



SWAT

Bidirectional Grammar Transformations

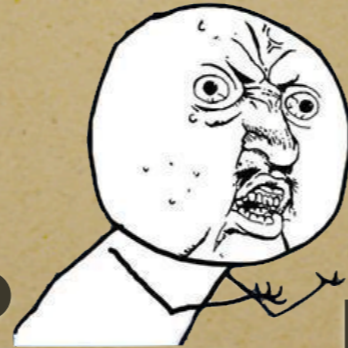
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In a nutshell

- Grammars in the broad sense
 - protocols, metamodels, schemata, ontologies, structures
- Grammars evolve
 - language evolution, error fixes, dialects, etc
- Systematic way
 - programmable transformation steps
 - operator suite
 - bidirectionality



Motivation

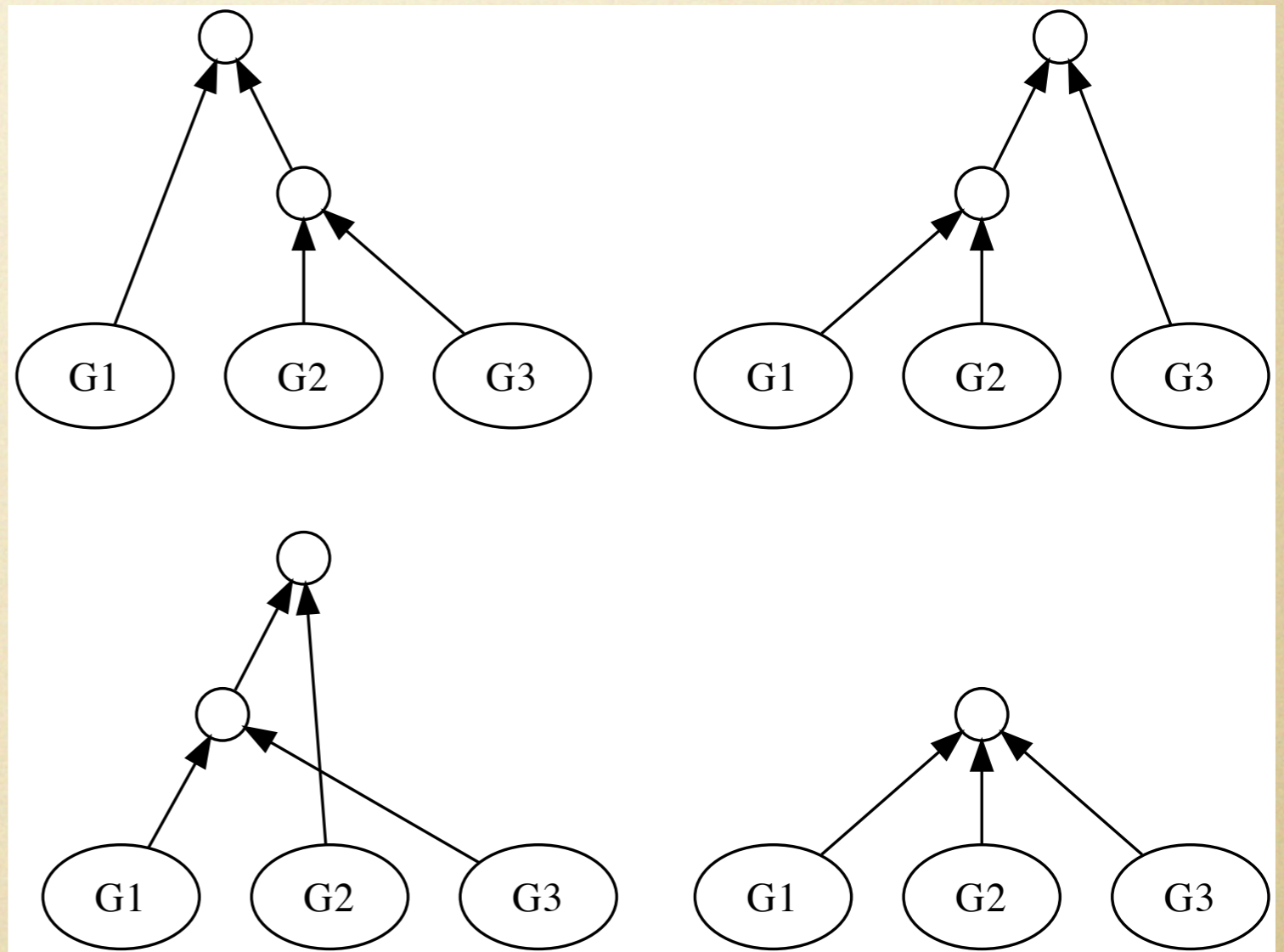
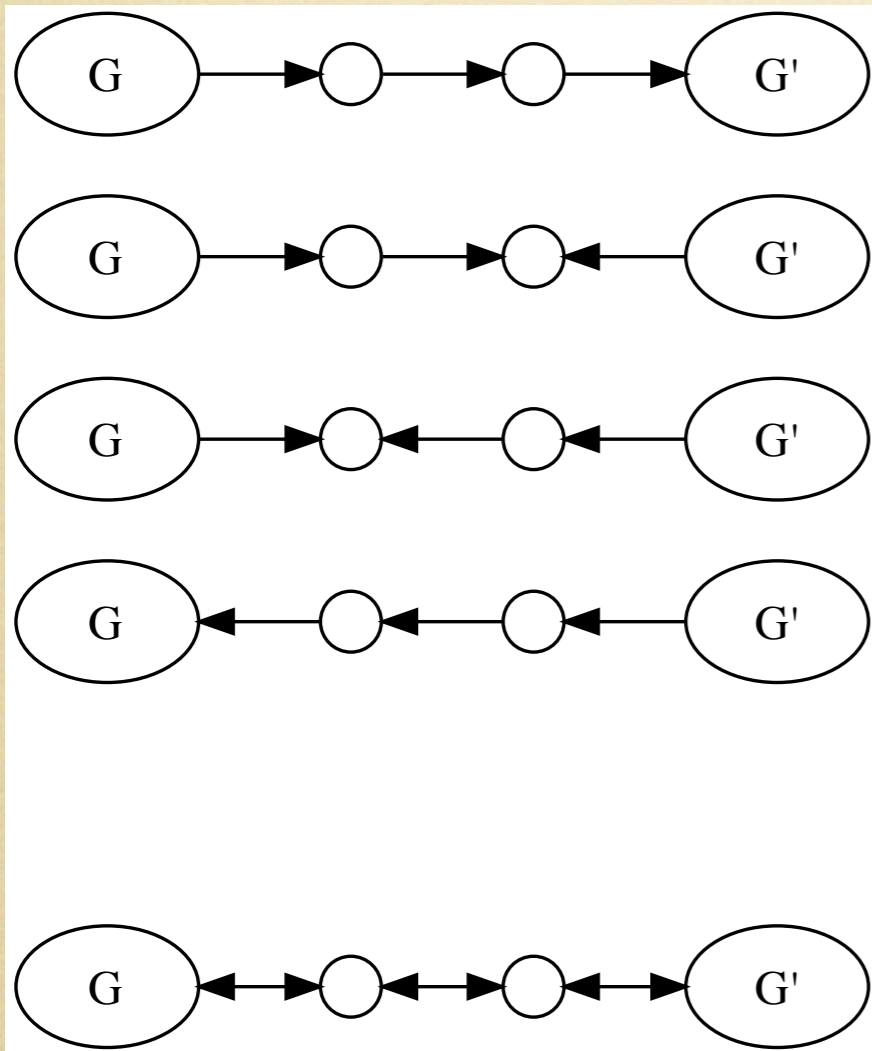
Why transformation?

- Edit as text
 - no (technical) problems during editing
 - some problems afterwards
 - no automation
 - no replication
- Program transformations
 - somewhat more complex technically
 - more maintainable, reproducible, transformable, ...

Why bidirectional?

