



SWAT

Language Evolution, Metasyntactically

First International Workshop on Bidirectional Transformations (BX 2012)

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Introduction

- Every language document employs its own notation
- We focus on **metalanguage evolution**
 - the language itself does not evolve
 - the notation in which it is written, does
- We limit ourselves to grammarware technical space
- Working prototypes are a part of SLPS

Metalinguage evolution

- Notation correction
 - Misused notation
 - Overused notation
- Notation evolution
 - Notations are languages, they evolve
- Mapping between notations
 - Many notations are equivalent
 - Most are (almost EBNF) + (little extra)

Metasyntactic
evolution
megamodel

Megamodel

Syntactic
notation
specification

EBNF Dialect Definition

`program ::=`

`function+;`

`function ::=`

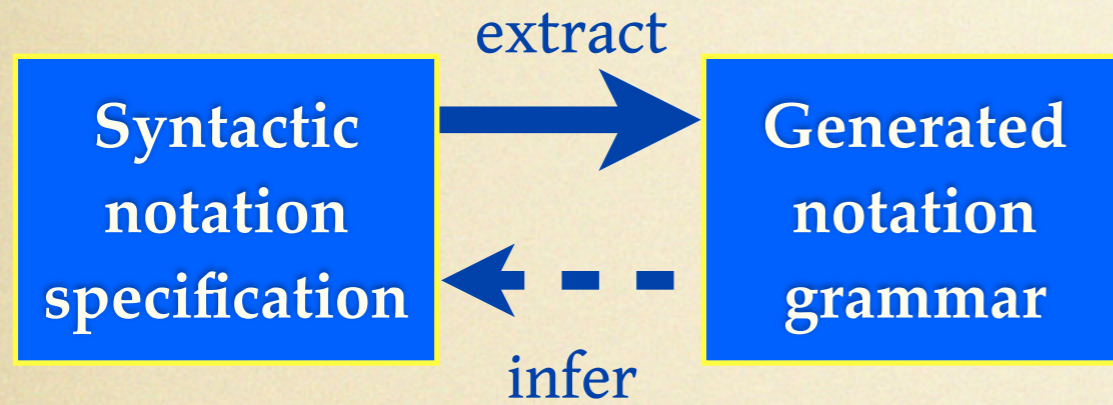
`name argument* "=" expr?;`

- List of indicators
- Together form a notation specification

EDD example

defining metasymbol	:	definition separator metasymbol	
terminator metasymbol	;	postfix optionality metasymbol	?
postfix star metasymbol	*	postfix plus metasymbol	+
start terminal metasymbol	"	end terminal metasymbol	"

Megamodel



Generated notation grammar (in Rascal)

```
module LLL
import util::IDE; // needed only for advanced IDE support (see last two lines)
start syntax LLLGrammar = LLLLayoutList LLLProduction * LLLLayoutList;
syntax LLLProduction = LLLNonterminal ":" {LLLDefinition "|" }+ ";";
syntax LLLDefinition = LLLSymbol+;
syntax LLLSymbol
= @category="Identifier" nonterminal: LLLNonterminal
| @category="Constant" terminal: LLLTerminal
| group: "(" LLLDefinition ")"
| optional: LLLSymbol "?"
| star: LLLSymbol "*"
| plus: LLLSymbol "+"
| sepliststar: "{" LLLSymbol LLLSymbol "}" "*"
| seplistplus: "{" LLLSymbol LLLSymbol "}" "+";
lexical LLLTerminal = "\" LLLTerminalSymbol* "\"";
lexical LLLTerminalSymbol = ![""];
lexical LLLNonterminal = [A-Za-z_01-9\-\-]+ !>> [A-Za-z_01-9\-\-];
layout LLLLayoutList = LLLLayout* !>> [\t-\n \r \ ] !>> "#";
lexical LLLLayout = [\t-\n \r \ ] | LLLComment ;
lexical LLLComment = @category="Comment" "#" ![\n]* [\n];
Tree getLLL(str s,loc z) = parse(#LLLGrammar,z);
public void registerLLL() = registerLanguage("LLL","lll",getLLL);
```

Grammar internal representation

$grammar(Rs, Ps) \Leftarrow mapoptlist(n, Rs), maplist(prod, Ps).$
 $prod(p(L, N, X)) \Leftarrow mapopt(label, L), atom(N), expr(X).$
 $label(l(X)) \Leftarrow atom(X).$
 $expr(true).$
 $expr(fail).$
 $expr(a).$
 $expr(t(T)) \Leftarrow atom(T).$
 $expr(n(N)) \Leftarrow atom(N).$
 $expr(', '(Xs)) \Leftarrow maplist(expr, Xs).$
 $expr('; '(Xs)) \Leftarrow maplist(expr, Xs).$
 $expr('?'(X)) \Leftarrow expr(X).$
 $expr('*'(X)) \Leftarrow expr(X).$
 $expr('+'(X)) \Leftarrow expr(X).$
 $expr(slp(X, Y)) \Leftarrow expr(X), expr(Y).$
 $expr(sls(X, Y)) \Leftarrow expr(X), expr(Y).$
 $expr(s(S, X)) \Leftarrow atom(S), expr(X).$

grammar = start symbols + productions
production = label + lhs + rhs
production labels
 ϵ
empty language
universal type
terminal symbols
nonterminal symbols
sequential composition
choice
optionality
Kleene star
transitive closure
 Y -separated list with 1 or more elements
 Y -separated list with 0 or more elements
selectable expressions

Generated notation grammar (in BGF)

LLL1 Grammar:

LLL1 Production *

LLL1 Production:

LLL1 Nonterminal ":" {LLL1 Definition "|" }+ ";"

LLL1 Definition:

LLL1 Symbol+

[nonterminal] LLL1 Symbol:

LLL1 Nonterminal

[terminal] LLL1 Symbol:

LLL1 Terminal

[optional] LLL1 Symbol:

LLL1 Symbol "?"

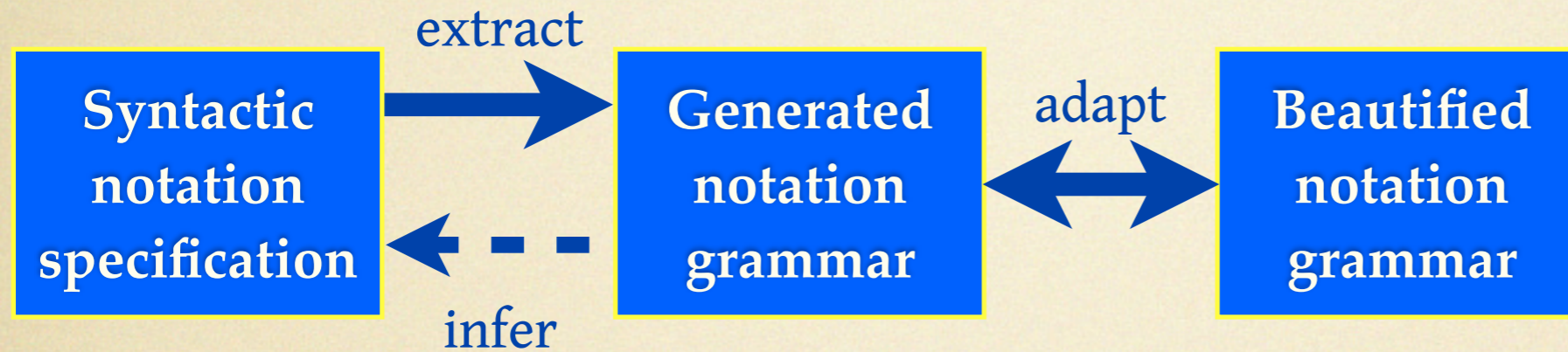
[star] LLL1 Symbol:

