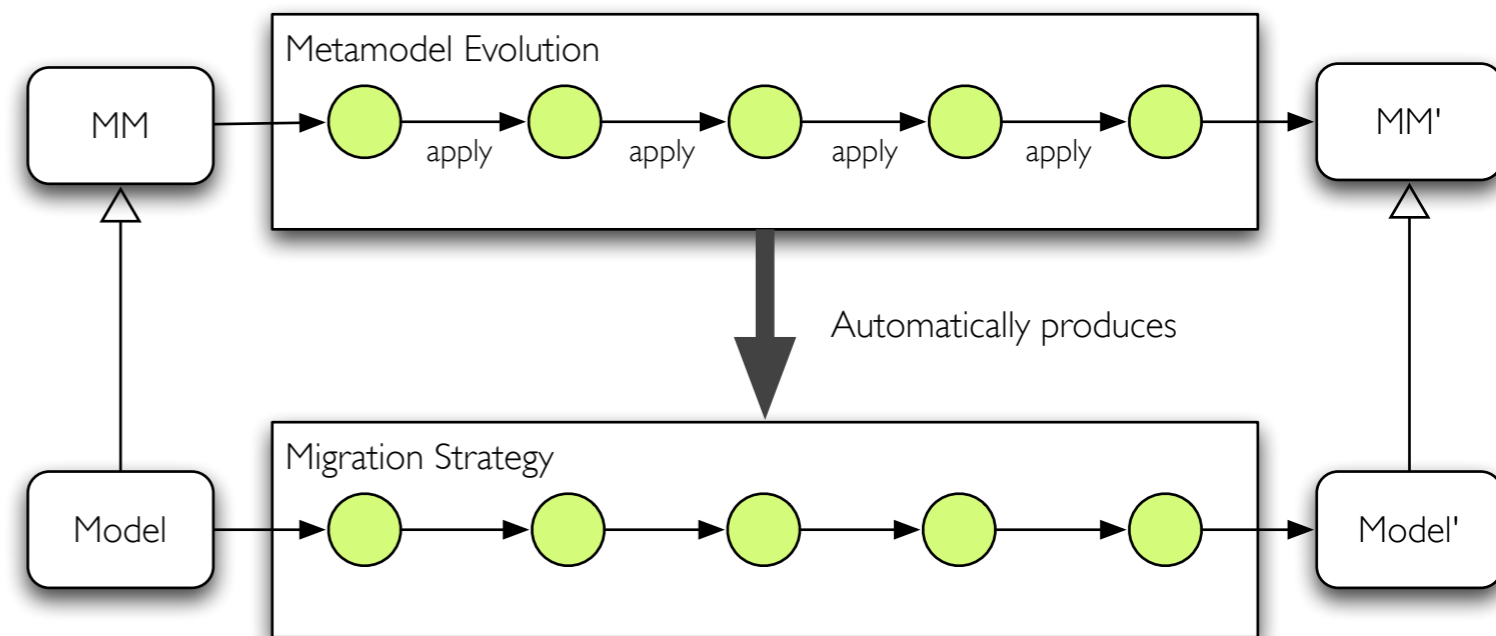


Meanwhile in modelware...

Model Migration Approaches:

Operator-based co-evolution

- + migration strategy is “free”
- tool lock in
- operator set completeness
- reverse engineering challenging



A transformation sequence

expr ::= ...;
atom ::= ID | INT | '(' expr ')';

expr ::= ...;
expr ::= ID;
expr ::= INT;

abstractize

abridge

expr ::= ...;
atom ::= ID | INT | expr;

expr ::= ...;
expr ::= ID;
expr ::= INT;
expr ::= expr;

vertical

unite

expr ::= ...;
atom ::= ID;
atom ::= INT;
atom ::= expr;

RelaxNG2BGF.rsc

DOM.rsc

XBGF.rsc

XMLSchemaDefinitio

XSD.rsc

Differ.rsc

Location.rsc

```
683 bool case_deyaccify_left_plus(bool debug)
684 {
685     GGrammar bgf = grammar(["foo"],
686         [production("foo",
687             nonterminal("bar")),
688         production("foo",
689             sequence([
690                 nonterminal("foo"),
691                 nonterminal("bar")])]),
692         []);
693     XSequence xbgf =
694         [deyaccify("foo")];
695     GGrammar bl = grammar(["foo"],
696         [production("foo",
697             plus(nonterminal("bar"))]),
698         []);
699     return run_case(bgf, xbgf, bl, debug);
700 }
701 test bool test_deyaccify_left_plus() = case_deyaccify_left_plus(false);
702 void show_deyaccify_left_plus() {case_deyaccify_left_plus(true);}
```

Package Explo Plug-ins

- AmibiDexter
- Antlr
- grammarlab
- gw1rp4sd
- java2bgf
- org.eclipse.imp.pdb
- org.eclipse.imp.pdb.ui
- org.eclipse.imp.pdb.values
- org.eclipse.imp.runtime
- org.xtext.example.ebnf
- org.xtext.example.ebnf.generator
- org.xtext.example.ebnf.ui
- rascal
- rascal-eclipse
- rascal-shell

```
683 bool case_deyaccify_left_plus(bool debug)
684 {
685     GGrammar bgf = grammar(["foo"],
686         [production("foo",
687             nonterminal("bar")),
688         production("foo",
689             sequence([
690                 nonterminal("foo"),
691                 nonterminal("bar")])]),
692         []);
693     XSequence xbgf =
694         [deyaccify("foo")];
695     GGrammar bl = grammar(["foo"],
696         [production("foo",
697             plus(nonterminal("bar")))],
698         []);
699     return run_case(bgf xbgf bl debug);
```

Error Log Tasks Problems Console Tutor

Store history Terminate Interrupt Trace

Rascal [DEBUG, grammarlab]

rascal>import tests::transform::XBGF;