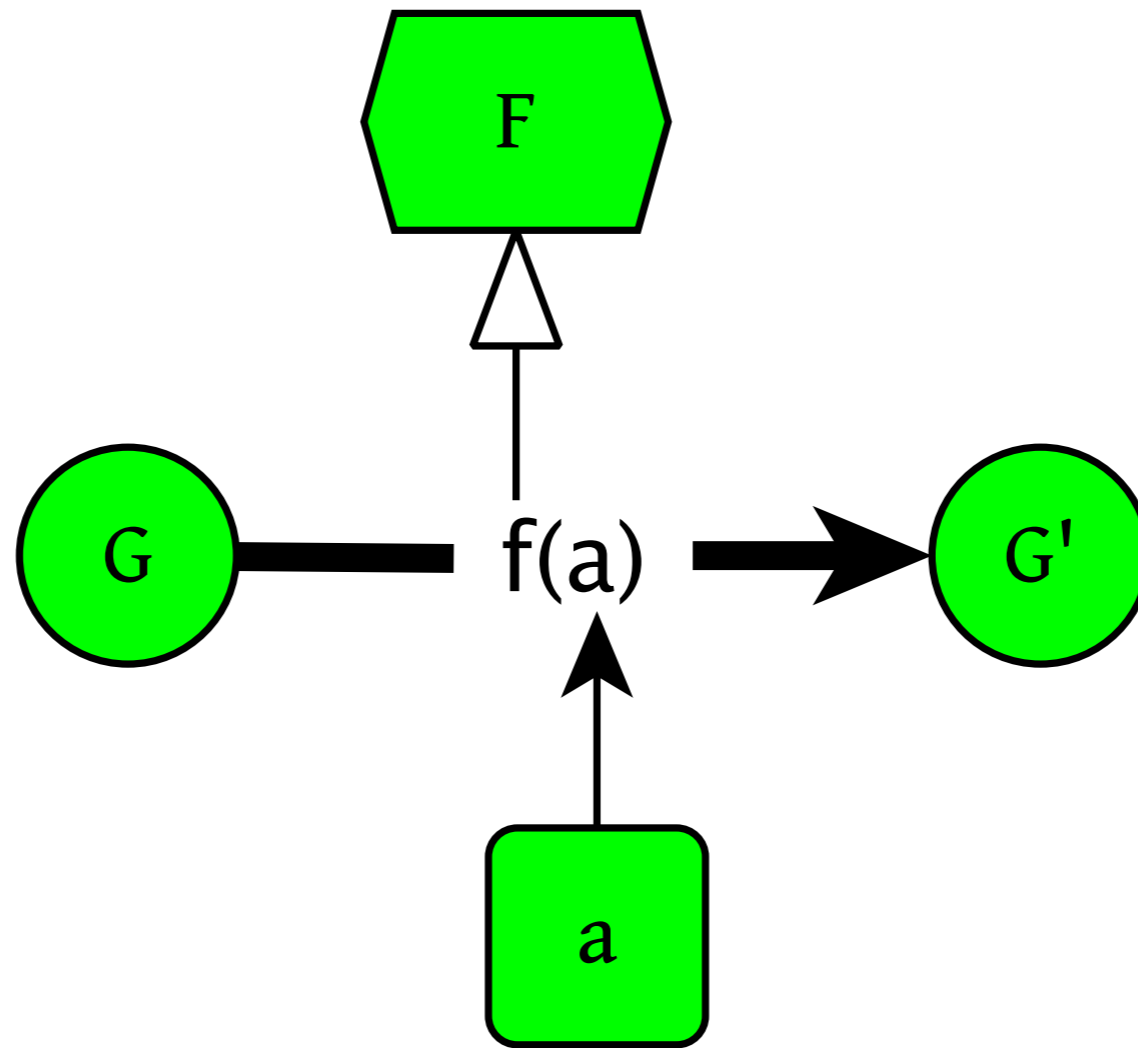
- Sweet spot between
 - functional approach
 - $y = f(x)$
 - declarative approach
 - $f(x,y) :- \dots$
- Fixed application order, direction is flexible