LLL1 Symbol "*"

[plus] LLL1 Symbol:

LLL1 Symbol "+"

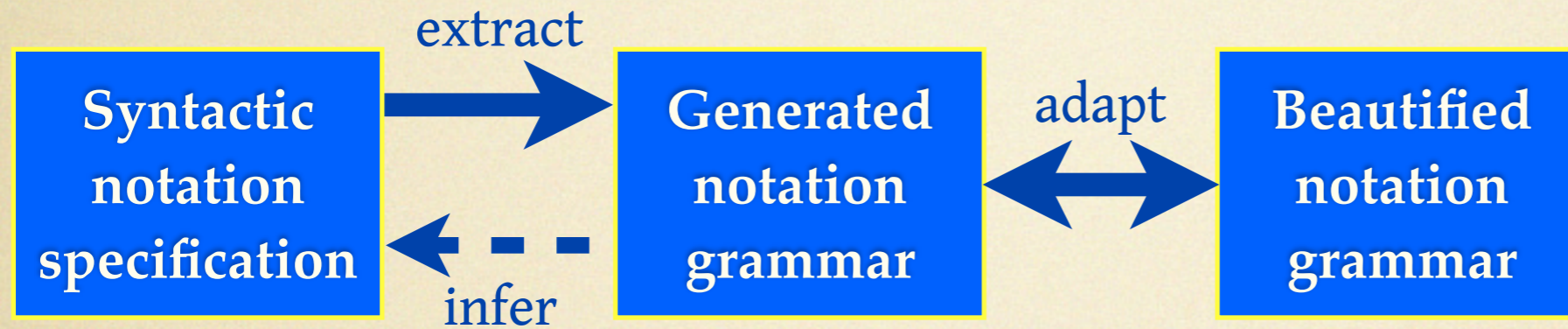
Megamodel



Beautified notation grammar (in BGF)

```
grammar:  
  rule+  
rule:  
  sort ":" alts ";"  
alts:  
  alt alts-tail*  
alts-tail:  
  "|" alt  
alt:  
  term*  
term:  
  basis repetition?  
basis:  
  literal  
  sort  
repetition:  
  "*" "  
  "+" "  
  "?" "
```

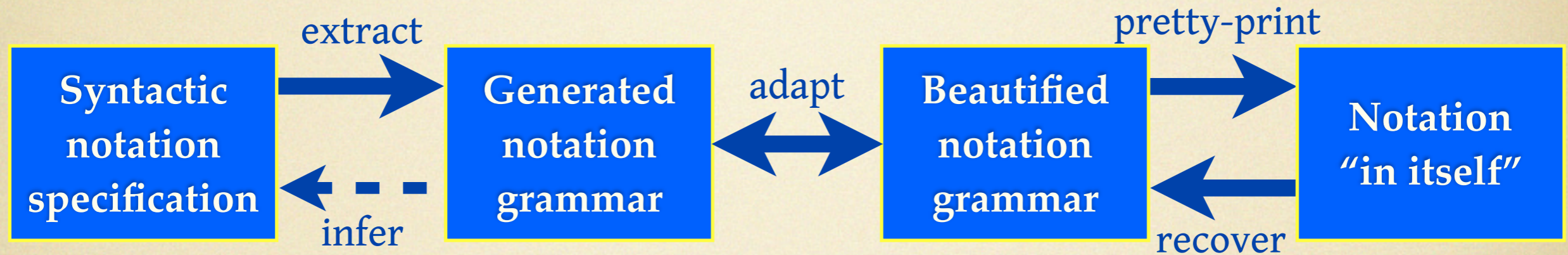
Megamodel



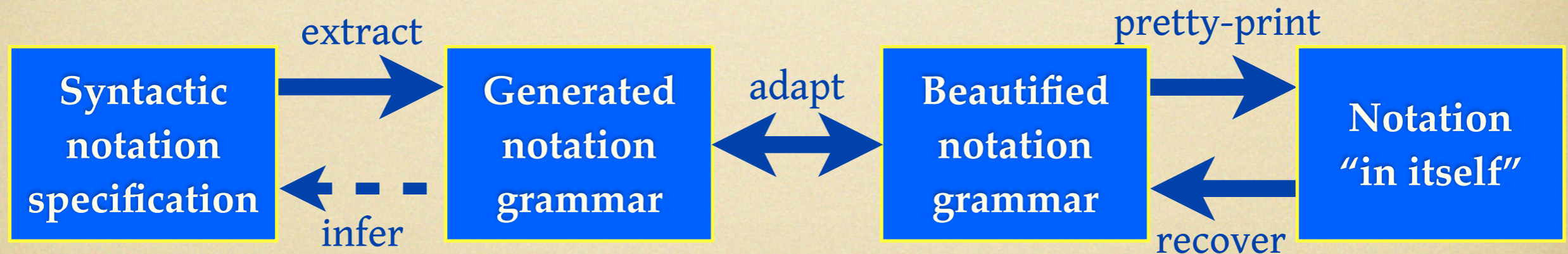
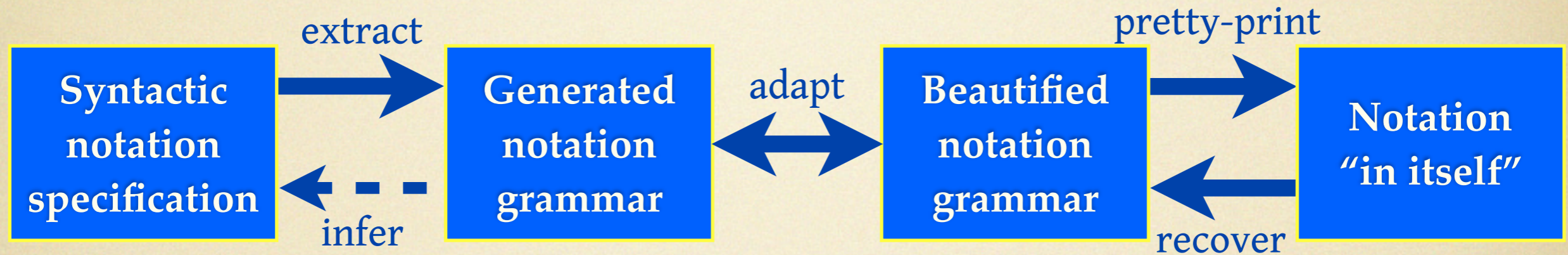
Bidirectional grammar adaptation