ok

rascal>:test

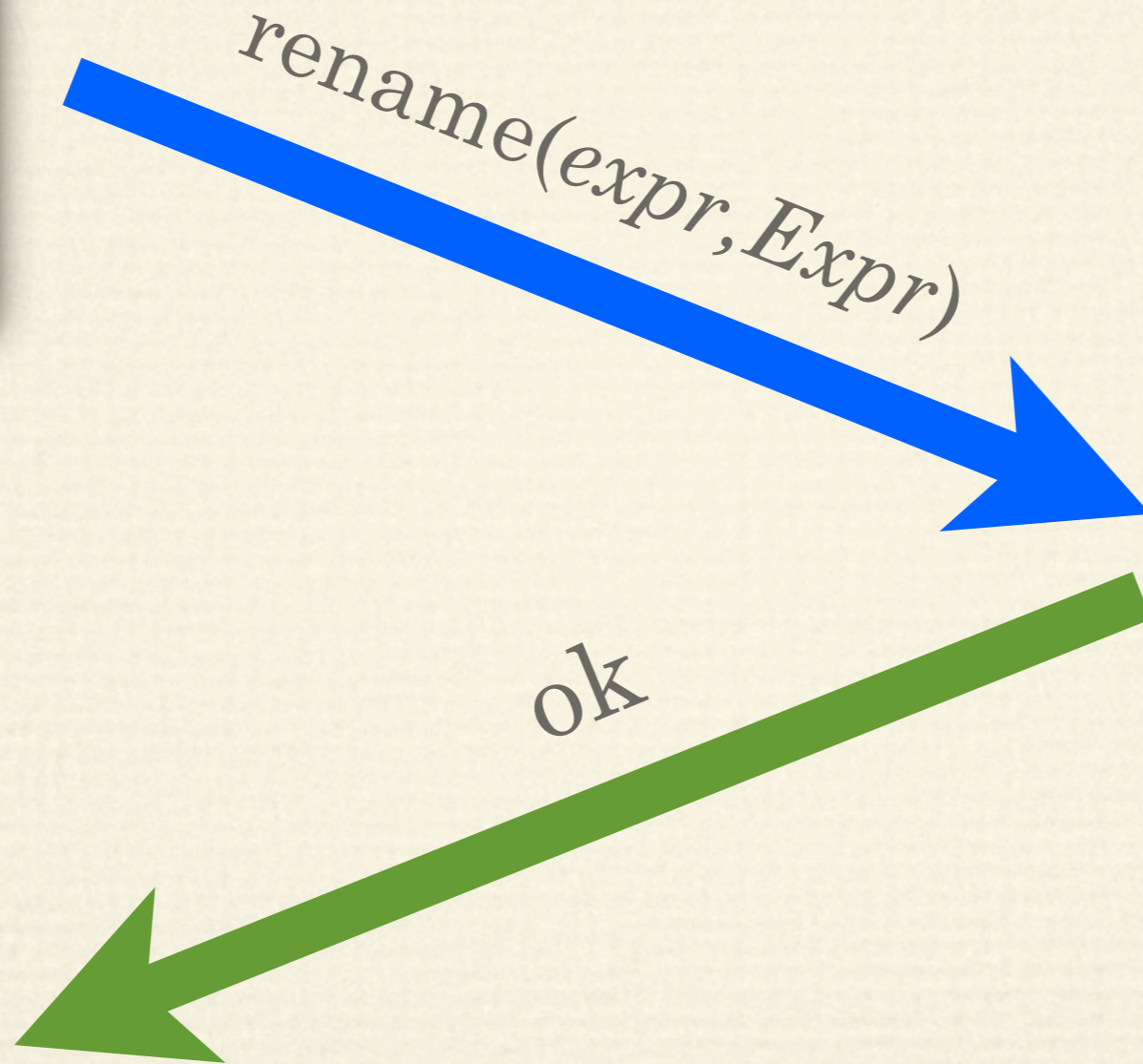
ok

rascal>



Negotiated transformations

Rename a nonterminal: success



Rename a nonterminal: failure



Rename a nonterminal: negotiation



Rename a nonterminal: negotiation



adjustment

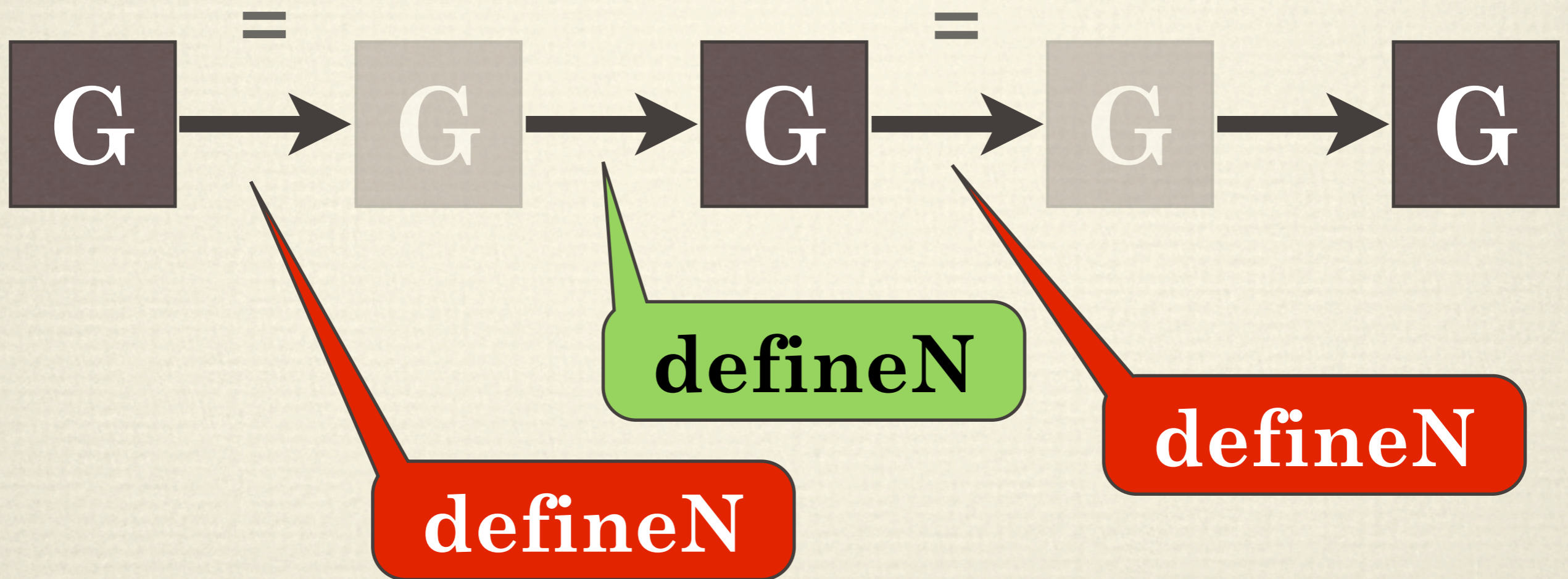
KEEP
CALM
AND
CONSIDER
IT
DONE

Pending evolution

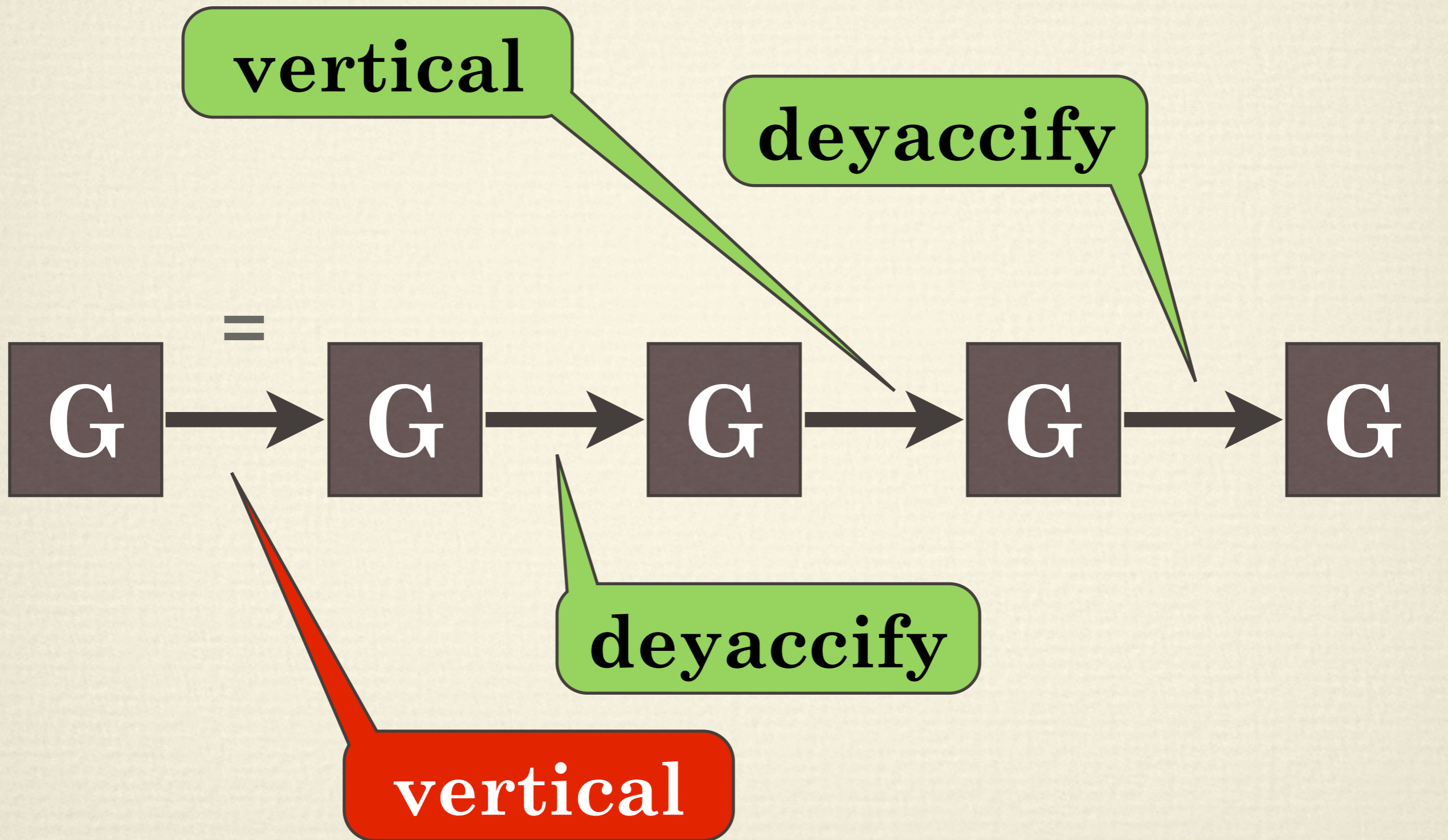