Bidirectionality profits

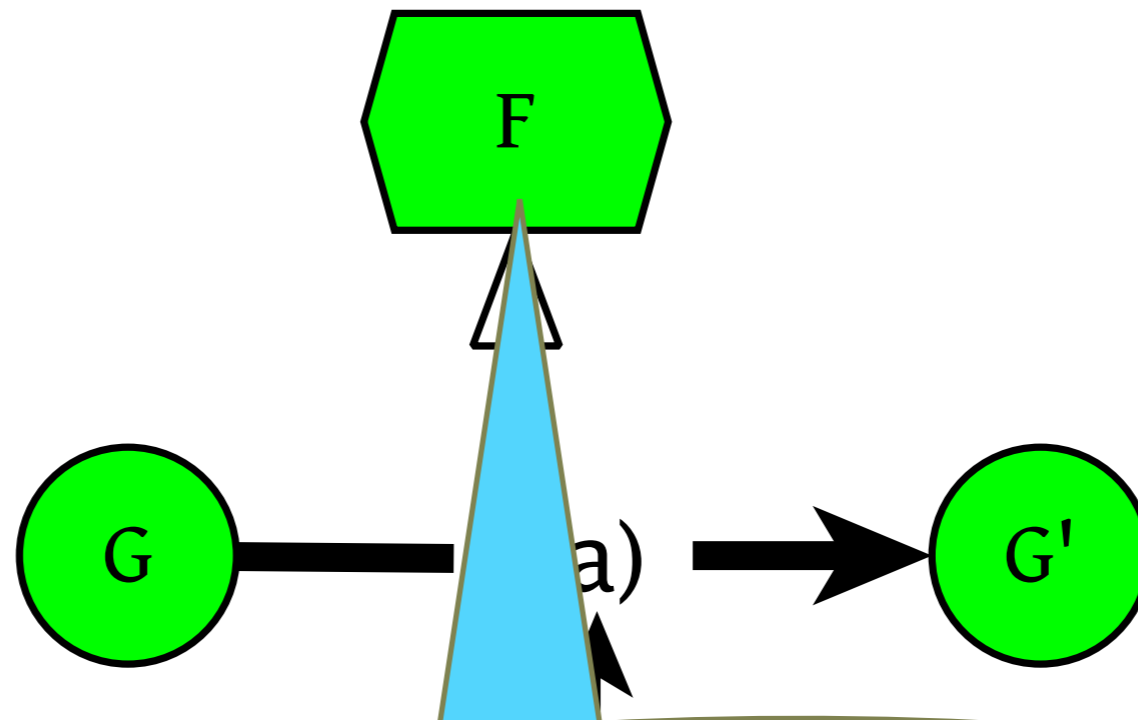


Unidirectional Transformations

Transformation components



Transformation components



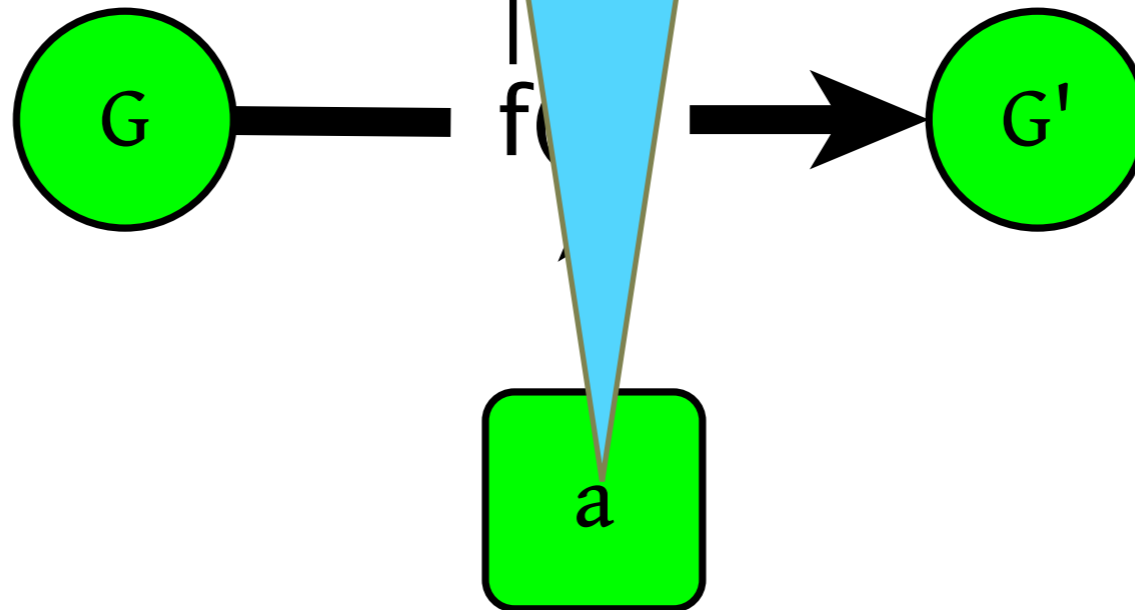
Operator