- XBGF \Rightarrow Ξ BGF:
 - renameN, factor, etc: flip arguments
 - addV/removeV, narrow/widen: form pairs
 - extract/inline, unlabel/designate: asymmetry
 - distribute: removed from the language
 - unite, equate: tricky, superposition of others
- BX is a stable way to represent grammar relationship

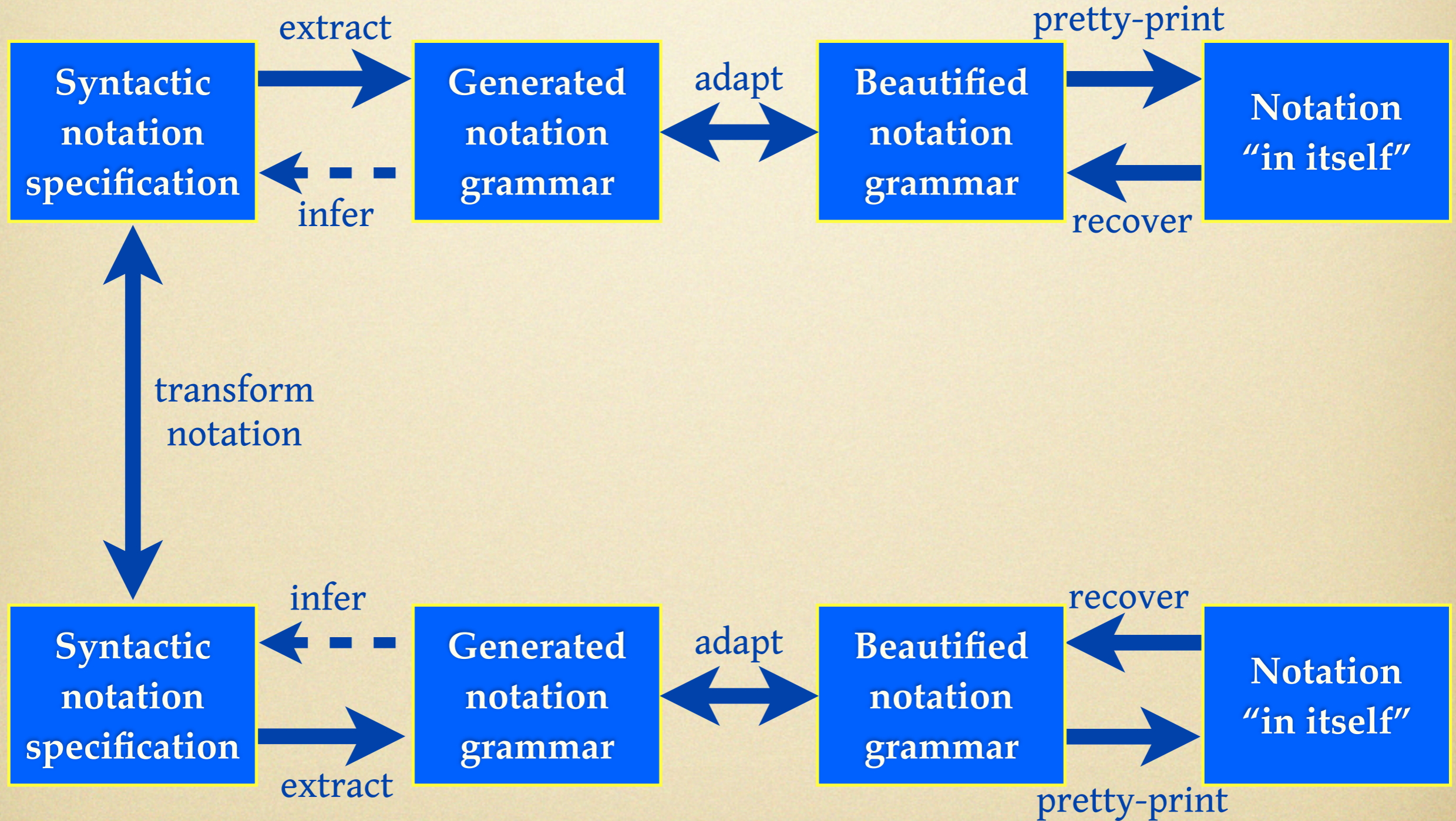
Megamodel



Megamodel