Grammar programming



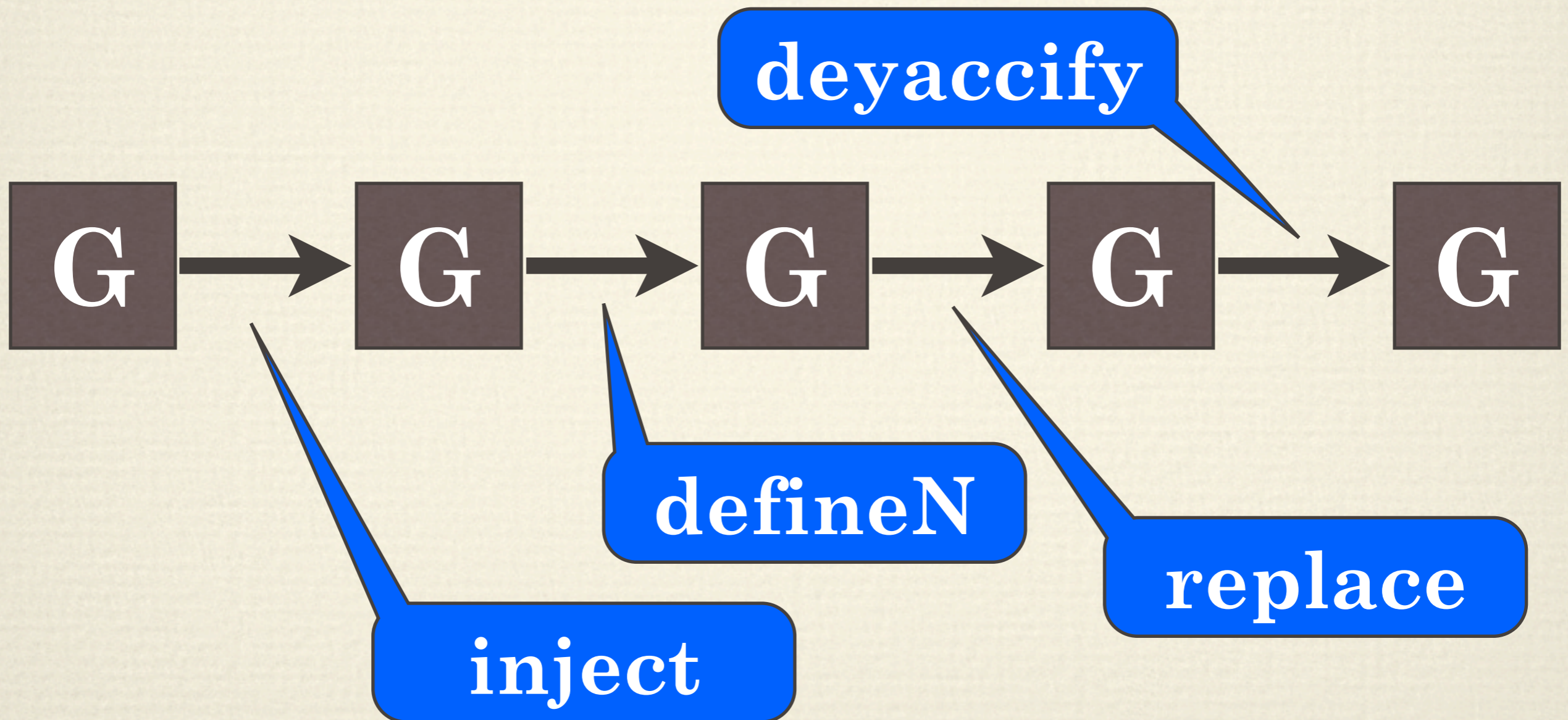
Optional execution



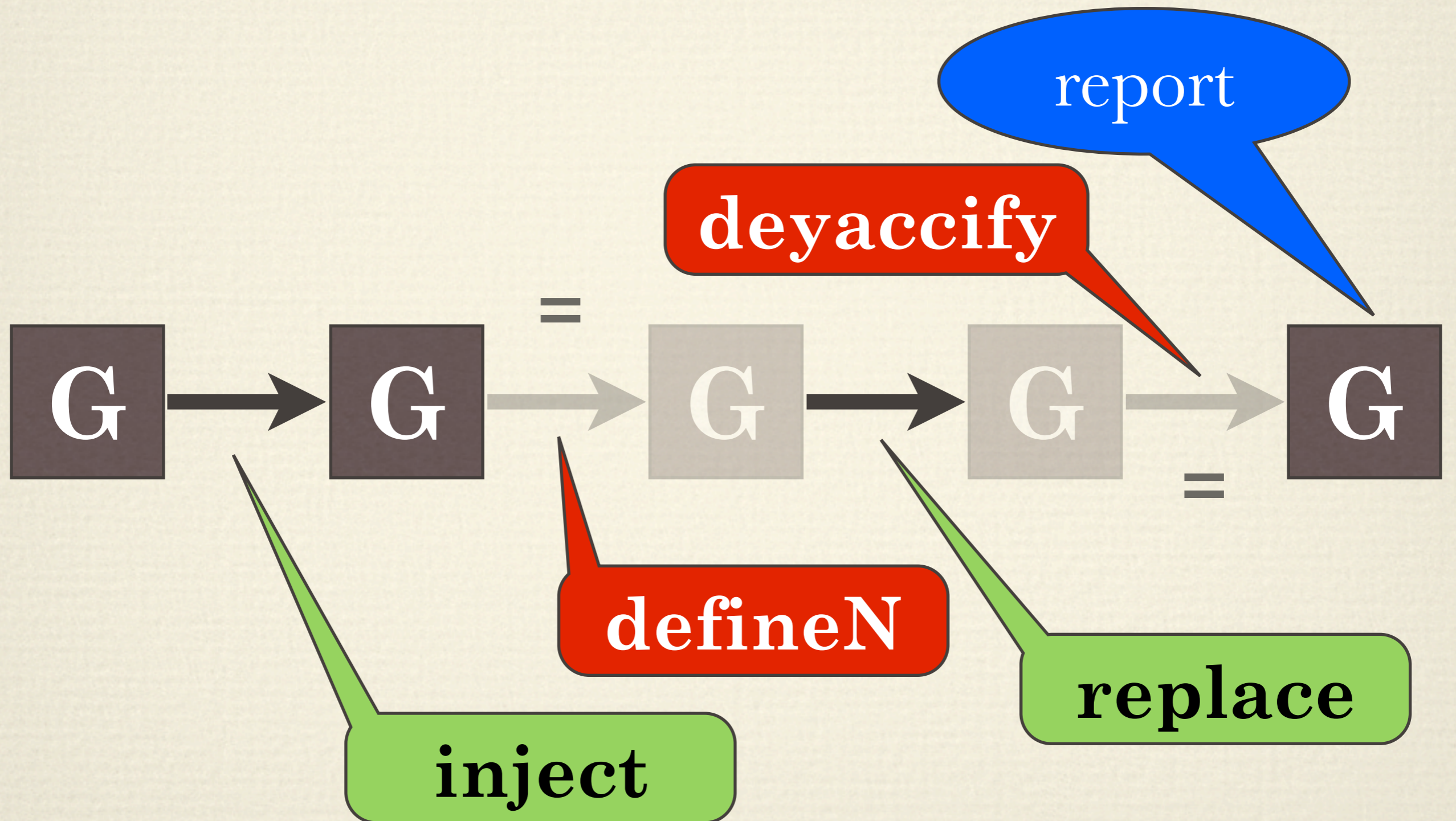
Asserting preconditions



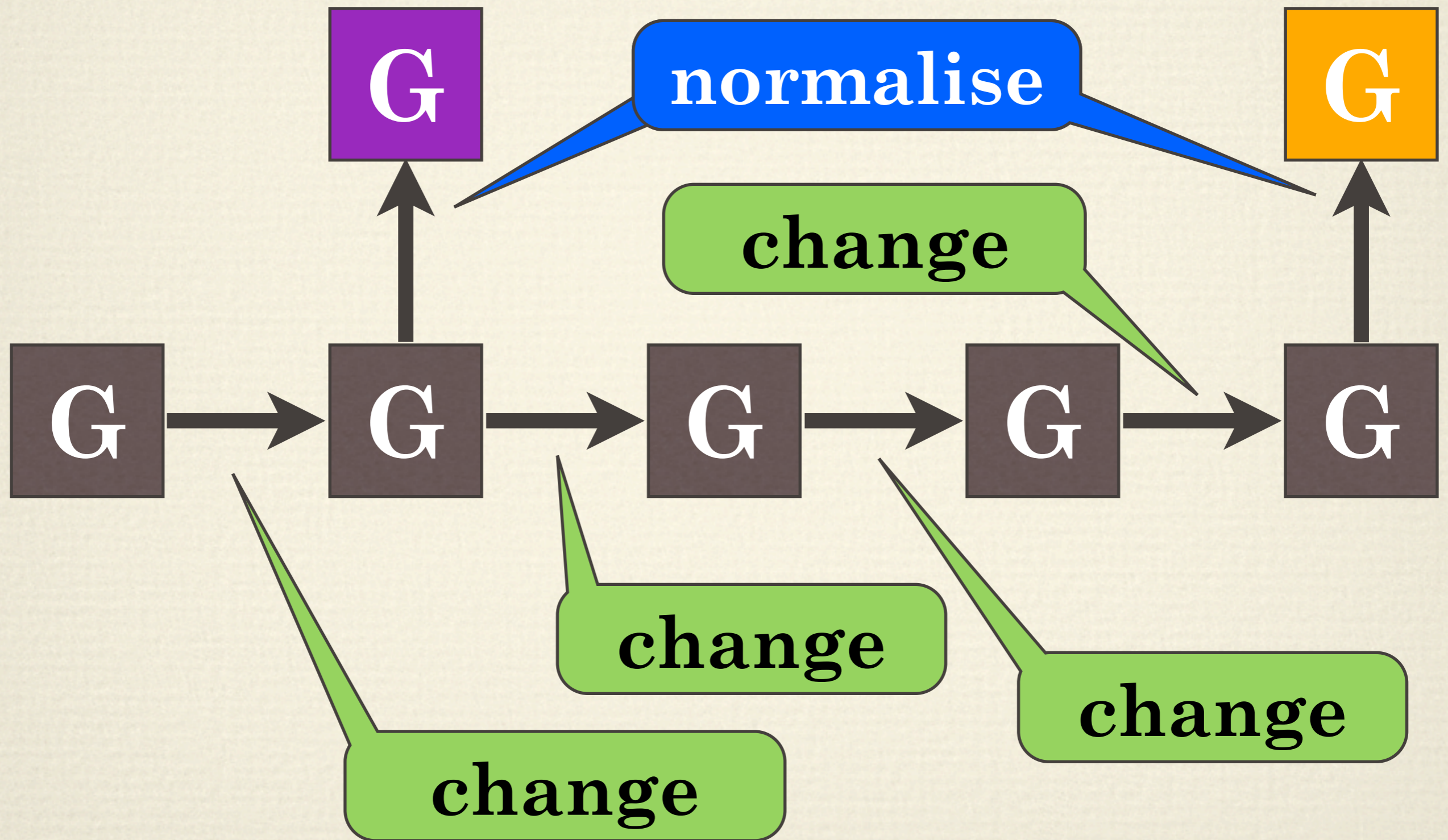
Reuse correction scripts



Reuse correction scripts



Normalise before export

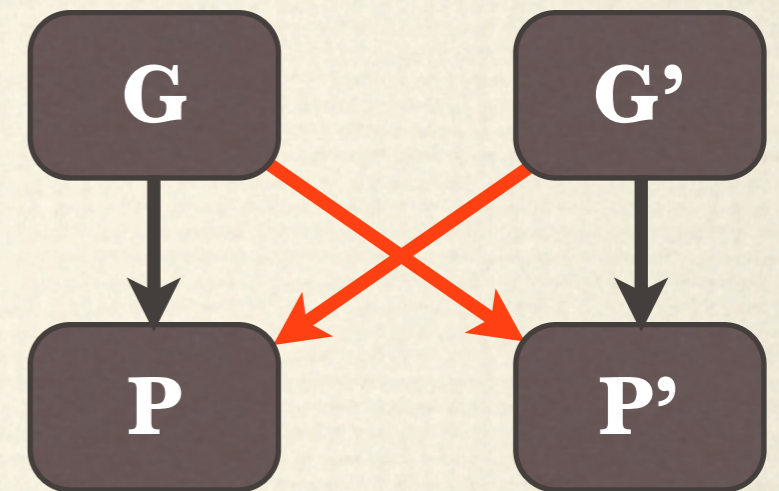


Operator-based evolution analysis

Grammar testing

Grammar-based testing