- known semantics, well-defined algorithm
 - rename, fold, factor, inject, remove, ...

Transformation components

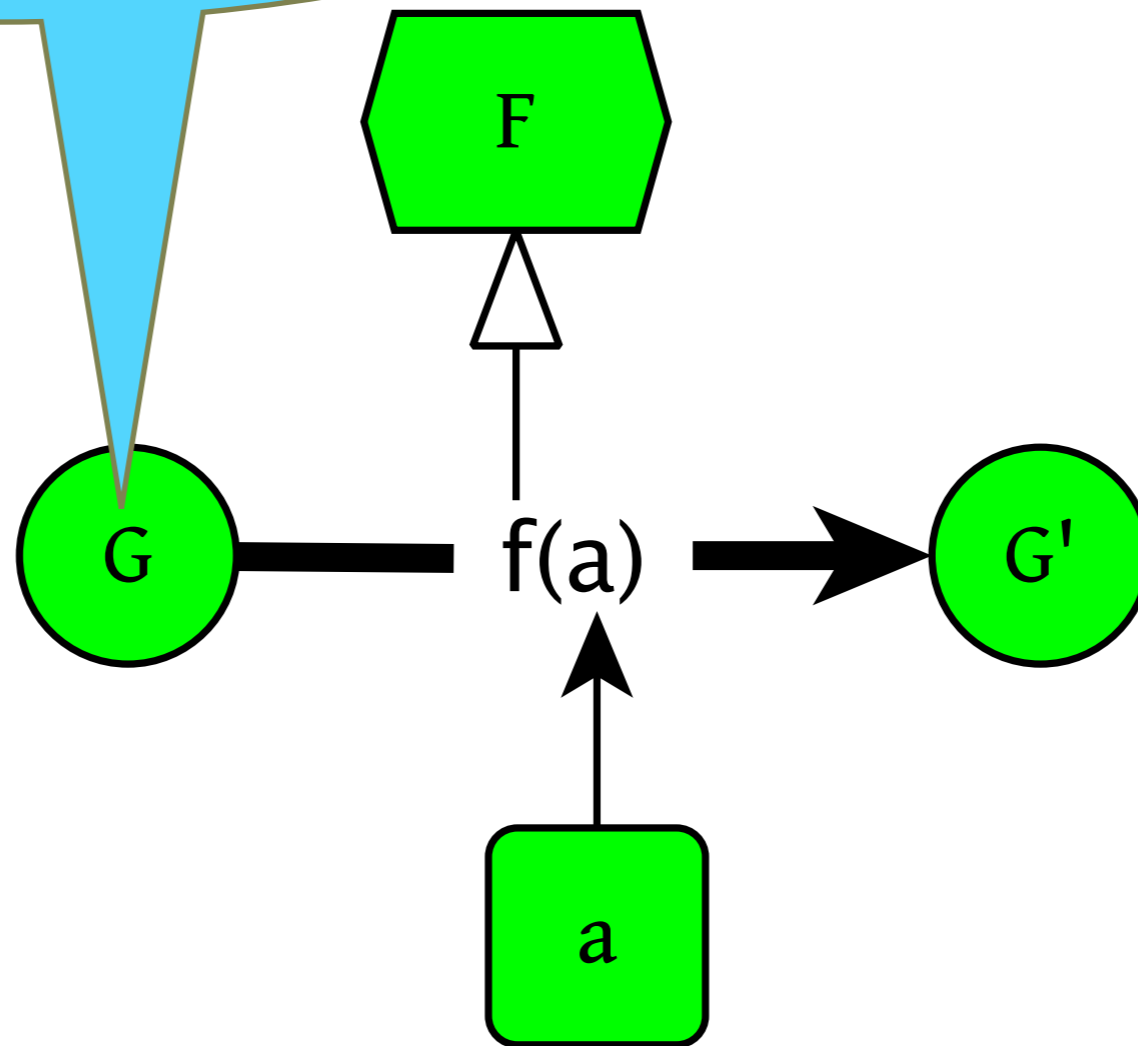
Arguments

- what exactly to rename/factor/inject/...?