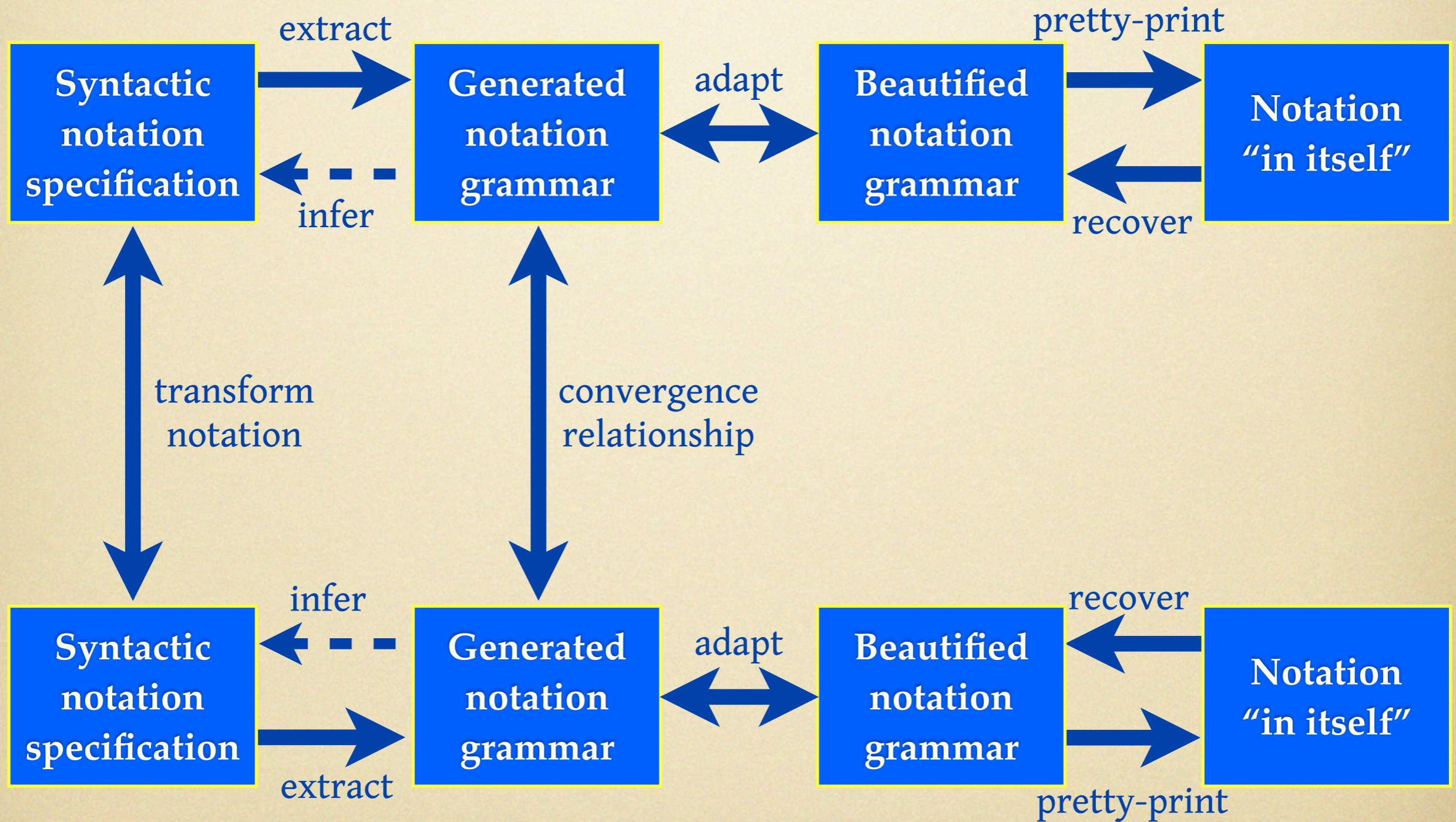
Megamodel



Notation transformation

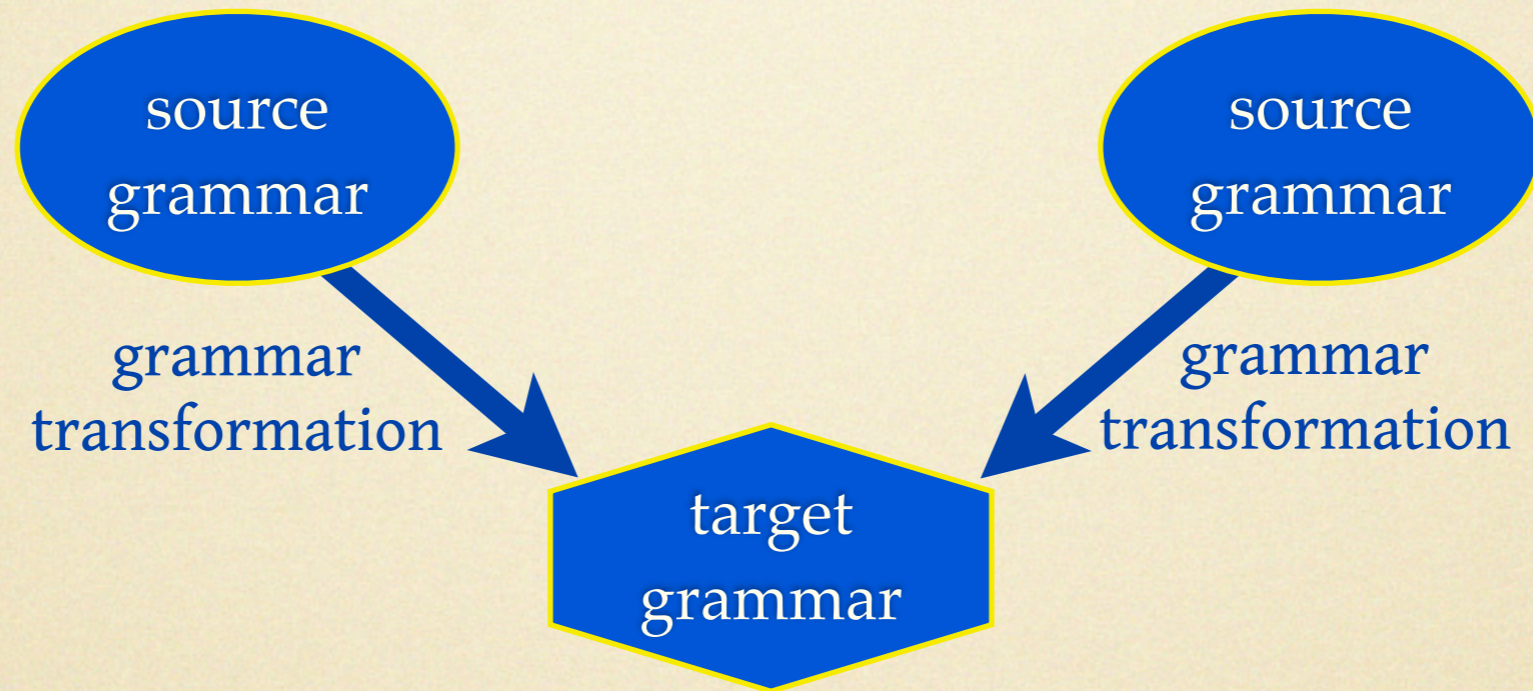
- EDD — notation, consists of metasymbols
- XEDD — transformation language
 - **rename-metasymbol**($s, v1, v2$)
 - e.g., change defining metasymbol from “:” to “::=”
 - **introduce-metasymbol**(s, v)
 - e.g., bring a terminator metasymbol to a notation
 - **eliminate-metasymbol**(s, v)