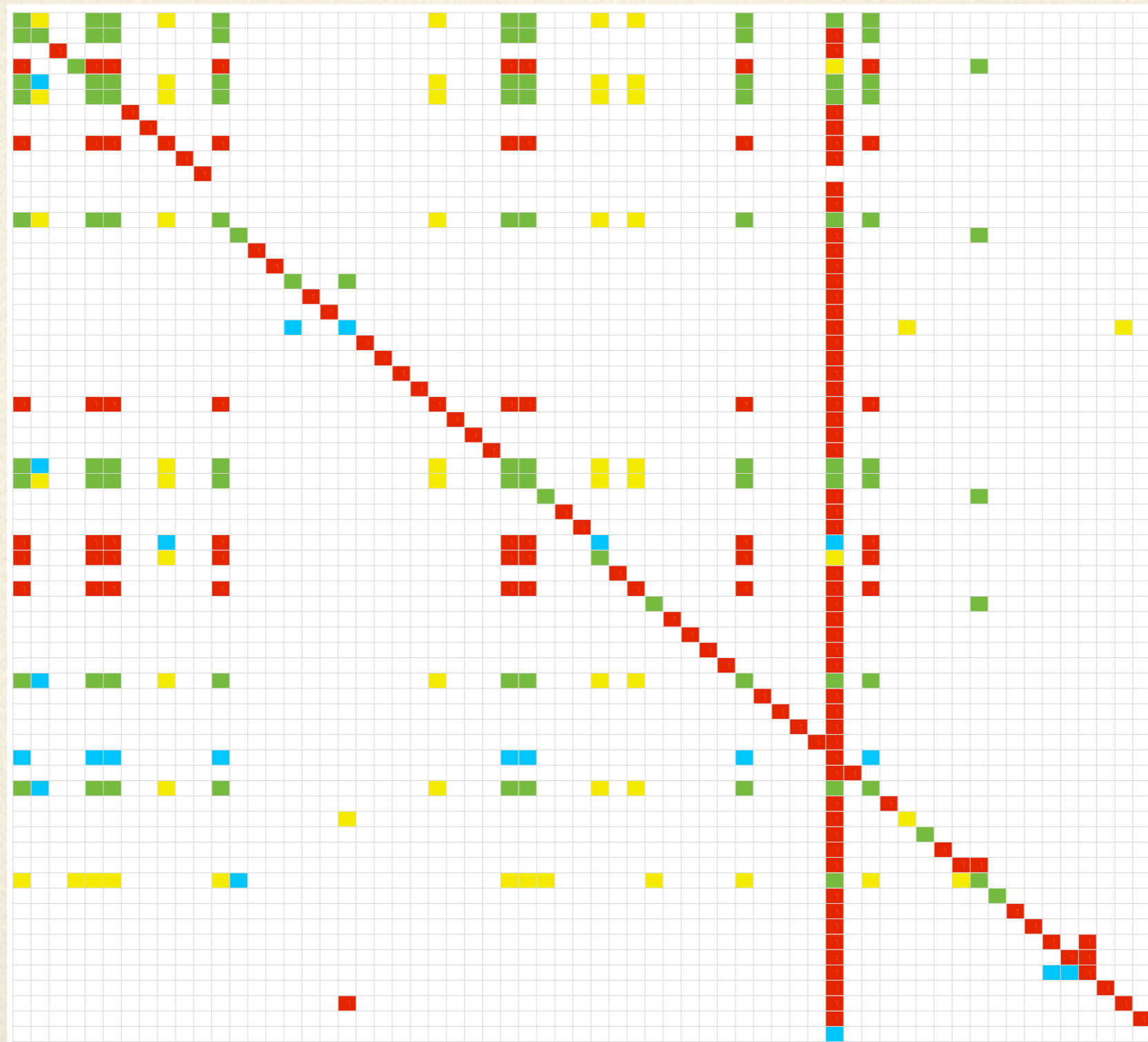
- ◆ Asymmetric comparison:
 - ◆ Reference grammar vs. parser under test
- ◆ Symmetric comparison:
 - ◆ Differential testing
 - ◆ Systematic test data generation
 - ◆ Controlled combinatorial coverage
 - ◆ Larger sets of smaller test data items
 - ◆ Nonterminal matching
 - ◆ Non-context-free effects



Coverage criteria

- ◆ **Trivial coverage:** if the test data set is not empty.
- ◆ **Nonterminal coverage:** if each nonterminal is exercised at least once.
- ◆ **Production coverage:** if each production rule in the grammar is exercised at least once.
- ◆ **Branch coverage:** each branch of $? | * +$
- ◆ **Unfolding coverage:** each production of each right hand side nonterminal occurrence
- ◆ **Context-dependent branch coverage!**