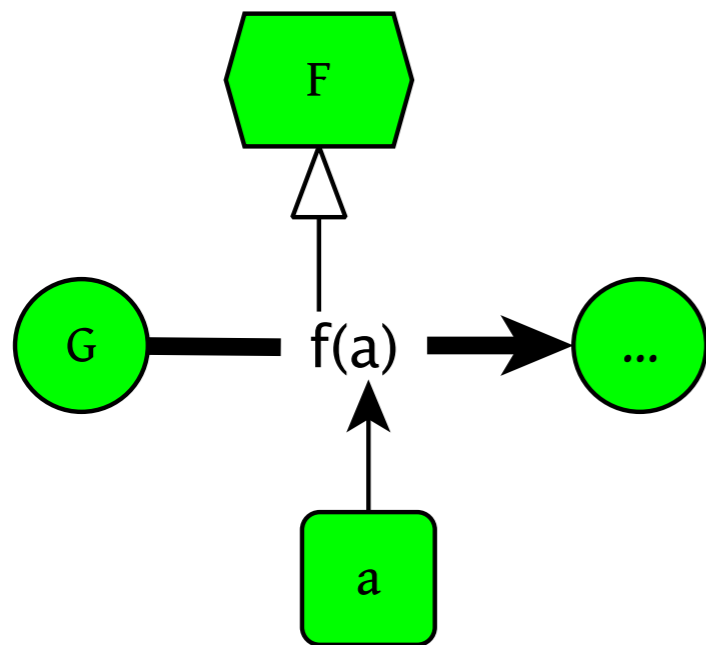


Input grammar

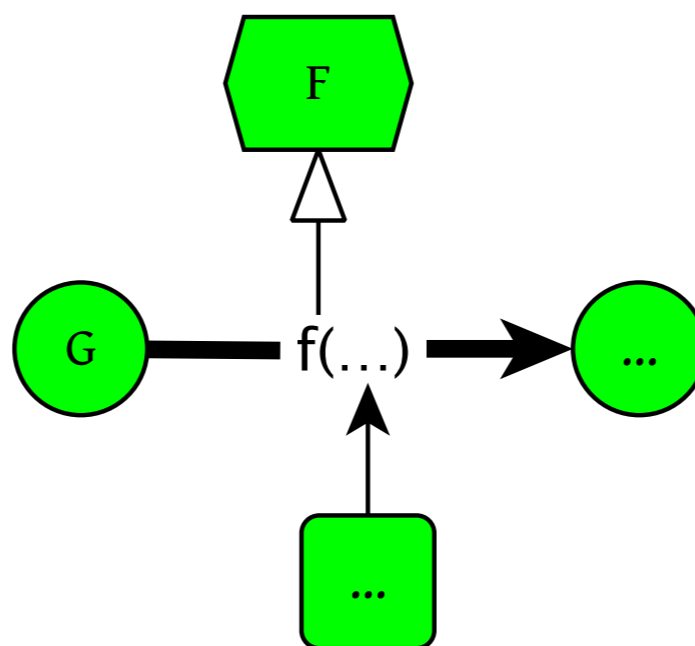
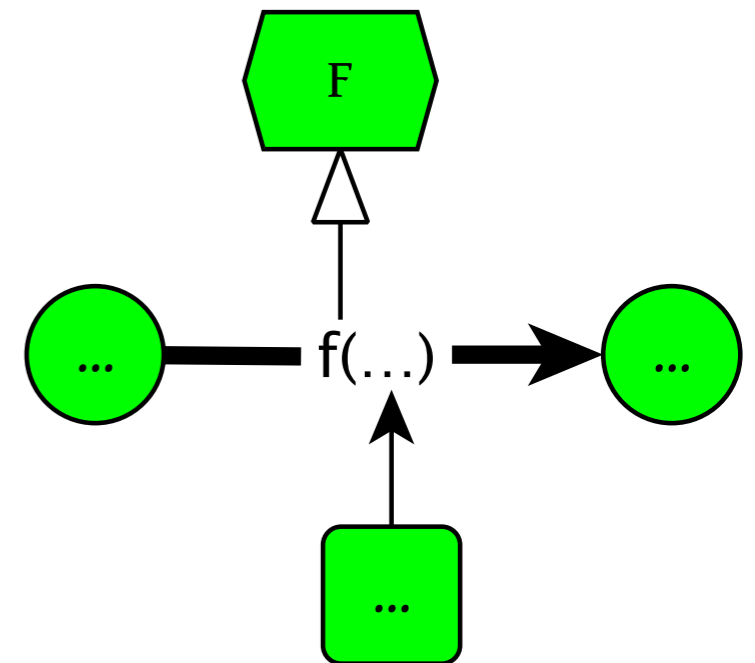
- determines applicability