Megamodel

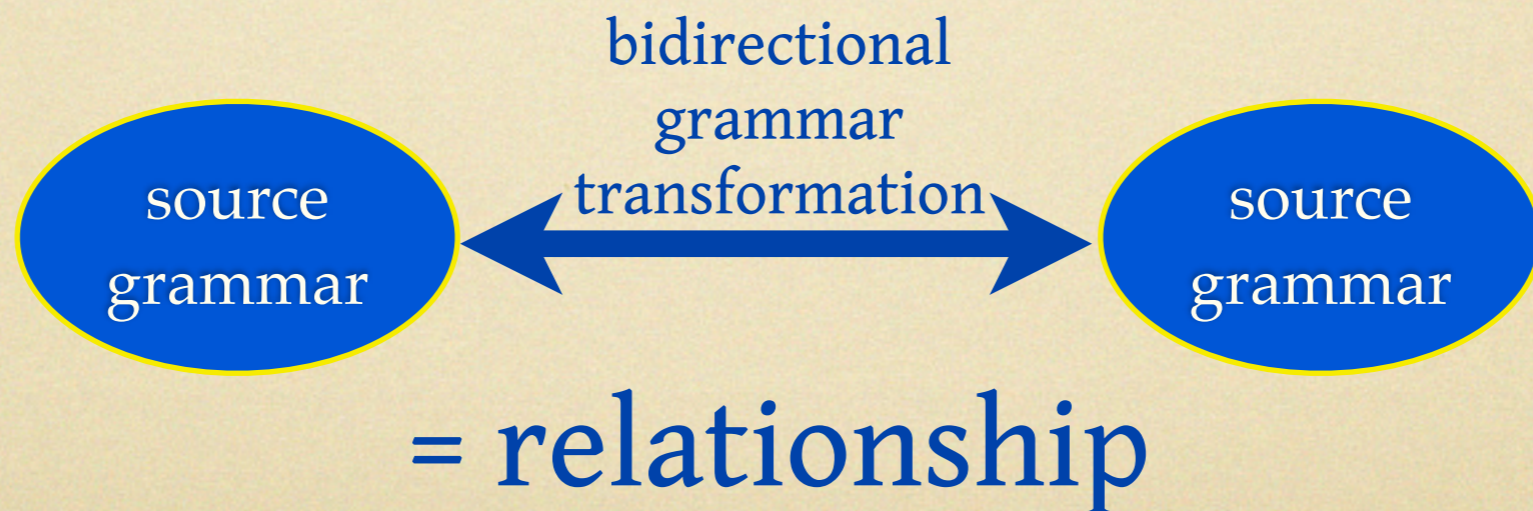


Grammar convergence

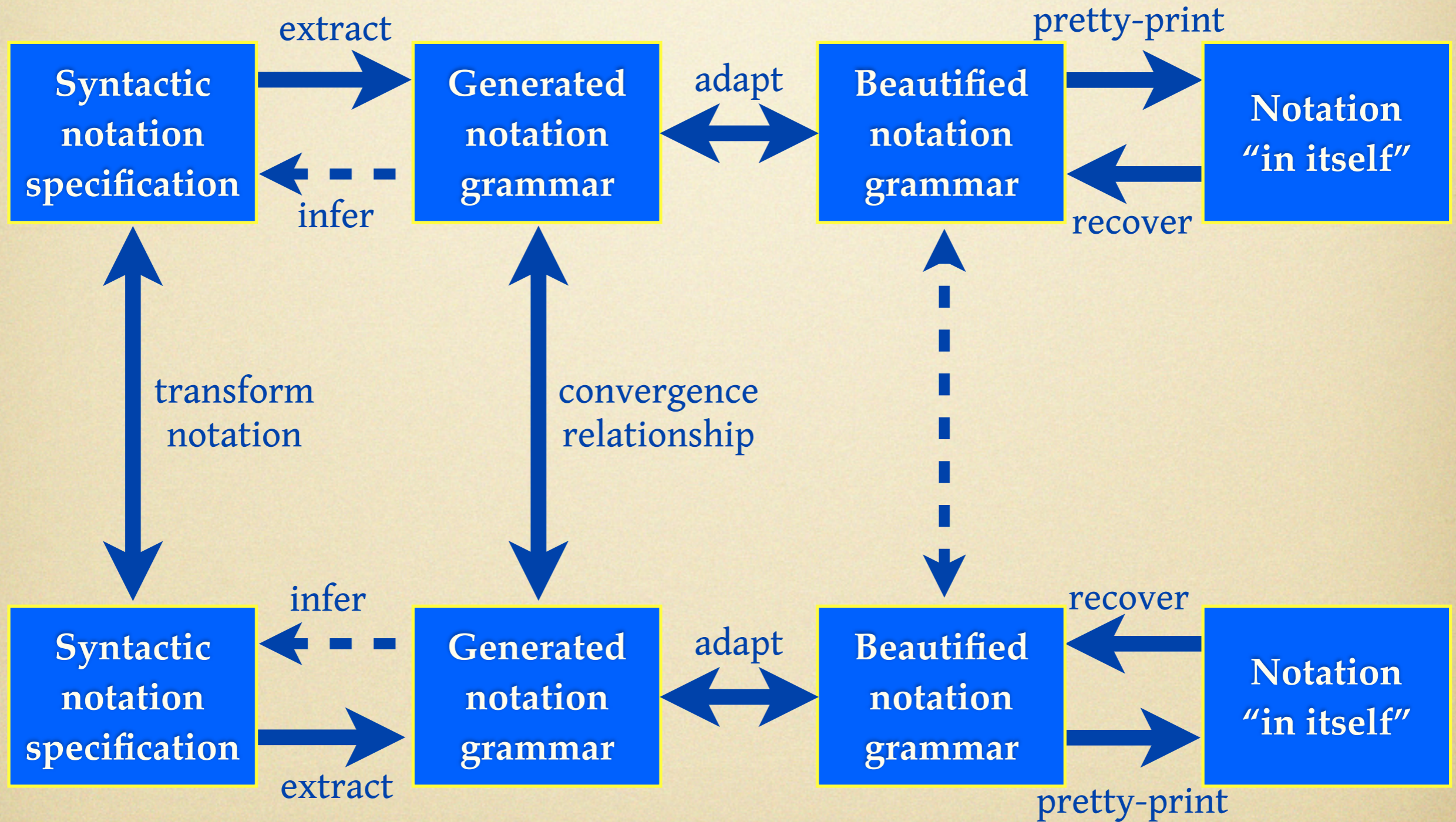
XBGF



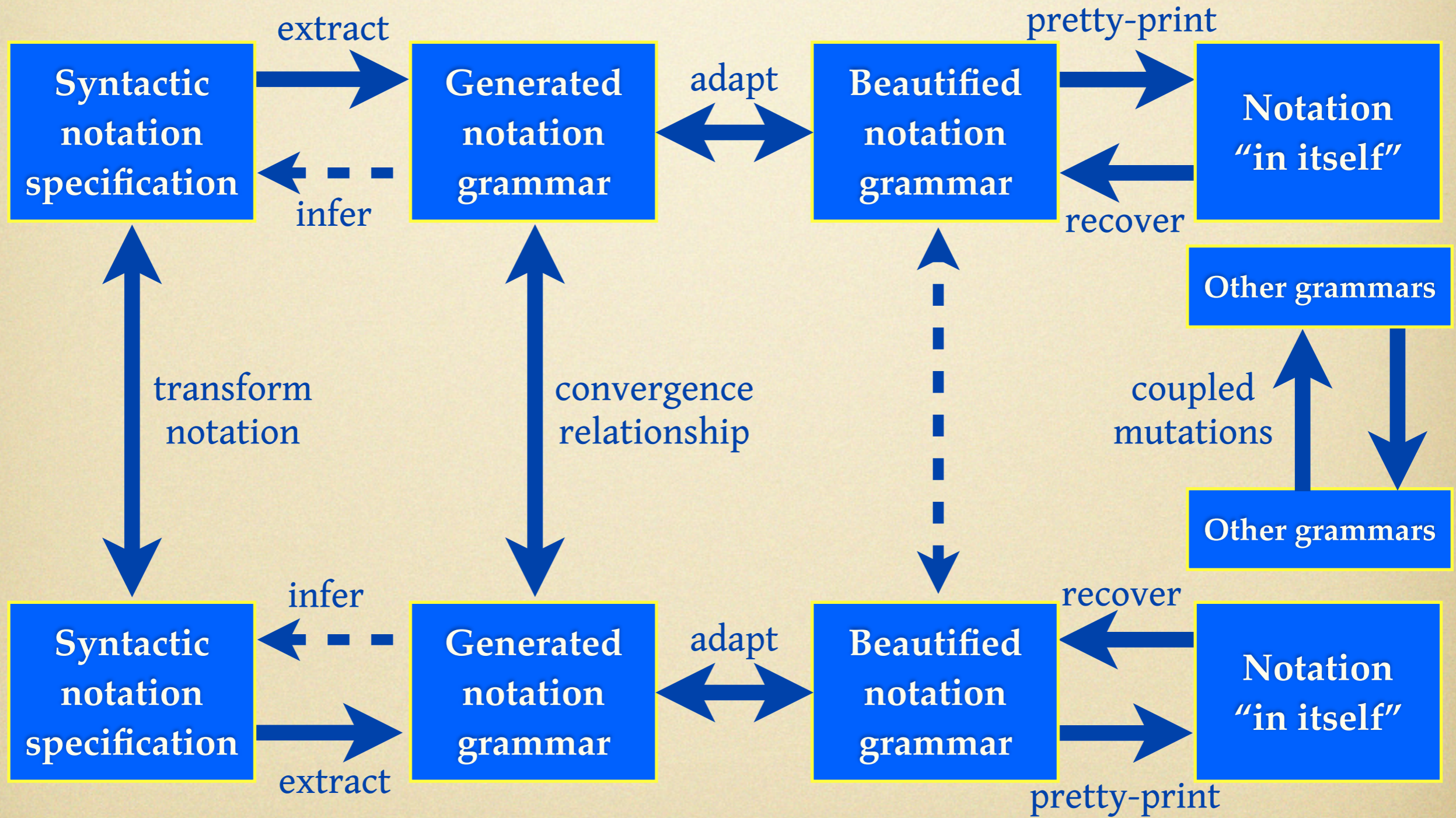
EBGF



Megamodel



Megamodel



Grammar transformation vs. grammar mutation

- A **grammar transformation operator** τ can be formalised as a triplet:
 $\tau = \langle c_pre, t, c_post \rangle$.
- A **grammar transformation** then is $\tau_a_i(G)$, resulting in G' .
 - if a_i are of incorrect types and quantity than expected by t
 $\Rightarrow \tau$ is **incorrectly called**;
 - if the constraint c_pre does not hold on G
 $\Rightarrow \tau_a_i$ is **inapplicable** to G ;
 - if the constraint c_post holds on G
 $\Rightarrow \tau_a_i$ is **vacuous** on G ;
 - if the constraint c_pre holds on G and c_post does not hold on G'
 $\Rightarrow t$ is **incorrectly implemented**;