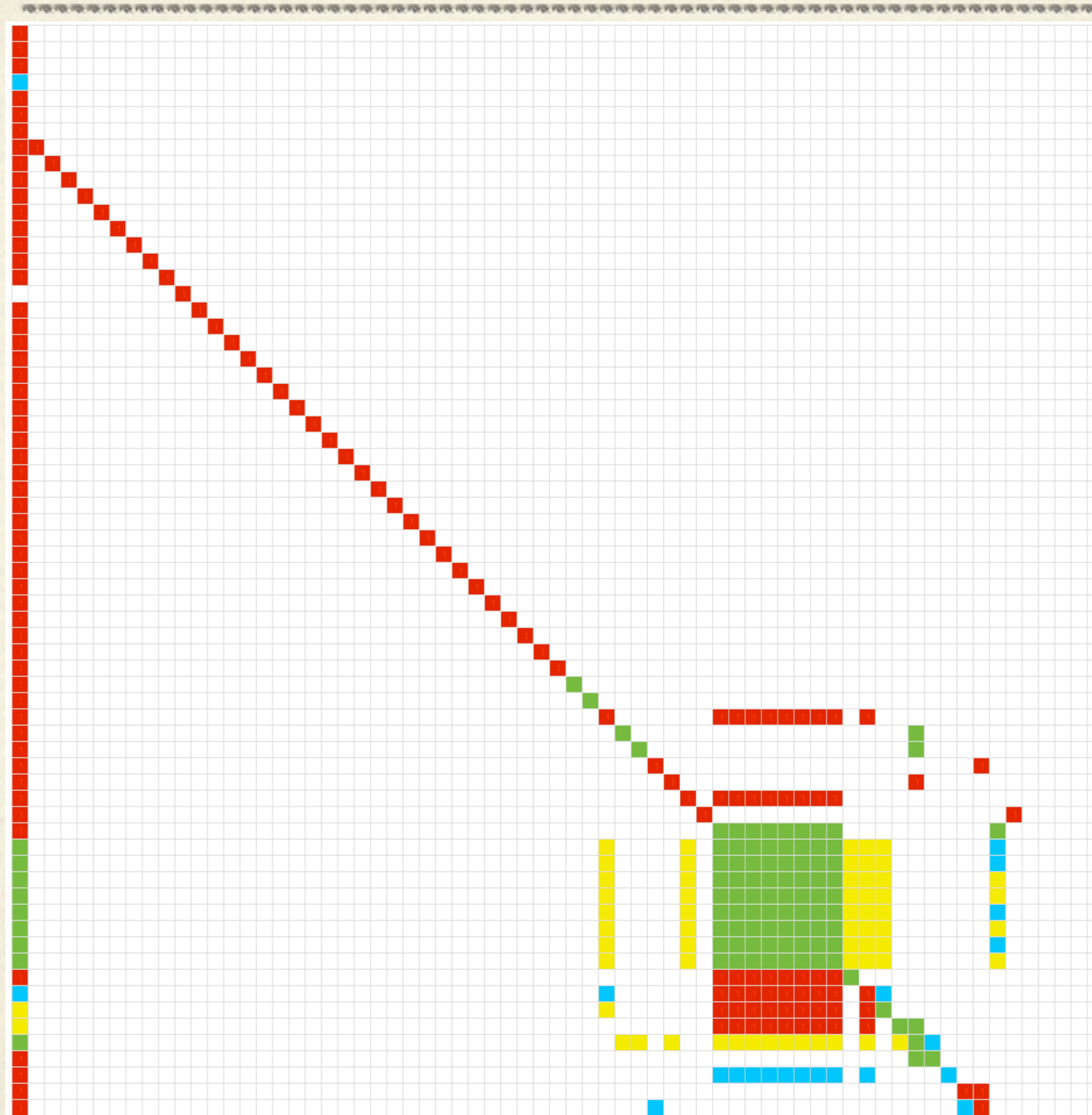
Nonterminal matching



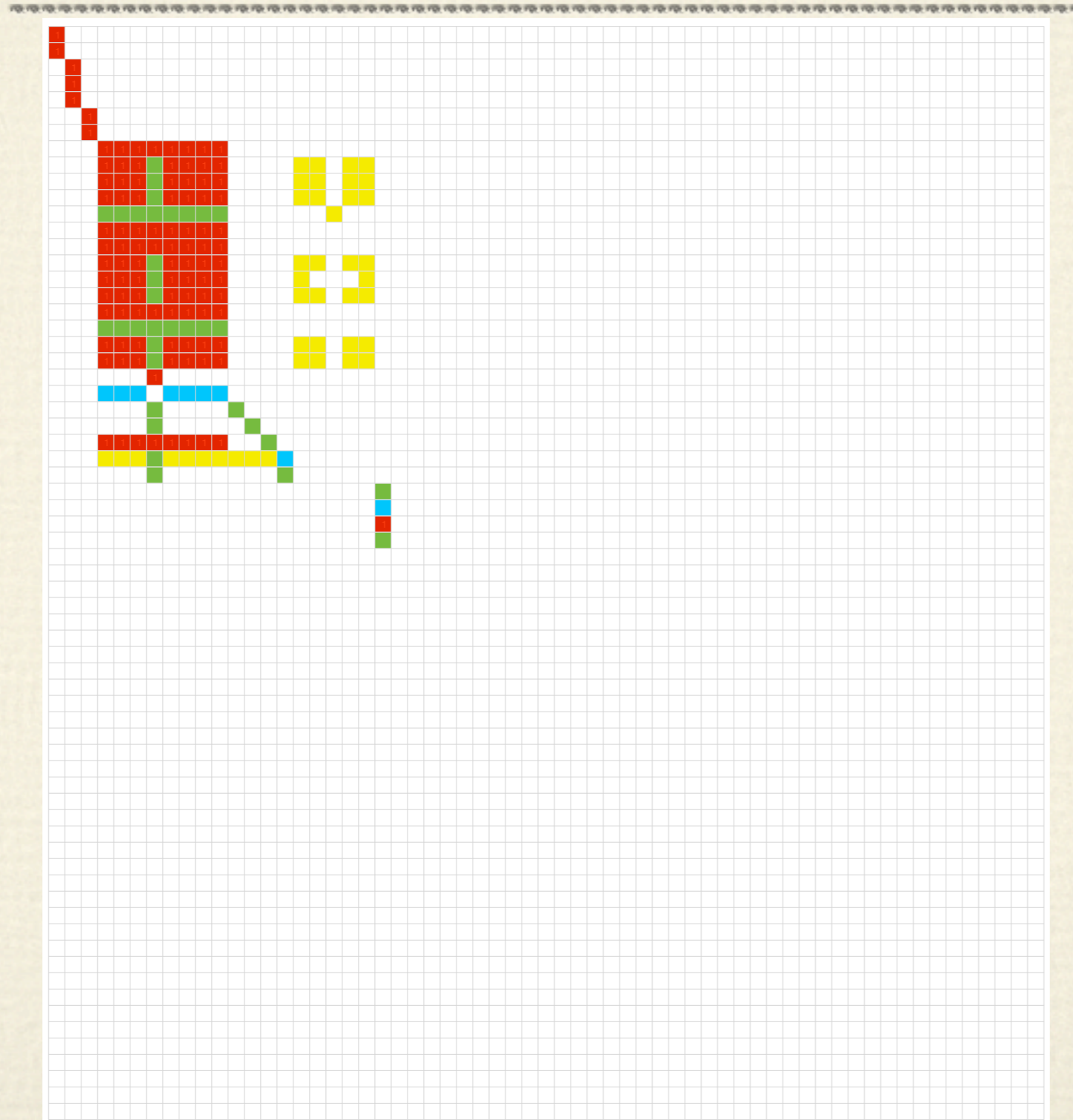
Nonterminal matching



Nonterminal matching



Nonterminal matching

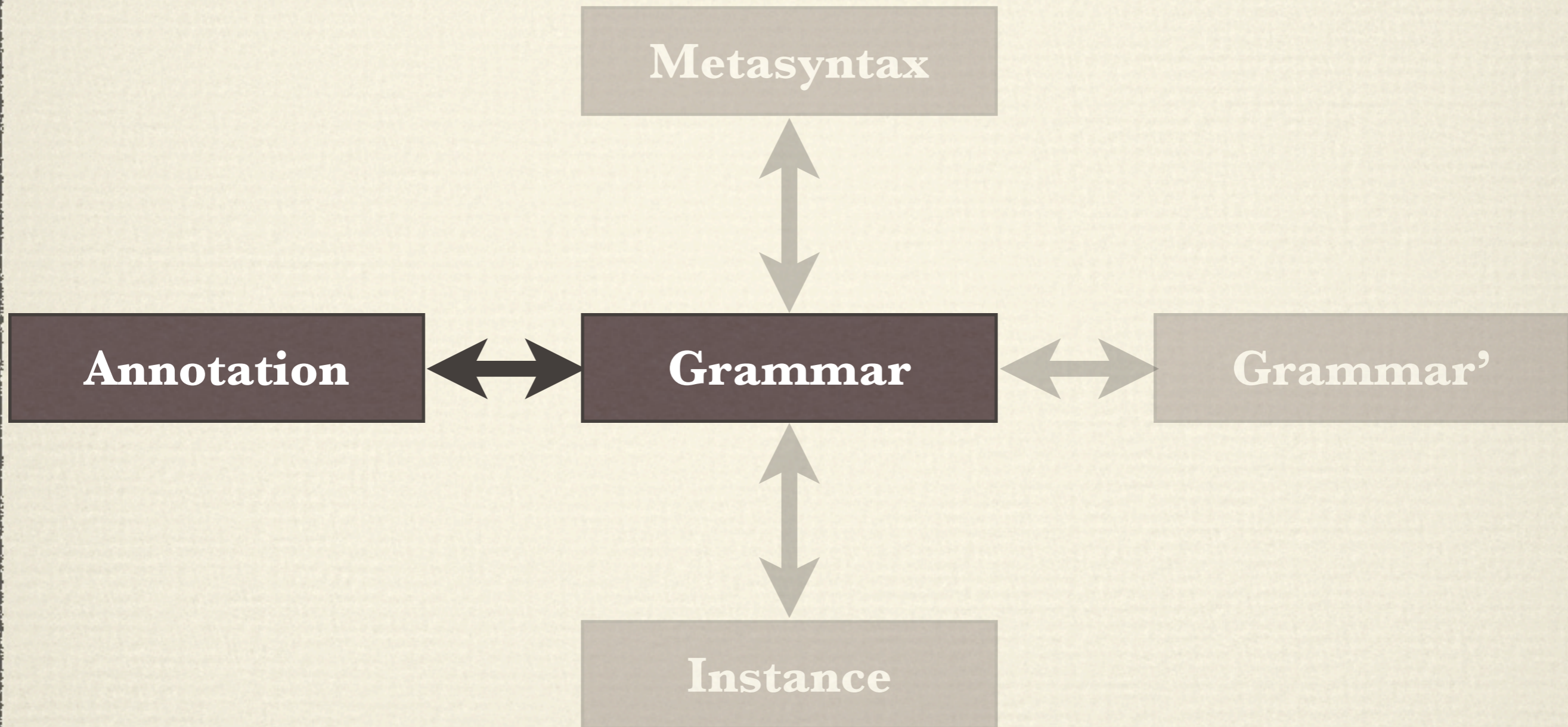


**Language
documentation**

Language documentation

- ◆ Given are:
 - ◆ a grammar for a software language
 - ◆ explanations in a natural language
 - ◆ executable code samples
 - ◆ known relationships between concepts
- ◆ How to do language documentation properly?
- ◆ By generating it from structured data!

Language documentation



Documenting grammars

Language Standardization Needs Grammarware

Steven Klusener, Vadim Zaytsev
Vrije Universiteit, De Boelelaan 1081a, NL-1081 HV Amsterdam
Email: {steven,vadim}@cs.vu.nl

October 4, 2005

Abstract

The ISO programming language standards are valuable documents that describe the syntax and semantics of mainstream languages. New features are proposed after thorough reviews by the standardization committees, leading to change documents that describe which modifications have to be enforced in the language standard document in order to actually add a new feature to the language. Maintaining these documents, both the language standard itself and all the change documents, is a time and resource consuming effort and in the evolution of these documents inconsistencies may be introduced. In this note we propose to utilize *grammarware*, a collection of new methods and new technology which can be used to support the advancement of these language documents in a more structured way. Besides, we will discuss how other tooling (like browsable language definitions, parser generators, pretty-printers, code checkers, etc.) can be obtained from the language standard. The final objective is threefold: (1) to facilitate the standardization committees in their activities and to raise the quality of the language standard documents; (2) to extend the usability of language standards by providing various presentations of each standard (in a human readable document, in a browsable form, in a machine readable BNF, etc.); (3) to help tool builders (compiler vendors, IDE vendors, etc.) in generating their parsing front-end, and to provide technology for tool builders to specify differences between their dialects and the actual standard.

1 Introduction

Software engineering is an established area, both in the IT industry and in the academic field of computer science. In the recent years a lot of progress within software engineering is made in (among others):