Variations on components



*other
variations
possible*



Example

```
expr : ...;  
atom : ID | INT | '(' expr ')';
```

A fragment of concrete syntax.

What if we want to derive the abstract syntax?

Example

```
expr : ...;  
atom : ID | INT | '(' expr '');
```

Need to merge “expr” & “atom”

Need to project away ‘(’ & ‘)’

Alternative needs to go entirely

A fragment of concrete syntax.

What if we want to derive the abstract syntax?

A transformation sequence

expr : ...;
atom : ID | INT | '(' expr ')';

abstractize

expr : ...;
atom : ID | INT | expr;

vertical

expr : ...;
atom : ID;
atom : INT;
atom : expr;

unite

expr : ...;
expr : ID;
expr : INT;

abridge

expr : ...;
expr : ID;
expr : INT;
expr : expr;

XBGF Operator Suite

$$L(G_1) = L(G_2)$$

- Semantics-preserving (refactoring)
 - rename, import, introduce, eliminate
 - fold, unfold, extract, inline
 - factor, distribute, horizontal, vertical
 - yaccify, deyaccify, massage
 - designate, unlabel
 - abridge, detour
 - ...

XBGF Operator Suite

- Semantics-in/decreasing
 - appear, disappear
 - narrow, widen
 - add, remove
 - upgrade, downgrade
 - unite
 - ...

$$L(G_1) \subseteq L(G_2)$$

$$\vee$$

$$L(G_2) \subseteq L(G_1)$$

XBGF Operator Suite

- Semantics-revising
 - undefine, define
 - inject, project, permute
 - abstractize, concretize
 - replace, redefine

$$L(G_1) \not\subseteq L(G_2)$$

$$\wedge$$

$$L(G_2) \not\subseteq L(G_1)$$

Bidirectional Transformations

≡BGF Operator Suite

- Pairs already formed
 - chain — unchain
 - add — remove
 - fold — unfold
 - ...
- Implicit pairs
 - factor — factor
 - message — message
 - replace — replace
 - ...

EBGF Operator Suite

- Incomplete pairs
 - rassoc — ???
 - unite — ???
 - equate — ???
 - ...
- Asymmetrical pairs
 - designate(prod) — unlabel(label)
 - deyaccify(nt) — yaccify(prods)
 - eliminate(nt) — introduce(prods)
 - ...

Ξ BGF Operator Suite

- Remaining problems
 - Lack of symmetry in execution
 - $\text{unfold}(A:\varepsilon) \text{ — fold}???$
 - $\text{unfold}(A:b \text{ in } bAb) \text{ — fold}???$
 - ...
- Hence, bidirectionality and not bijection

To summarise

- Grammar transformation steps describe software language evolution
- Programmable grammar transformation operators
- XBGF: unidirectional
- EBGF: bidirectional
- SLPS: implementations in Prolog and Rascal
- <http://grammarware.github.com>



Questions?

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