 - if c_pre holds on G , c_post holds on G'
 $\Rightarrow \tau$ has been **applied correctly** with arguments a_i to G resulting in G' .

Grammar transformation VS. grammar mutation

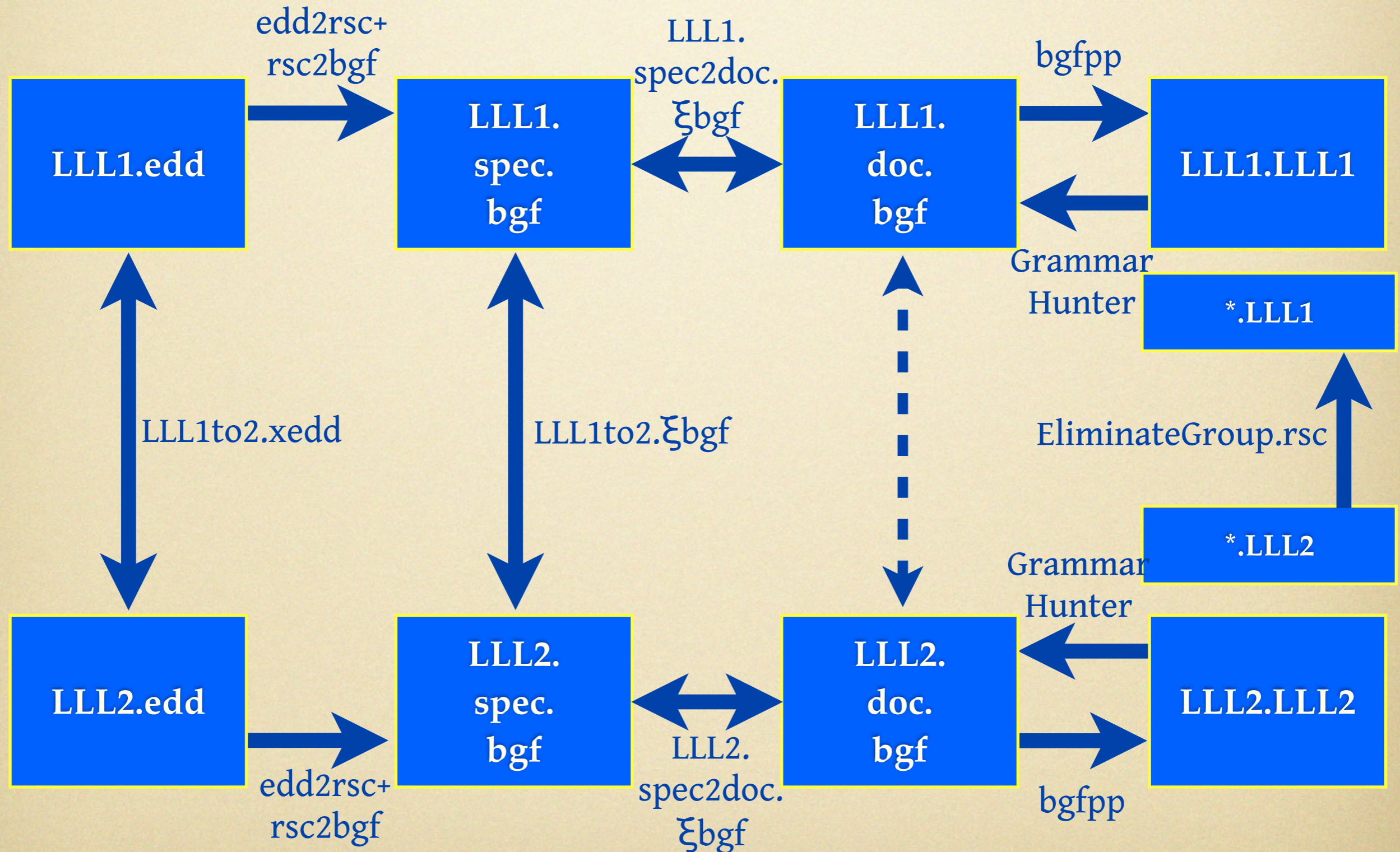
- A **grammar mutation** does not have a single precondition
- It has a set of preconditions that serve as **triggers**:
 $\mu = \langle \{c_i\}, \{t_i\}, c_post \rangle$.
- The mutation **terminates** once no trigger c_i holds and the postcondition c_post is met.
- A **bidirectional** grammar mutation:
 $\mu_bx = \langle c_pre, \{c_i\}, \{t_i\}, c_post \rangle$
will be an instantiation of a grammar mutation
- The family of spawned BMs does **not** define the original:
i.e., $\forall \mu \exists G \exists G' \nexists \mu_bx, G' = \mu(G) \wedge G' = \mu_bx(G) \wedge G = \mu^{-1}(G')$.