1. **Computer languages and specification formalisms**, with a continuous increase in the level of abstraction and declarativity;
2. **Tool support**. Software is not written in simple text editors anymore, but in advanced *Integrated Development Environments* (IDEs), like Microsoft Visual Studio, IBM WebSphere Studio and Eclipse, Borland Enterprise Studio, Anjuta DevStudio, etc.;

◆ IAL, FORTRAN

◆ ISO/IEC

◆ ECMA

◆ W3C

◆ OMG

◆ Design Patterns

◆ ...

Unified model for language docs

Domain concept	IAL [Bac60]	Jovial [MIL84]	Design Patterns [GHJV95]	Smalltalk [Sha97]	Informix [IBM03]	C# [Sta06]	MOF [MOF06]	XPath [BBC ⁺ 07]
synopsis	—	~	intent	synopsis	~	~	~	—
description	~	—	motivation	definition	usage	~	—	~
syntax	— ^a	syntax	structure	~	~	~	—	[NN] ^b
constraints	—	constraints	applicability	errors	restrictions	~	constraints	~
references	—	—	related patterns	—	references	~	—	~
relationship	—	—	consequences	return value, refinement	related	return type	—	~
semantics	—	semantics	collaborations	—	important	~	semantics	~
rationale	~	notes	implementation	rationale	GLS, ES ^c	note	rationale	note
example	examples	examples	sample code, known uses	—	~	example	—	~
update	—	—	—	—	—	— ^d	changes	—
default	—	—	—	—	note	default values	—	—
value	—	—	also known as	conforms to	—	—	—	—
list	~	—	—	messages, parameters	<i>terminals</i>	—	properties	~
section	~	—	—	—	~	~	—	~
subtopic	—	types	participants	—	fields	parameters, methods	operations	functions
Coverage of LDF								

Grammar export



Generated Rascal

```
@contributor{BGF2Rascal exporter - SLPS - http://github.com/grammarware/slps/wiki/BGF2Rascal}  
module CSharp
```

```
import ParseTree;  
import util::IDE;  
import IO;
```

```
layout WS = [\t-\n\r\ ]* !>> [\t-\n\r\ ];
```

```
start syntax CompilationUnit = UsingDirectives? GlobalAttributes? NamespaceMemberDeclarations?;
```

```
syntax UsingDirectives  
  = UsingDirective  
  | (UsingDirectives UsingDirective);
```

```
syntax UsingDirective  
  = UsingAliasDirective  
  | UsingNamespaceDirective;
```

```
syntax UsingAliasDirective  
  = "using" Identifier "=" NamespaceOrTypeName ";;";
```

...

Generated SDF

module Main

exports

context-free start-symbols Compilation-unit

sorts

...

context-free syntax

Using-directives? Global-attributes? Namespace-member-declarations? →
Compilation-unit

Using-directive → Using-directives

Using-directives Using-directive → Using-directives

Using-alias-directive → Using-directive

Using-namespace-directive → Using-directive

"using" Identifier "=" Namespace-or-type-name ";" → Using-alias-directive

...