Notation evolution summary

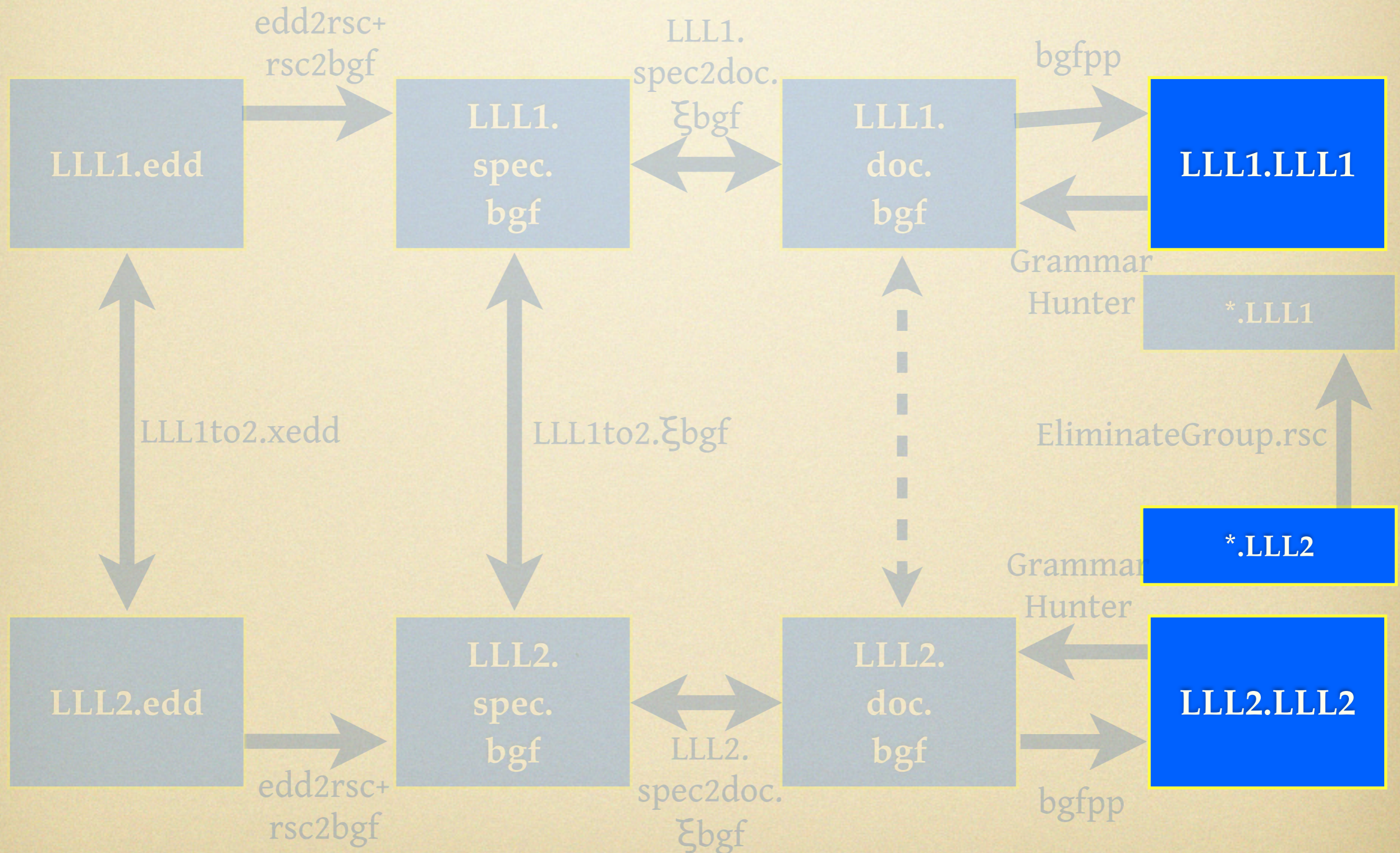
- A notation evolution step Δ consists of the following coupled components:
 - σ , a **bidirectional notation transformation** that changes the notation itself
 - δ , a **convergence relationship** that can transform the notation grammars
 - γ , a **bidirectional grammar adaptation** that prepares a beautified readable version of N'
 - μ , an unidirectional **coupled grammar mutation** that migrates the grammarbase according to notation changes
 - possibly μ' , an unidirectional **coupled grammar mutation** that migrates the grammarbase according to the inverse of the intended notation changes

Evaluation

Megamodel: case study



Megamodel: previously



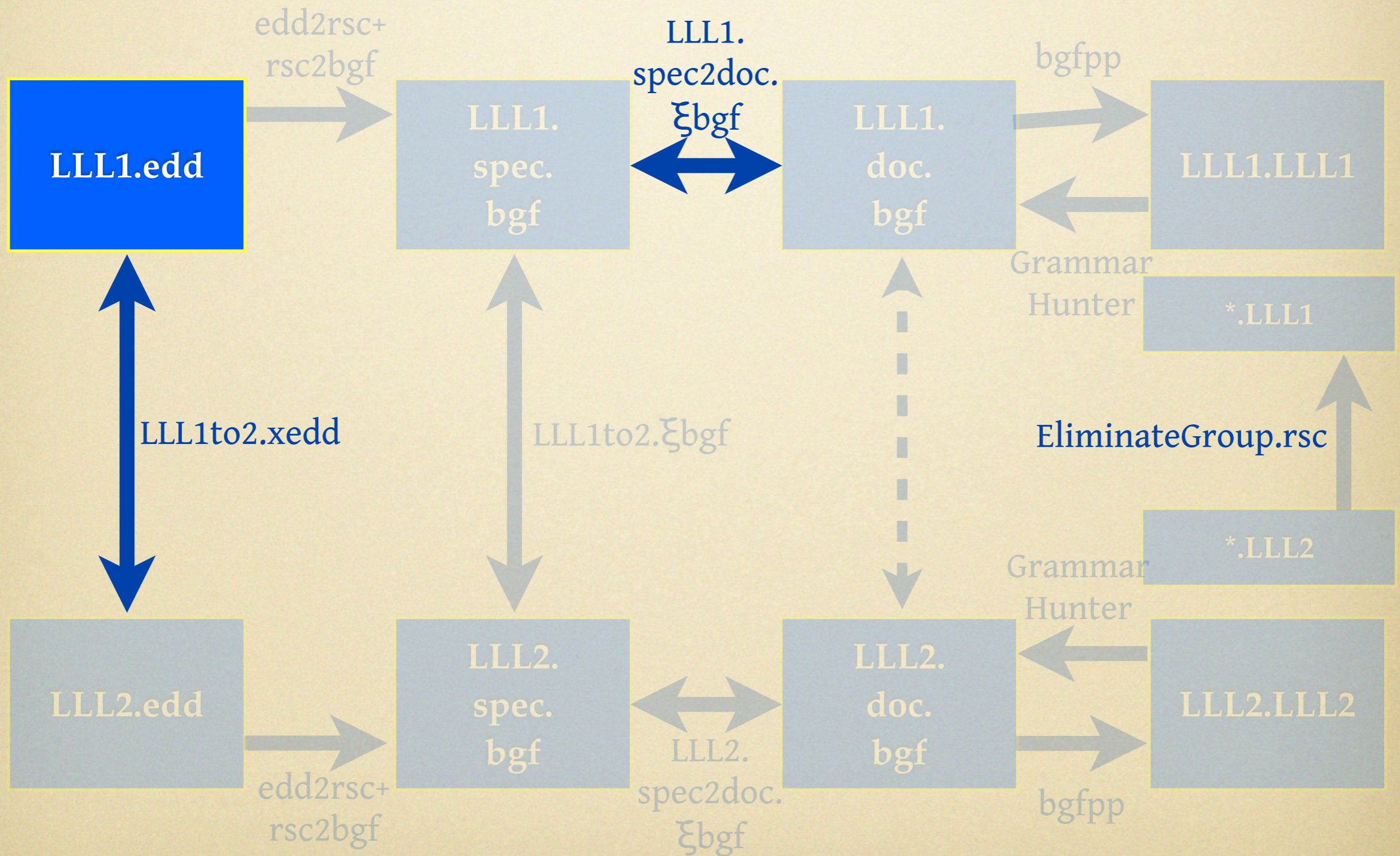
LLL1 in itself

```
grammar      : rule+;
rule         : sort ":" alts ";" ;
alts         : alt alts-tail*;
alts-tail    : "|" alt;
alt          : term*;
term         : basis repetition?;
basis        : literal | sort;
repetition   : "*" | "+" | "?" ;
```

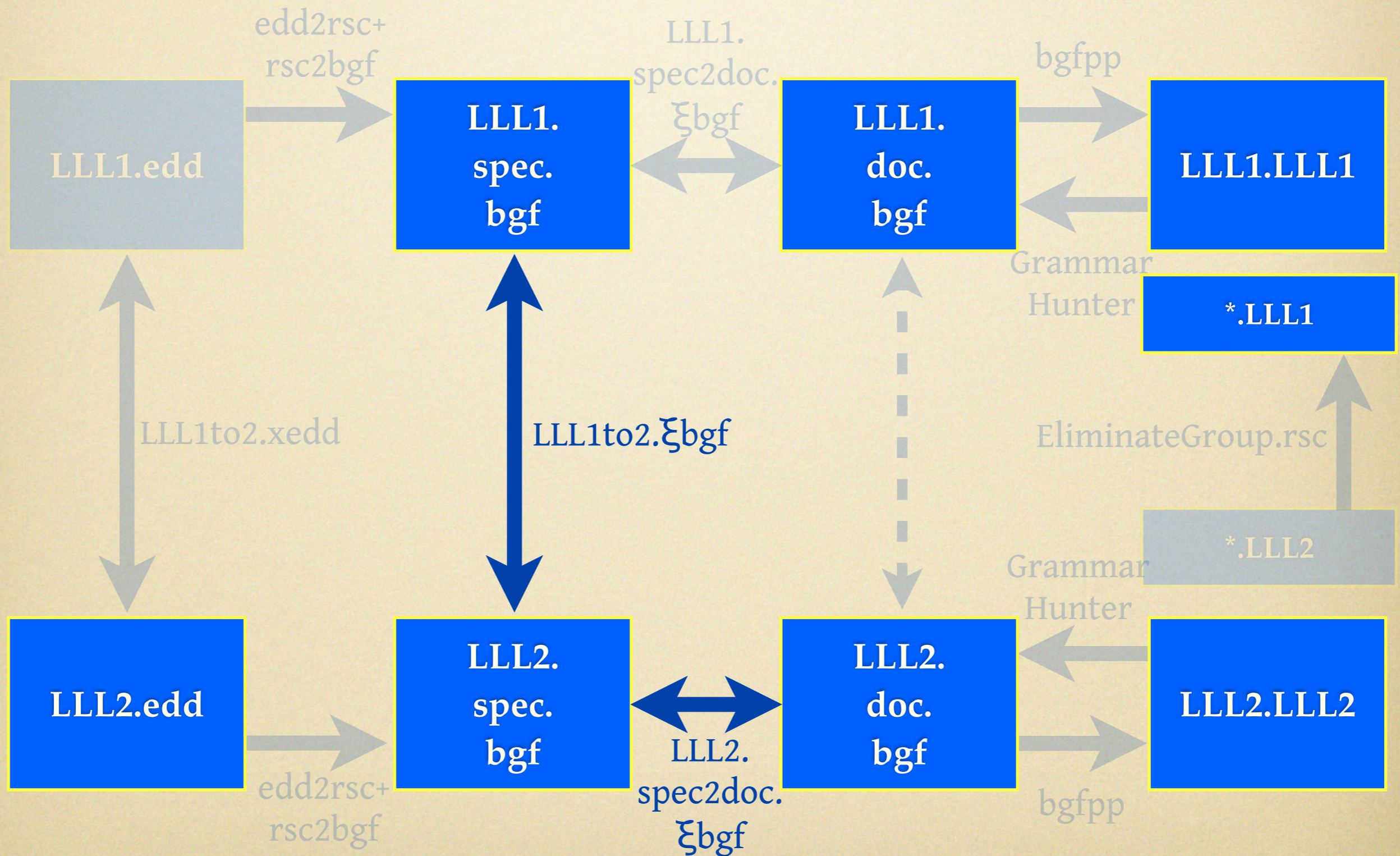
LLL2 in itself

```
specification : rule+;
rule          : ident ":" disjunction ";" ;
disjunction  : { conjunction "|" }+;
conjunction  : term*;
term         : basis repetition?;
basis        : ident | literal
              | alternation | group;
repetition   : "+" | "*" | "?";
alternation  : "{" basis basis "}" repetition;
group        : "(" disjunction ")" ;
```

Megamodel: manually



Megamodel: automated



Applying coupled mutation eliminate-metasymbol(group) to Grammar Zoo

ada-kellogg	108	csharp-iso-23270-2003	0	java-1-jls-read	0
ada-kempe	89	csharp-iso-23270-2006	0	java-2-jls-impl	36
ada-laemmel-verhoef	79	csharp-msft-ls-1.0	0	java-2-jls-read	0
ada-lncs-2219	89	csharp-msft-ls-1.2	0	java-5-habelitz	65
ada-lncs-4348	109	csharp-msft-ls-3.0	0	java-5-jls-impl	60
c-iso-9899-1999	0	csharp-msft-ls-4.0	0	java-5-jls-read	1
c-iso-9899-tc2	0	csharp-zaytsev	23	java-5-parr	95
c-iso-9899-tc3	0	dart-google	58	java-5-stahl	92
cpp-iso-14882-1998	0	dart-spec-0.01	56	java-5-studman	91
cpp-iso-n2723	0	dart-spec-0.05	62	mediawiki-bnf	32
csharp-ecma-334-1	0	eiffel-bezault	45	mediawiki-ebnf	30
csharp-ecma-334-2	0	eiffel-iso-25436-2006	345	modula-sdf	50
csharp-ecma-334-3	0	fortran-derricks	101	modula-src-052	65
csharp-ecma-334-4	0	java-1-jls-impl	0	w3c-xpath1	3

Conclusion

Conclusion

- Extended XBGF to bidirectionality, resulting in Ξ BGF.
- Proposed EDD and XEDD for notation & its evolution.
- Presented a case study of LLL evolution (GDK).
- Generalised transformers and generators to transformations and mutations; also formalised them.
- Implemented an XEDD processor for evolution, coevolution, change propagation and mutation.

Discussion

grammarware.net

slps.sf.net