Browsable grammars

Browsable C# 1.x Grammar



Extracted and/or recovered by [Vadim Zaytsev](#), see [Grammar Zoo](#) for details.

Source used for this grammar: *ECMA-334*, December 2001, Appendix A, pages 339–364

Summary

- Number of production rules: **240**
- Number of top alternatives: **584**
- Number of defined nonterminal symbols: **240**
- Root nonterminal symbols: —
- Other top nonterminal symbols: **1**: [compilation-unit](#)
- Bottom nonterminal symbols: **2**: [identifier](#), [literal](#)
- Number of used terminal symbols: **126**
- Special terminal symbols: **44**: `(".", "[", "]", ",", "(", ")", "++", "--", "+", "-", "!", "~", "*", "/", "%", "<<", ">>", "<", ">", "<=", ">=", "=="`, `"!="`, `"&", "^", "|", "&&", "||", "?", ":", "=", "+=", "-=", "*=", "/=", "%=", "&=", "|=", "^=", "<<=", ">>=", "{", "}", ";"`
- Keywords: **82**: `"bool", "decimal", "sbyte", "byte", "short", "ushort", "int", "uint", "long", "ulong", "char", "float", "double", "object", "string", "ref", "out", "this", "base", "new", "typeof", "void", "checked", "unchecked", "is", "as", "const", "if", "else", "switch", "case", "default", "while", "do", "for", "foreach", "in", "break", "continue", "goto", "return", "throw", "try", "catch", "finally", "lock", "using", "namespace", "class", "public", "protected", "internal", "private", "abstract", "sealed", "static", "readonly", "volatile", "virtual", "override", "extern", "params", "get", "set", "event", "add", "remove", "operator", "true", "false", "implicit", "explicit", "struct", "interface", "enum", "delegate", "assembly", "field", "method", "module", "param", "property"`

Syntax

```
namespace-name:  
  namespace-or-type-name
```

```
type-name:  
  namespace-or-type-name
```

<http://slps.github.io/zoo/cs/csharp-ecma-334-1.html>

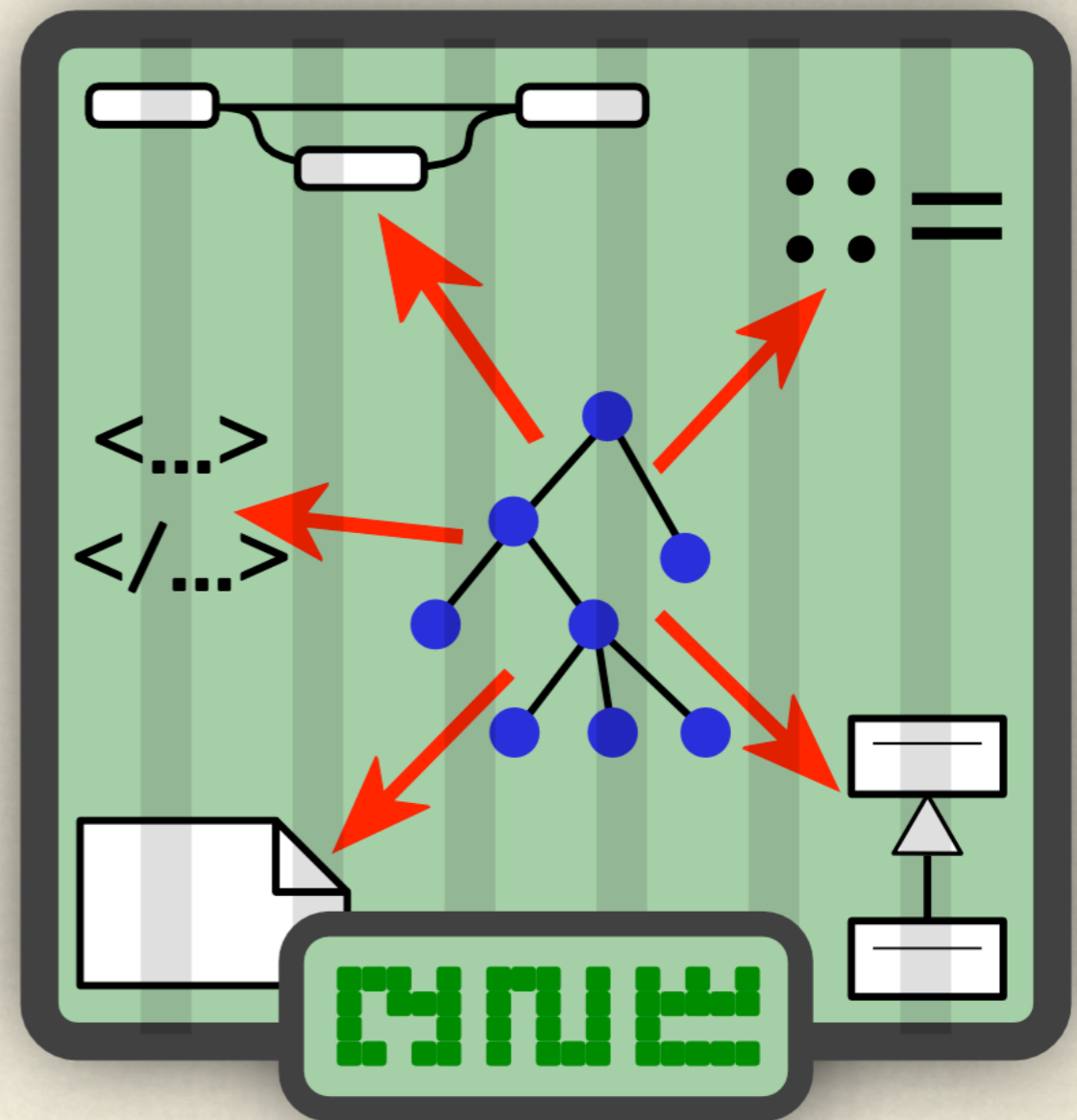
Grammar analysis

Grammar-based metrics

- ◆ PROD, VAR, TERM, ...
- ◆ VOC, LEN, EFF, PLEV, LLEV, ...
- ◆ LEV, CLEV, RLEV, NLEV, ...
- ◆ cf.
- ◆ J. F. Power, B. A. Malloy. *A Metrics Suite for Grammar-Based Software*. 2004.
- ◆ T. Alves, J. Visser. *Metrication of SDF Grammars*. 2005.
- ◆ J. Cervelle, M. Črepinšek, R. Forax, T. Kosar, M. Mernik, G. Roussel. *On Defining Quality Based Grammar Metrics*. 2009.
- ◆ M. Črepinšek, T. Kosar, M. Mernik, J. Cervelle, R. Forax, G. Roussel. *On Automata and Language Based Grammar Metrics*. 2010.

Micropatterns

- ◆ Recognisable
- ◆ Purposeful
- ◆ Prevalent
- ◆ Simple
- ◆ Scoped
- ◆ Observed



To summarise



- ◆ Extraction
- ◆ recovery
- ◆ Evolution
- ◆ transformation
- ◆ negotiated
- ◆ pending
- ◆ Documentation
- ◆ Export





GRAMMARLAB

Questions?

(What else would you like?)

`vadim@grammarware.net`

